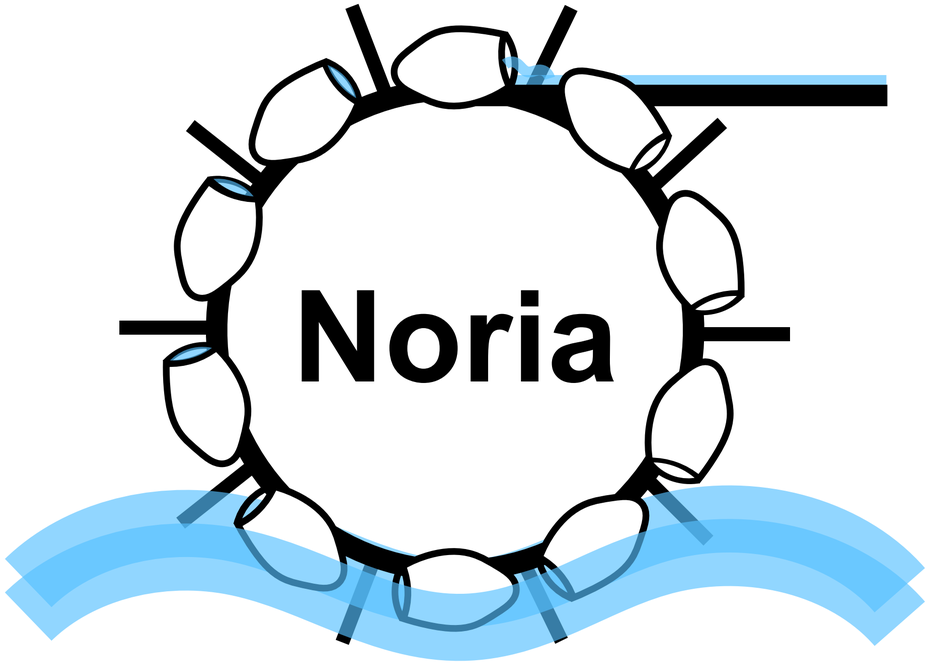


Dynamic, partially-stateful data-flow for high-performance Web applications

Jon Gjengset

Malte Schwarzkopf

Jonathan Behrens

Lara Timbó Araújo

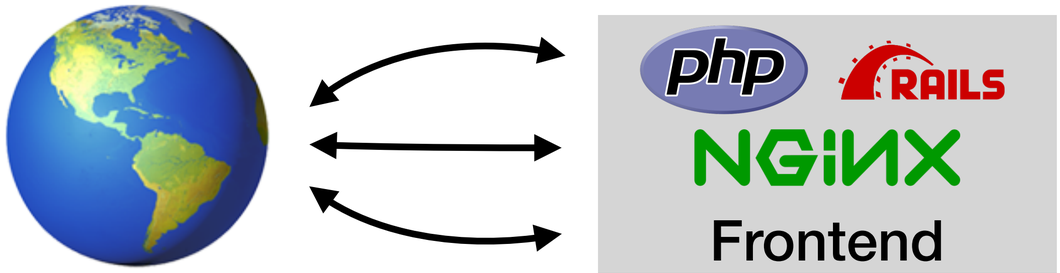
Martin Ek

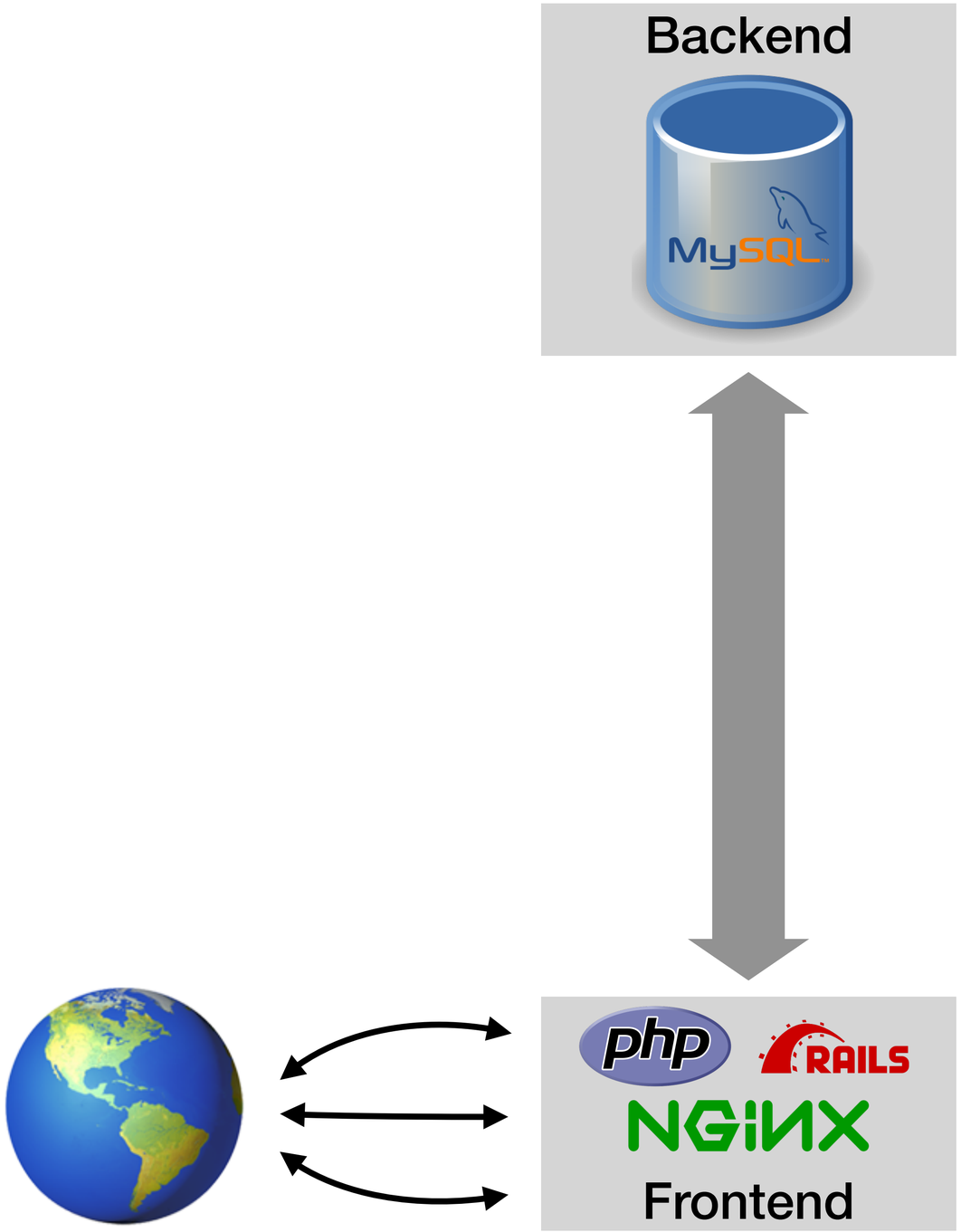
Eddie Kohler

M. Frans Kaashoek

Robert Morris


php
RAILS
NGINX
Frontend

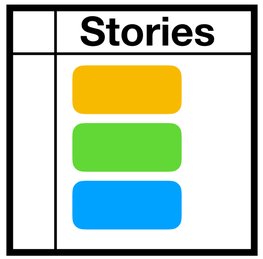

php
RAILS
NGINX
Frontend

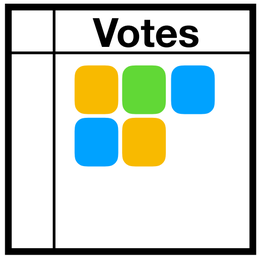

Backend
MySQL
php
RAILS
NGINX
Frontend

Backend

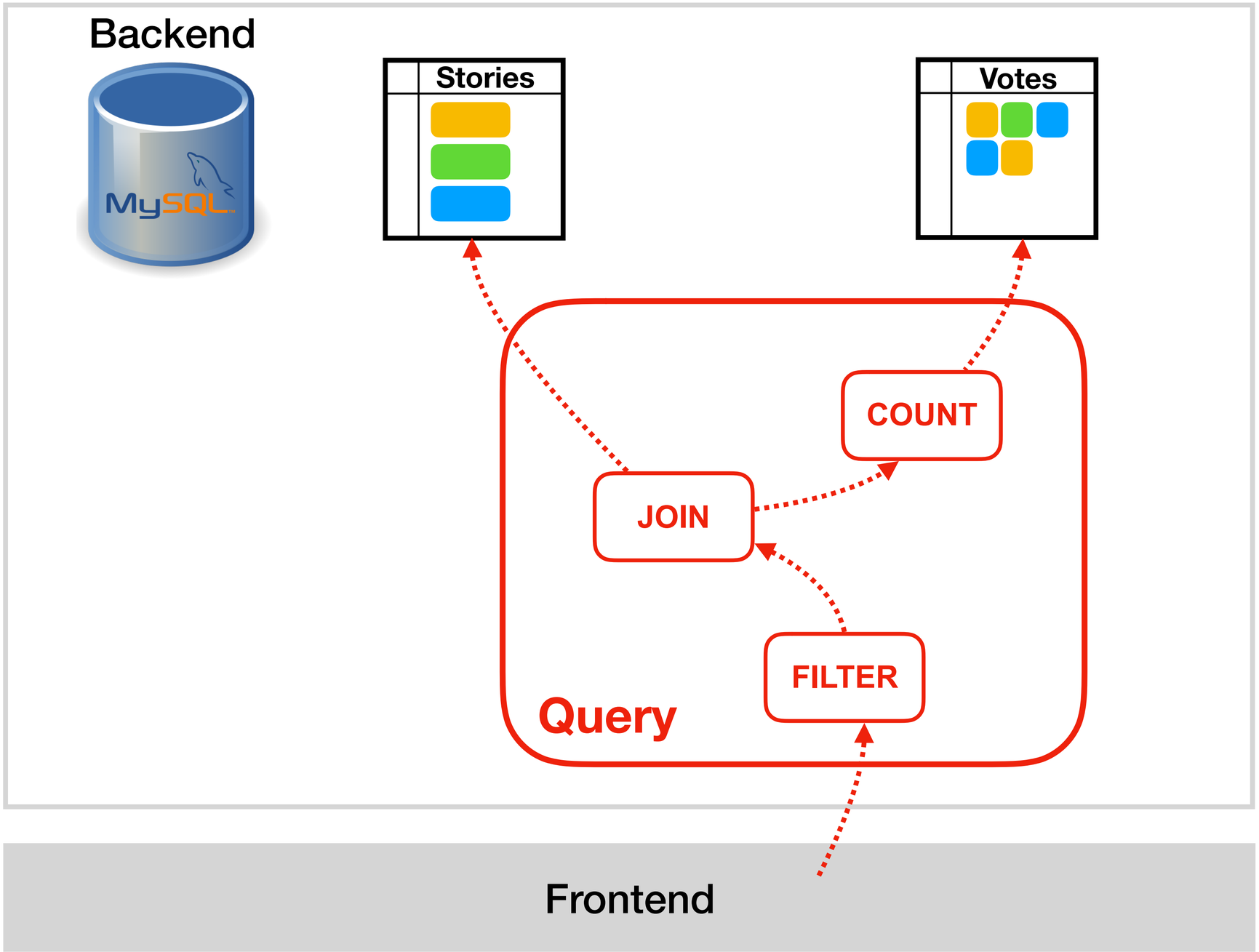
Frontend

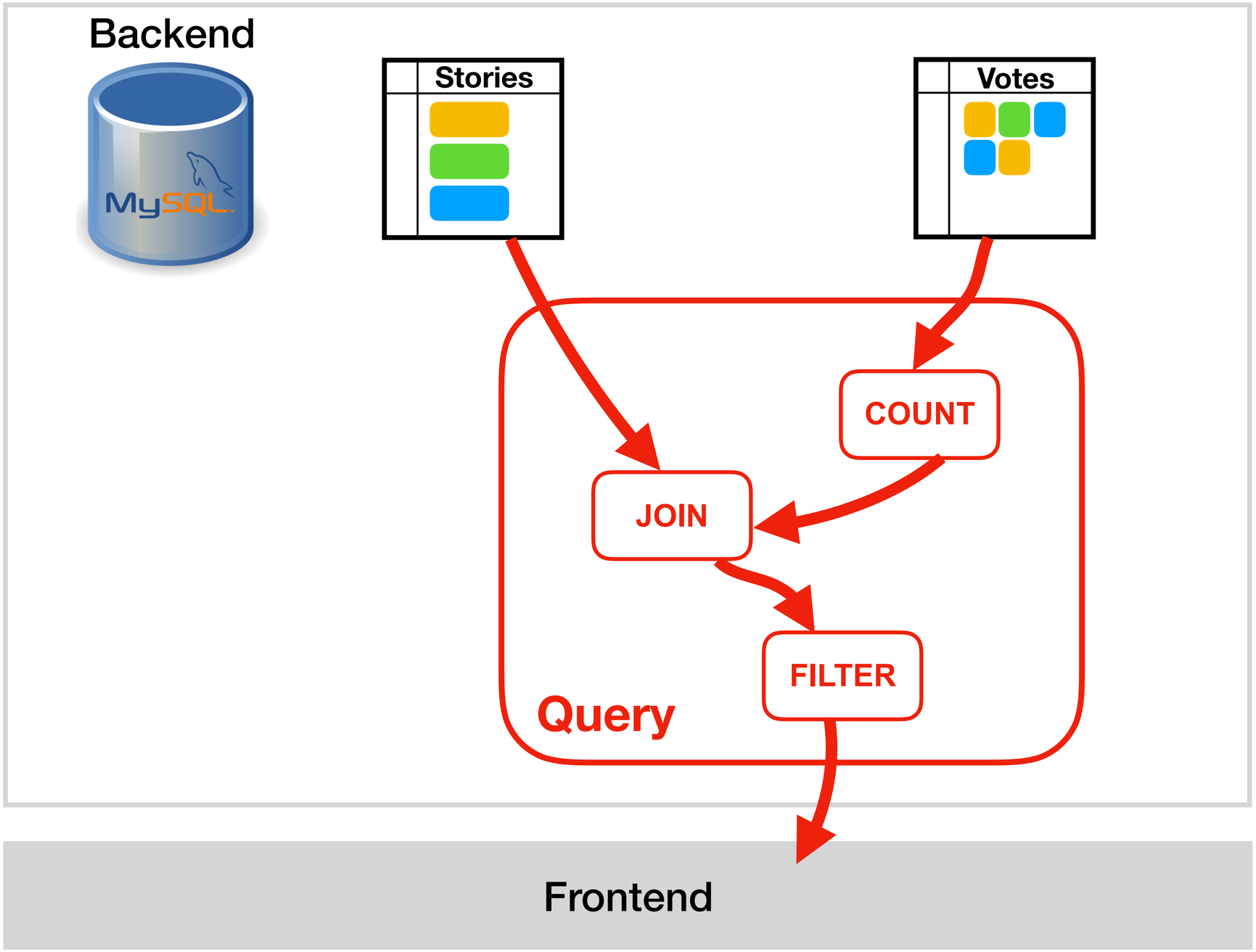
Backend





Frontend


Backend
MySQL
Stories
Votes
COUNT
JOIN
FILTER
Query
Frontend


Backend
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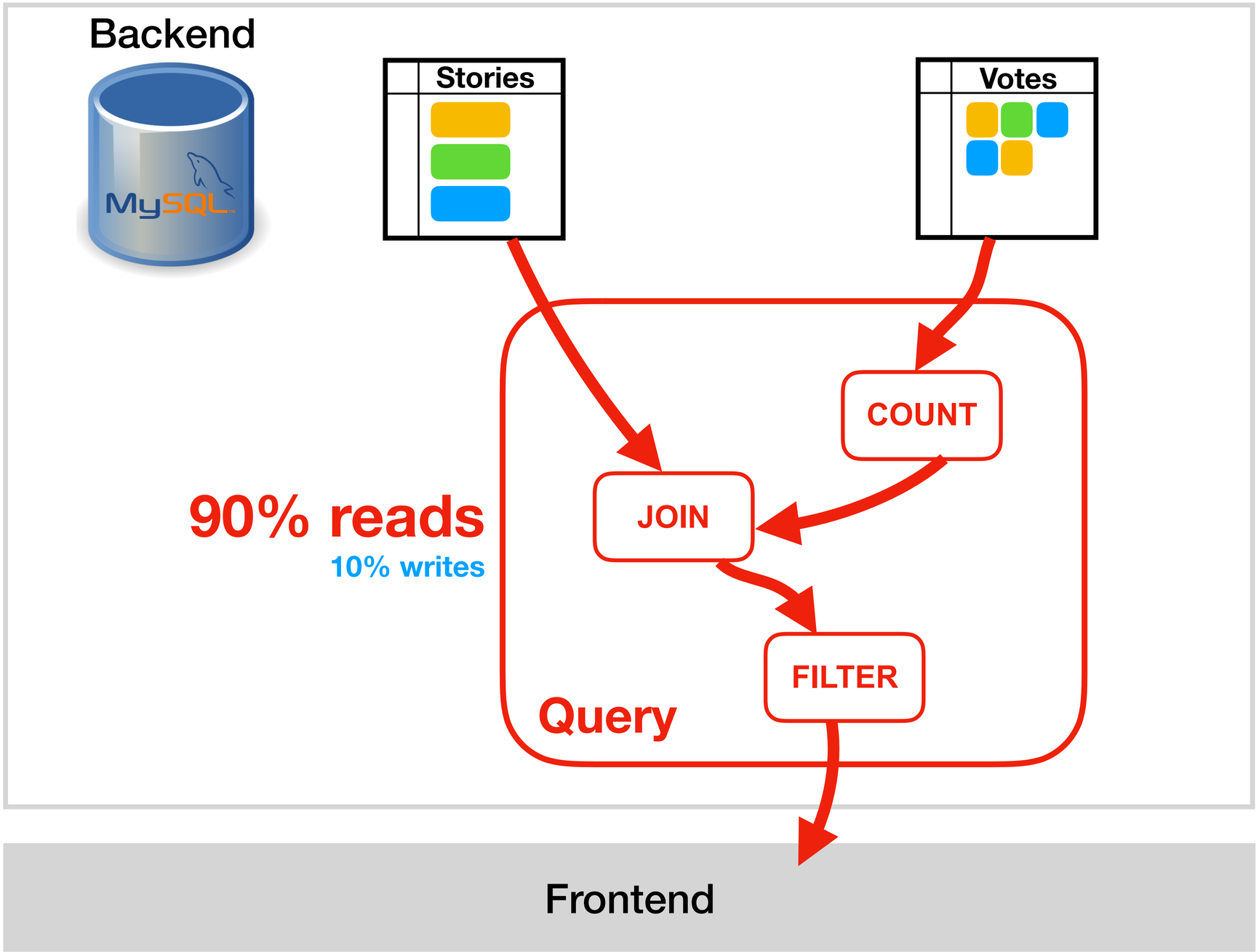
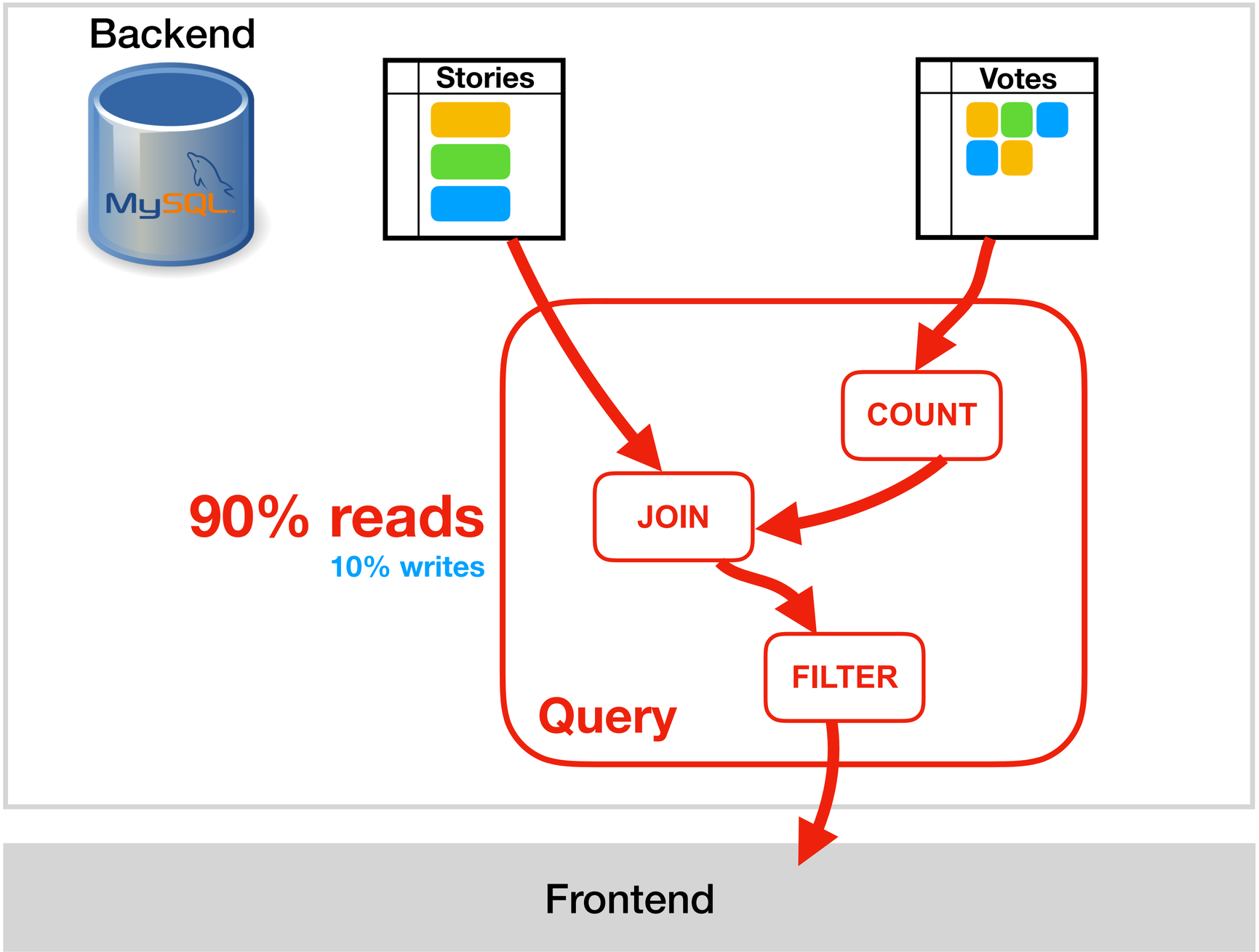
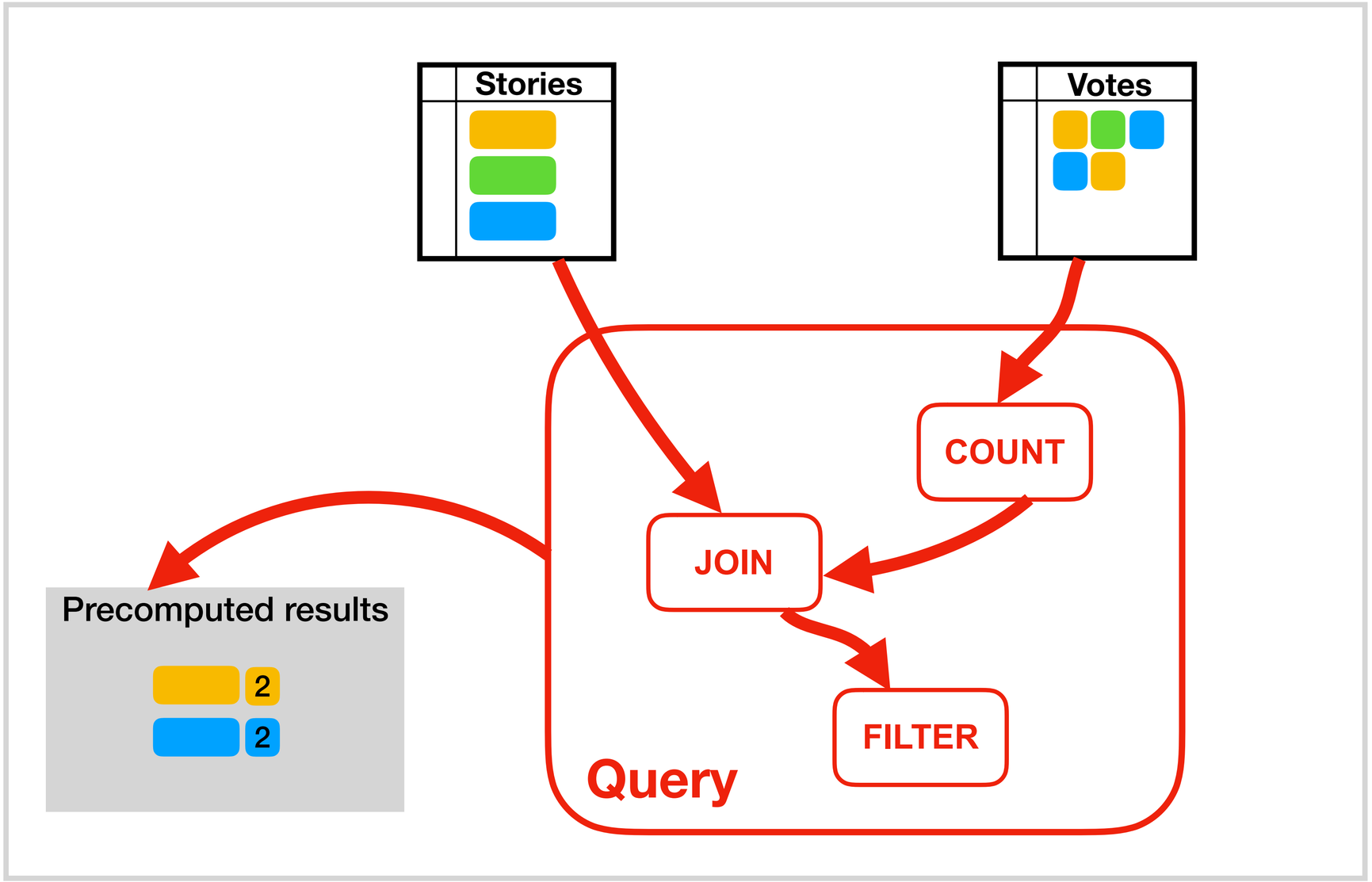
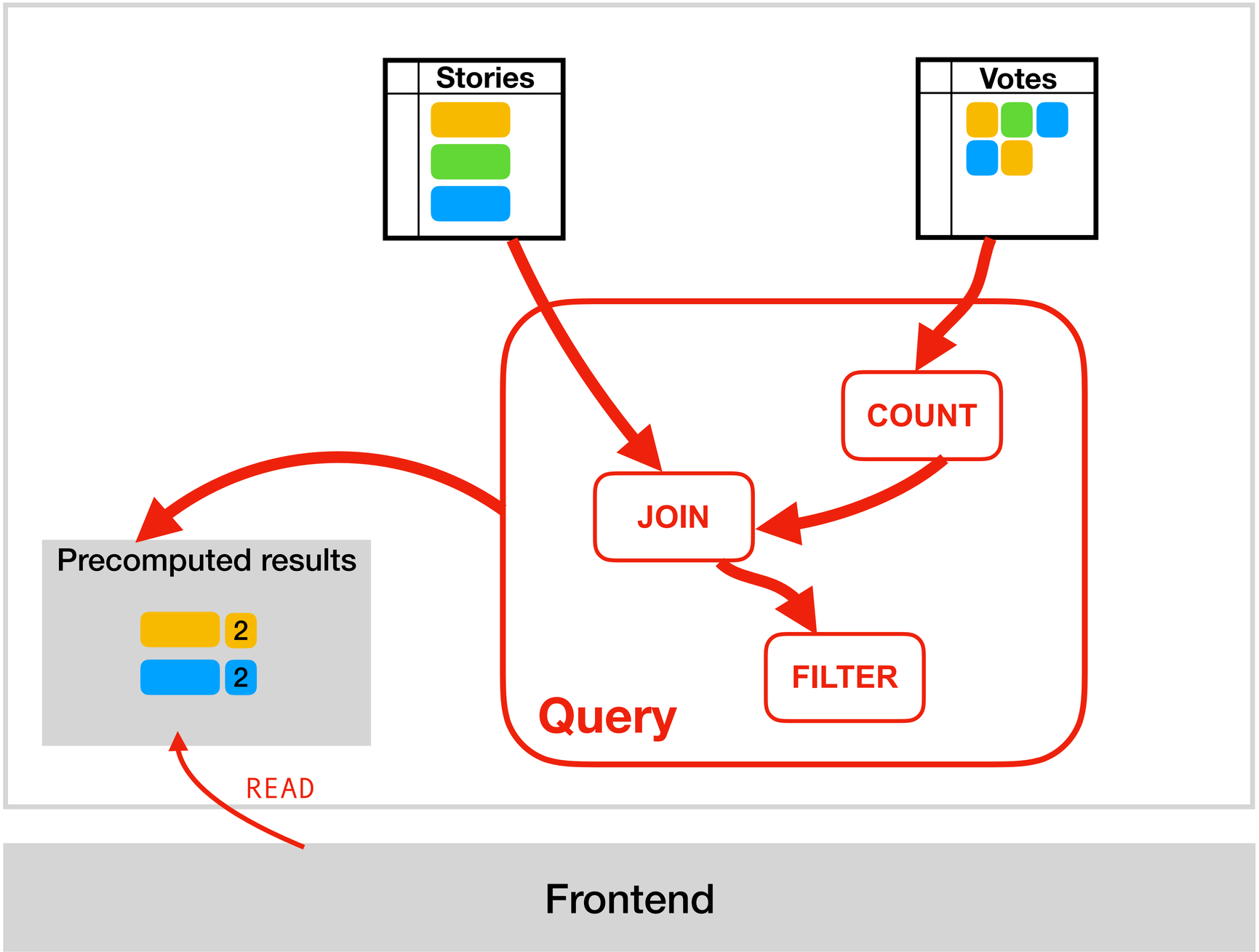

Backend
MySQL
Stories
Votes
COUNT
JOIN
FILTER
Query
90% reads
10% writes
Frontend

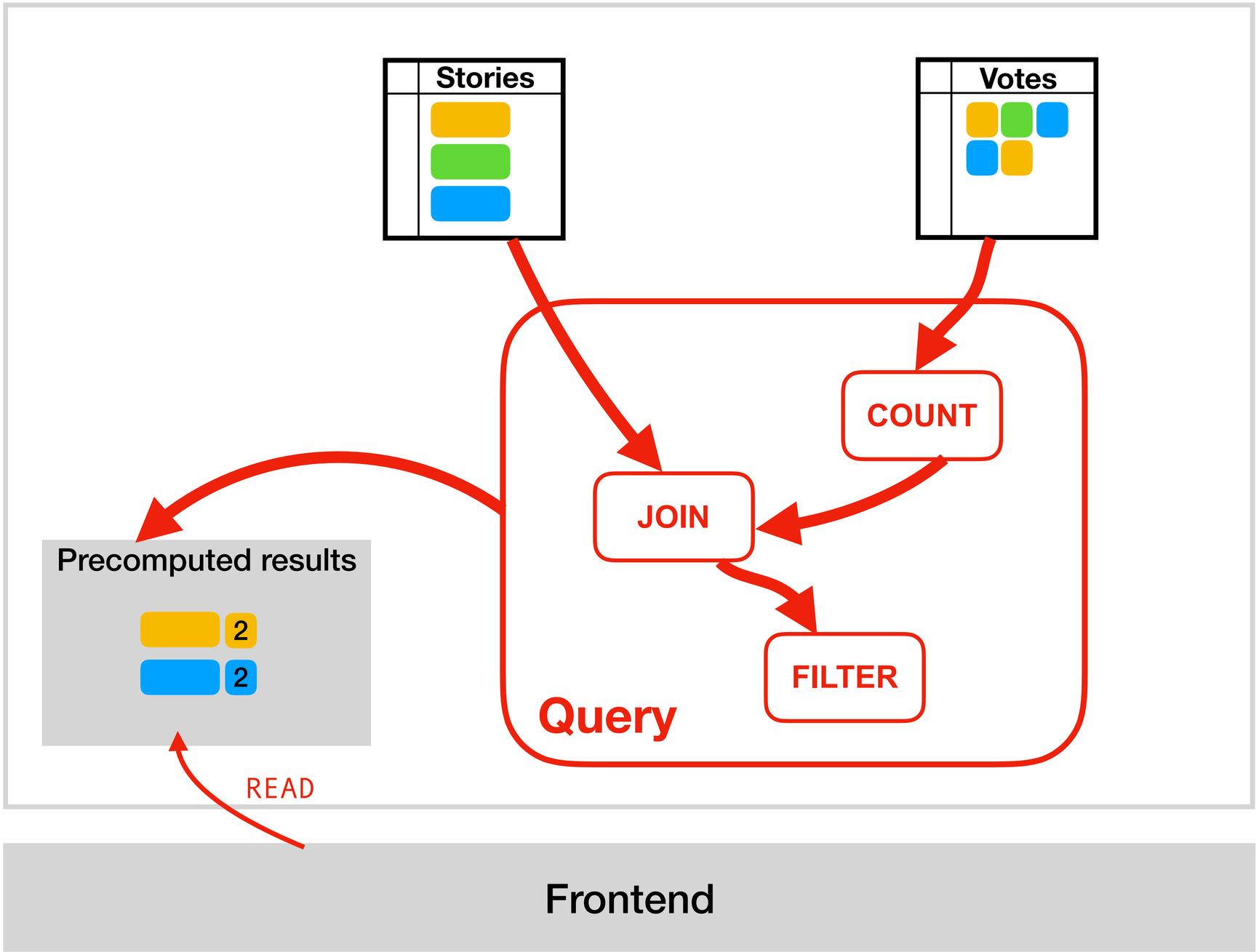
Slow reads, repeated work!




Stories
Votes
COUNT
JOIN
FILTER
Query
Precomputed results
2
2
Frontend

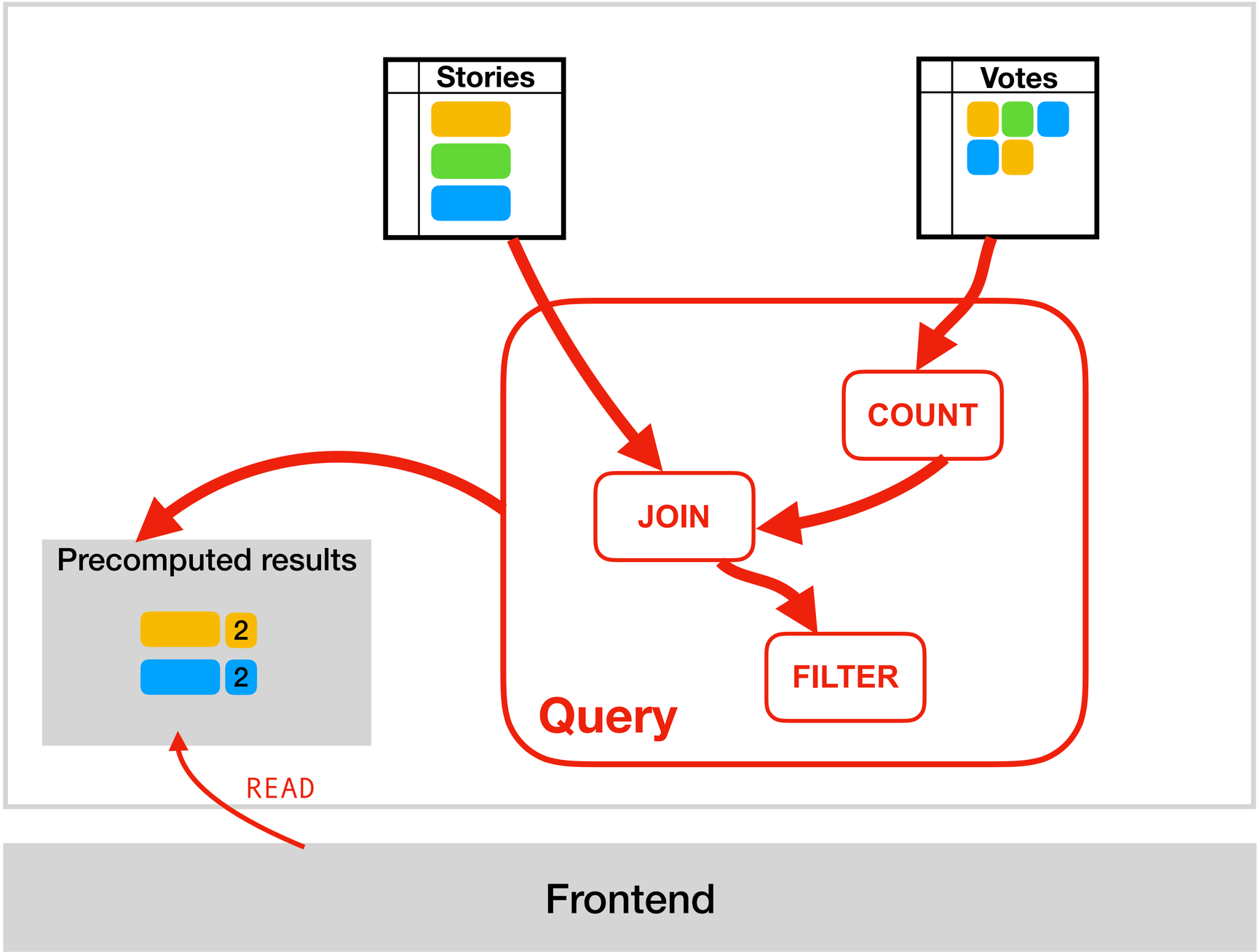

Stories
Votes
COUNT
JOIN
FILTER
Query
Precomputed results
2
2
READ
Frontend

Store in base table?
— **manual, slow**.



Store in base table?
— **manual, slow**.

memcached?
— **complex**
[Facebook NSDI'13].



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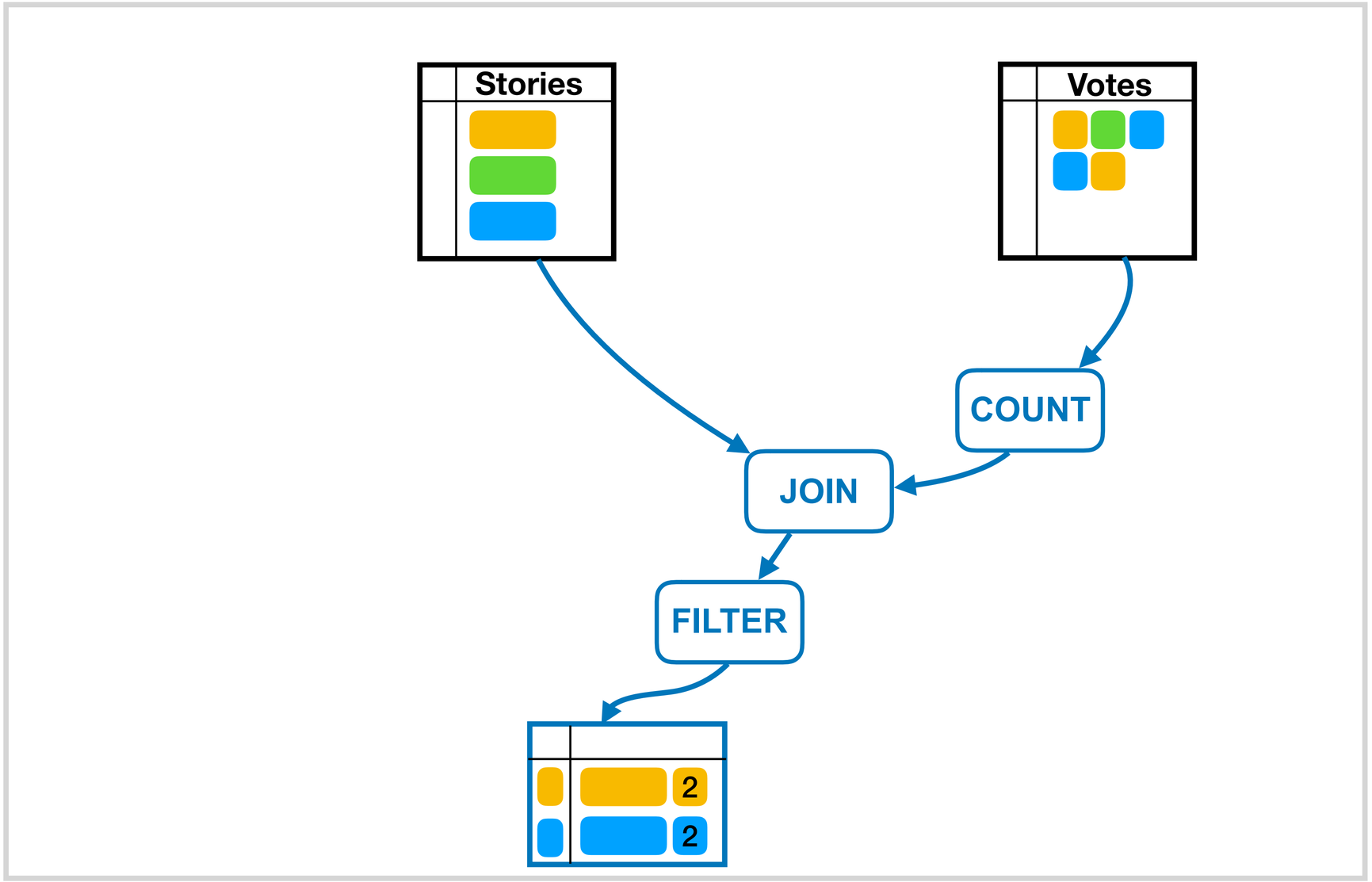
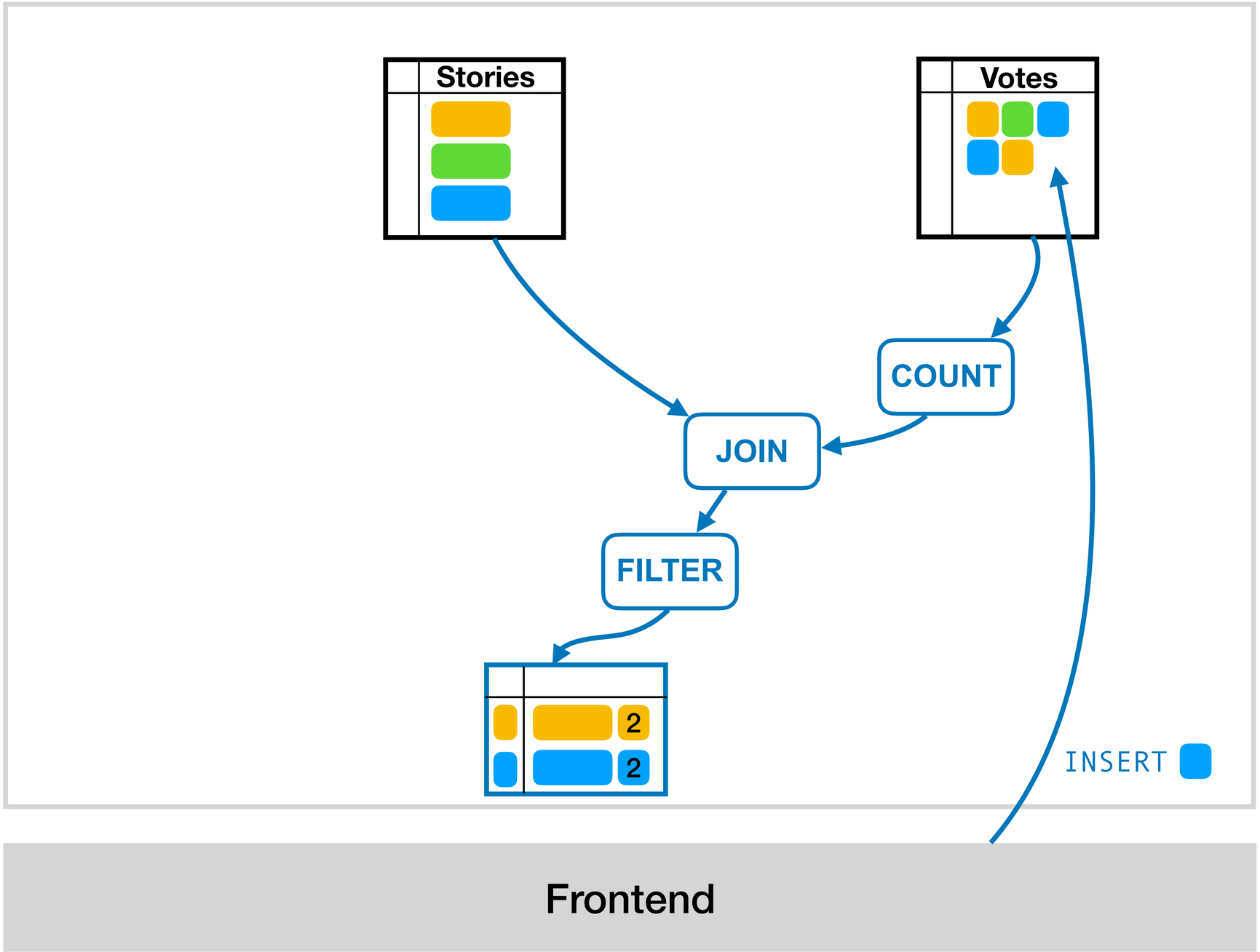
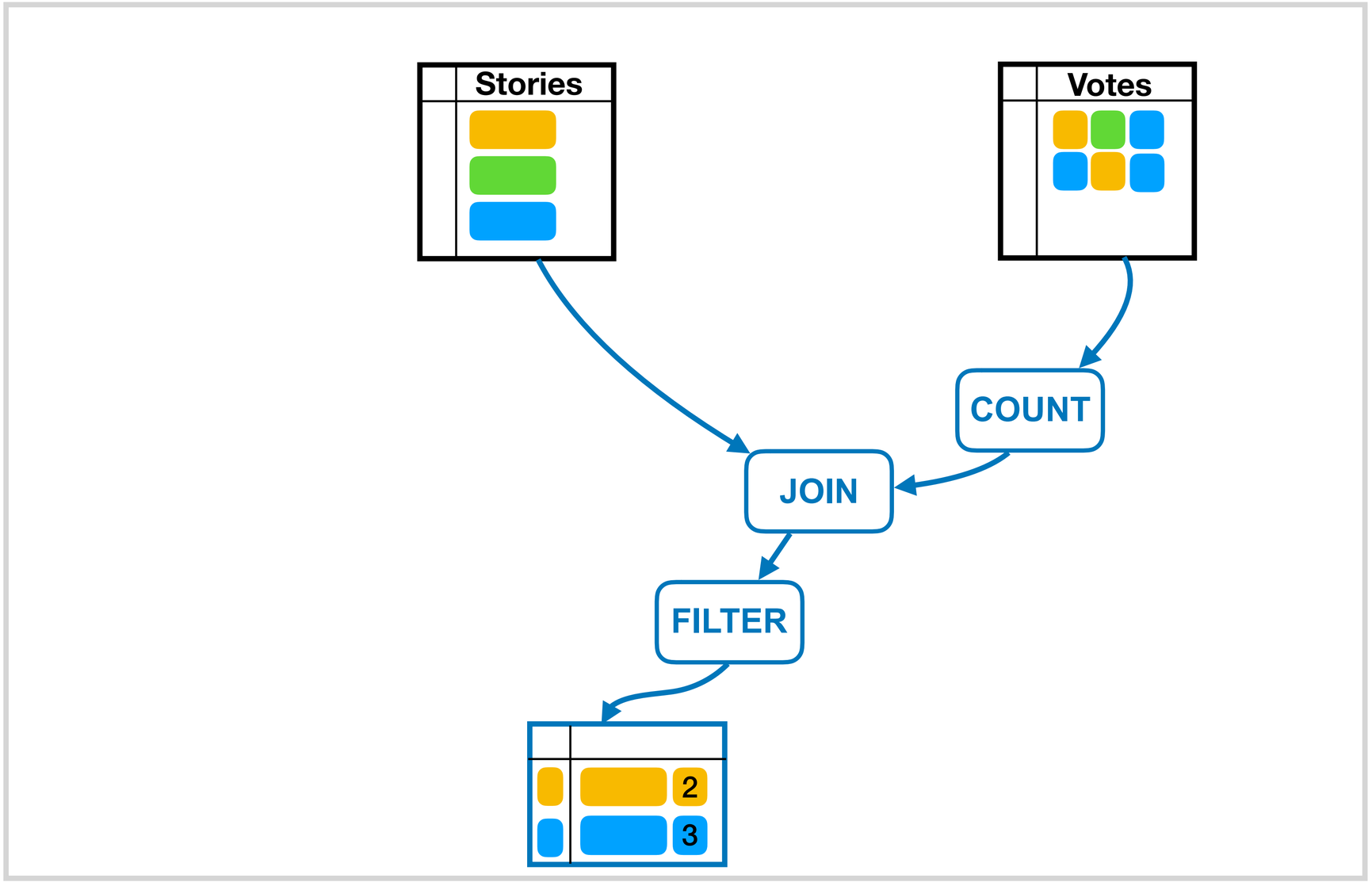
Streaming data-flow?



Streaming data-flow?

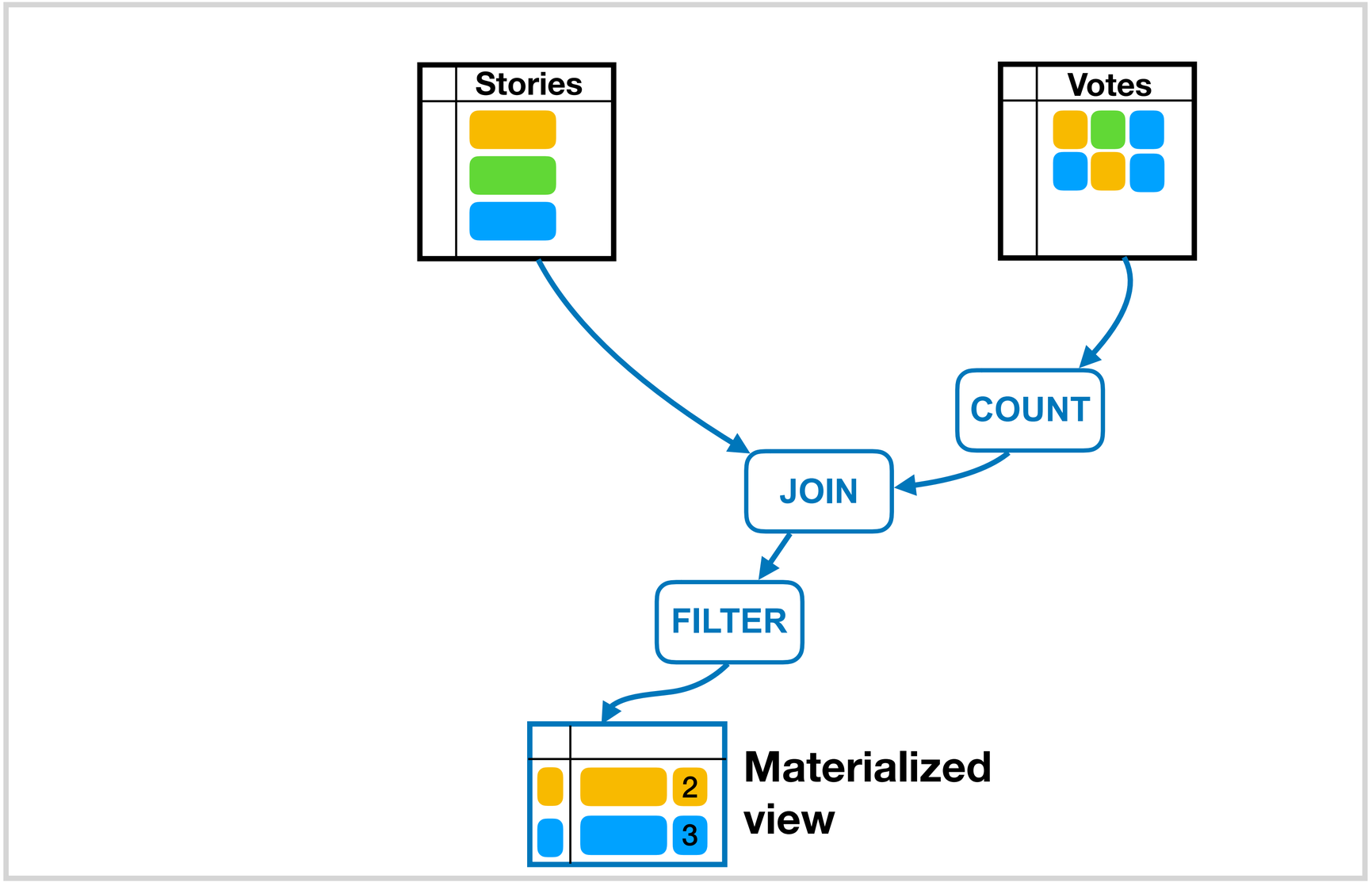


Streaming data-flow?



Frontend

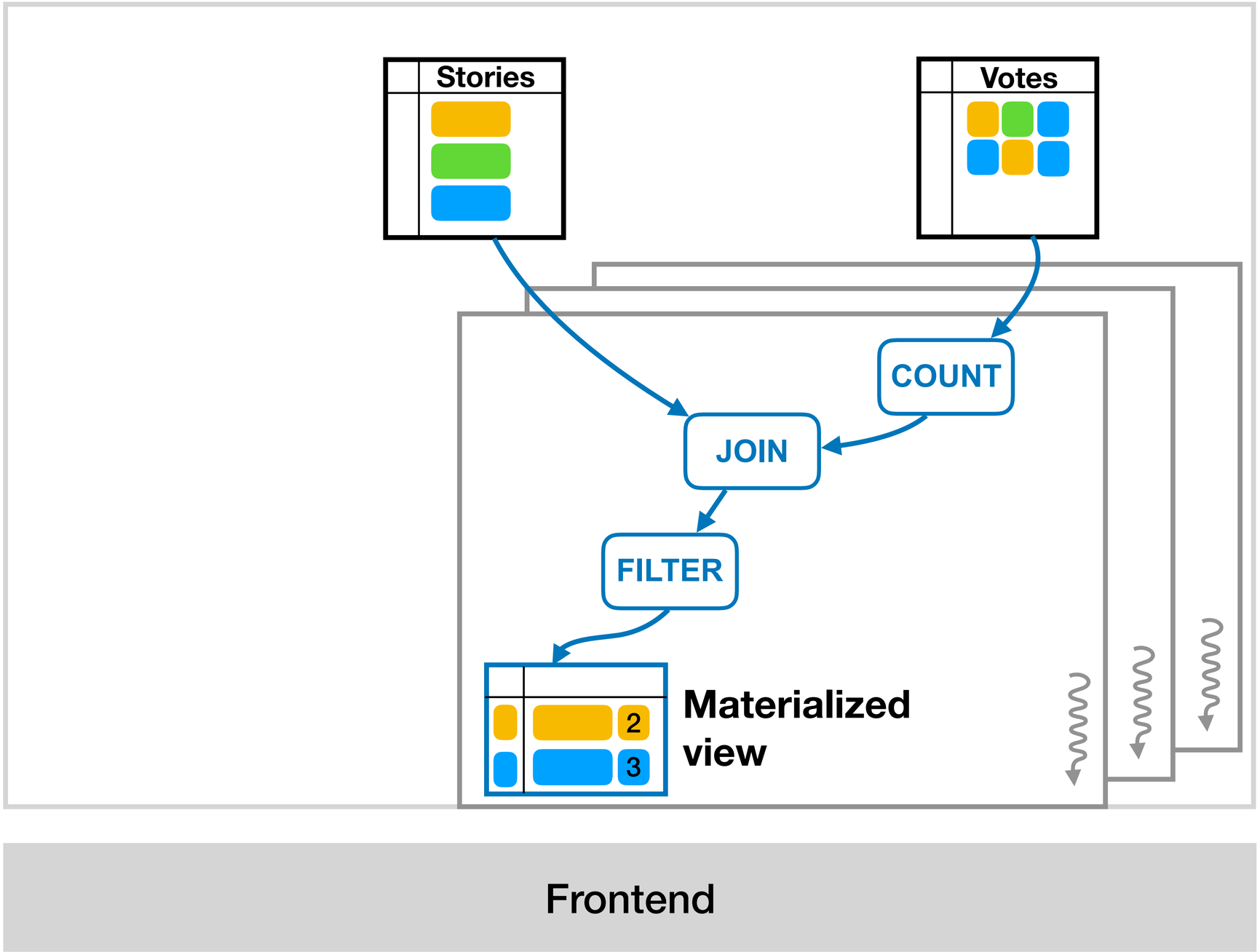
Streaming data-flow?



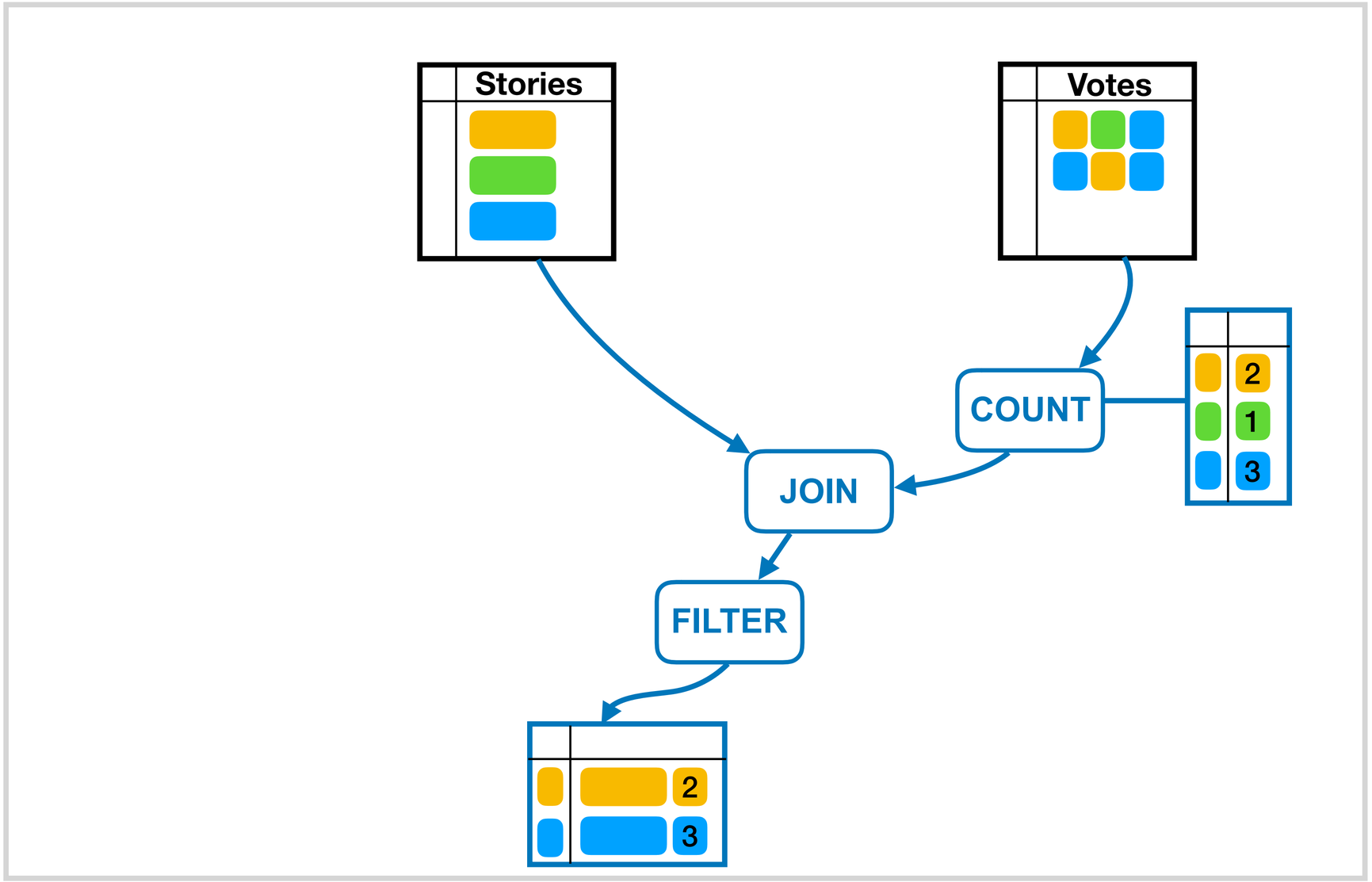
Frontend

Streaming data-flow?

Fast reads.
Efficient writes.
Parallelizes well.



Challenges

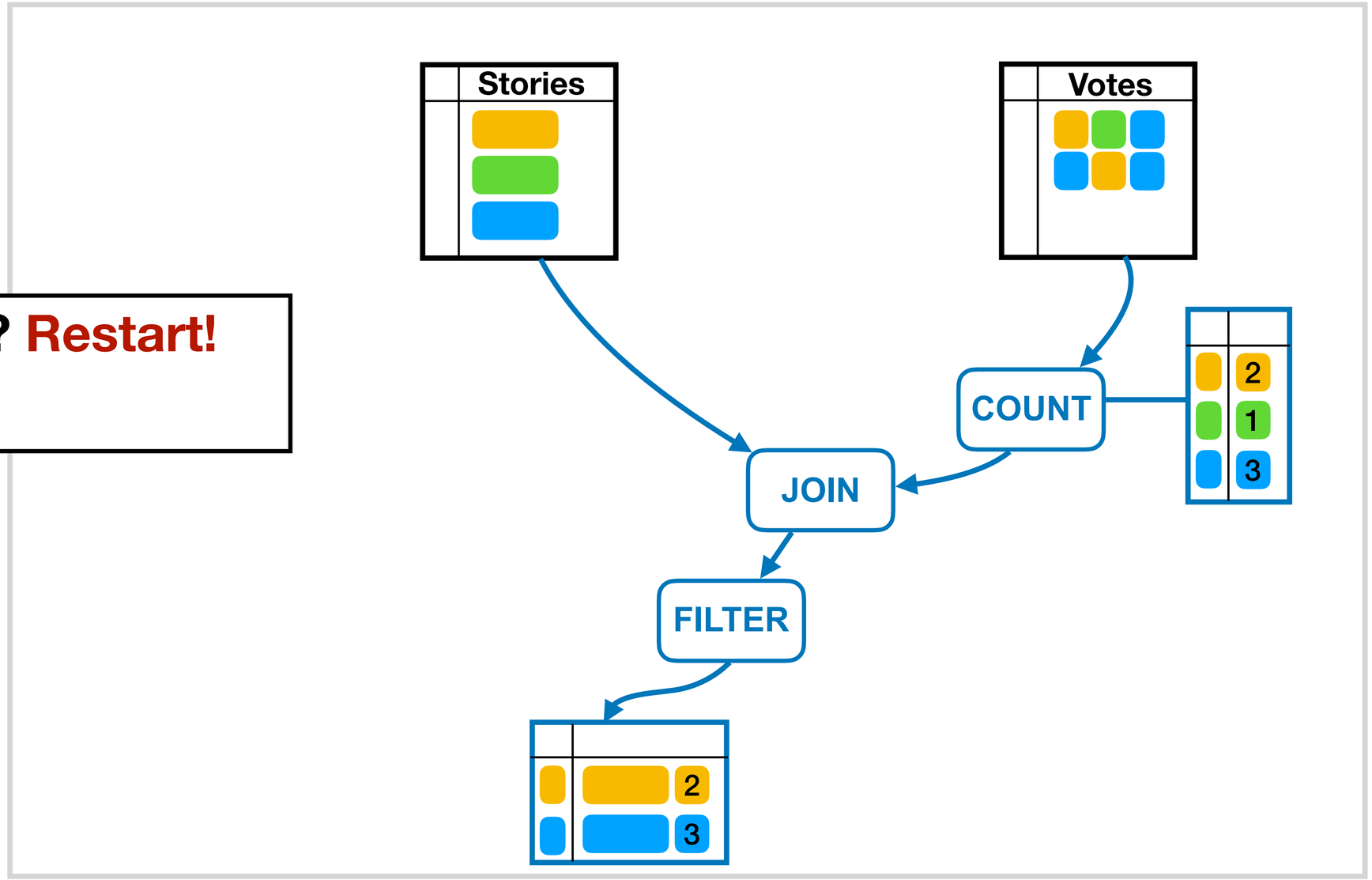


Frontend

Challenges

State-of-the-art data-flow systems:

- **Change queries? Restart!**



Challenges

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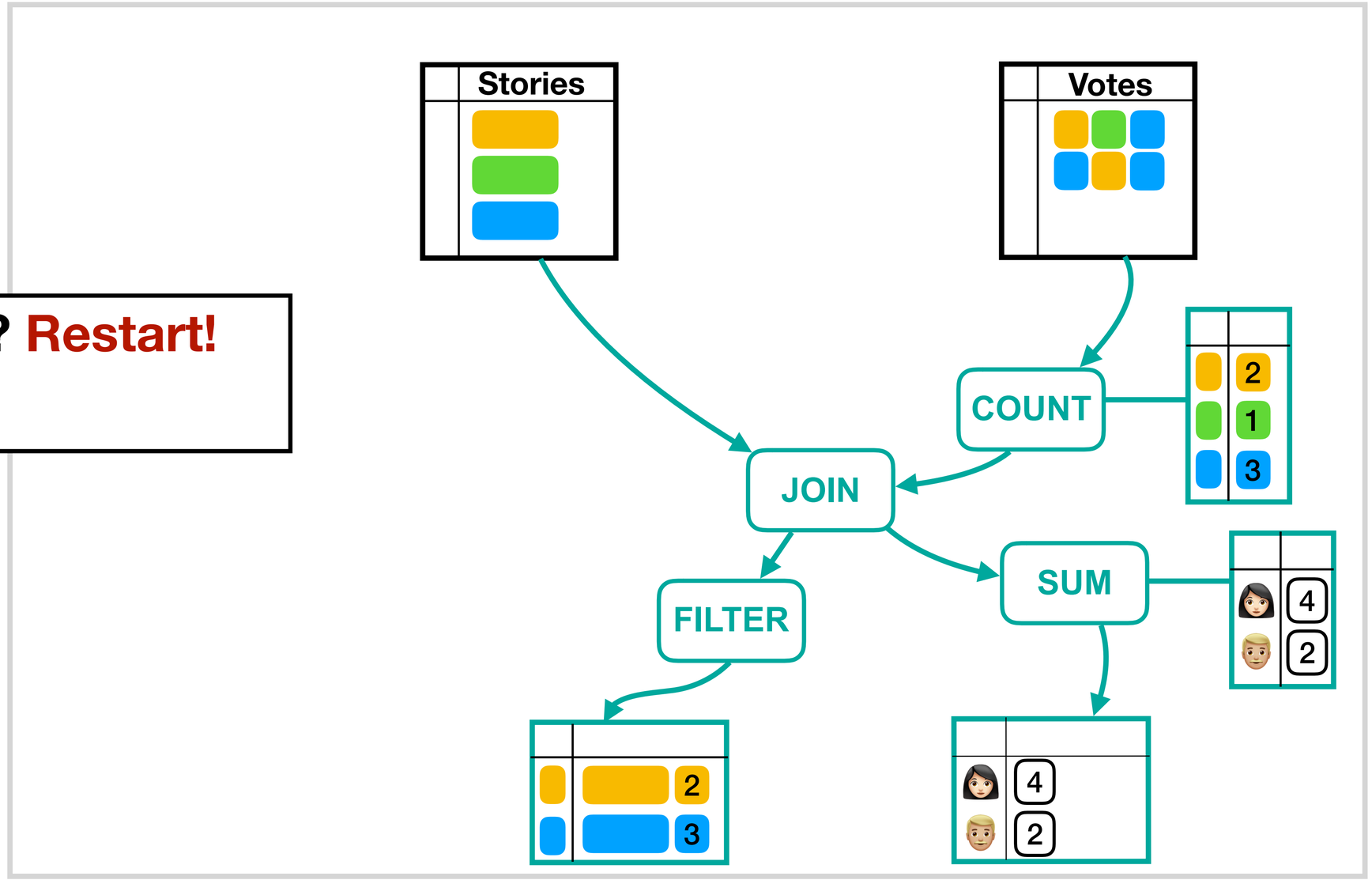
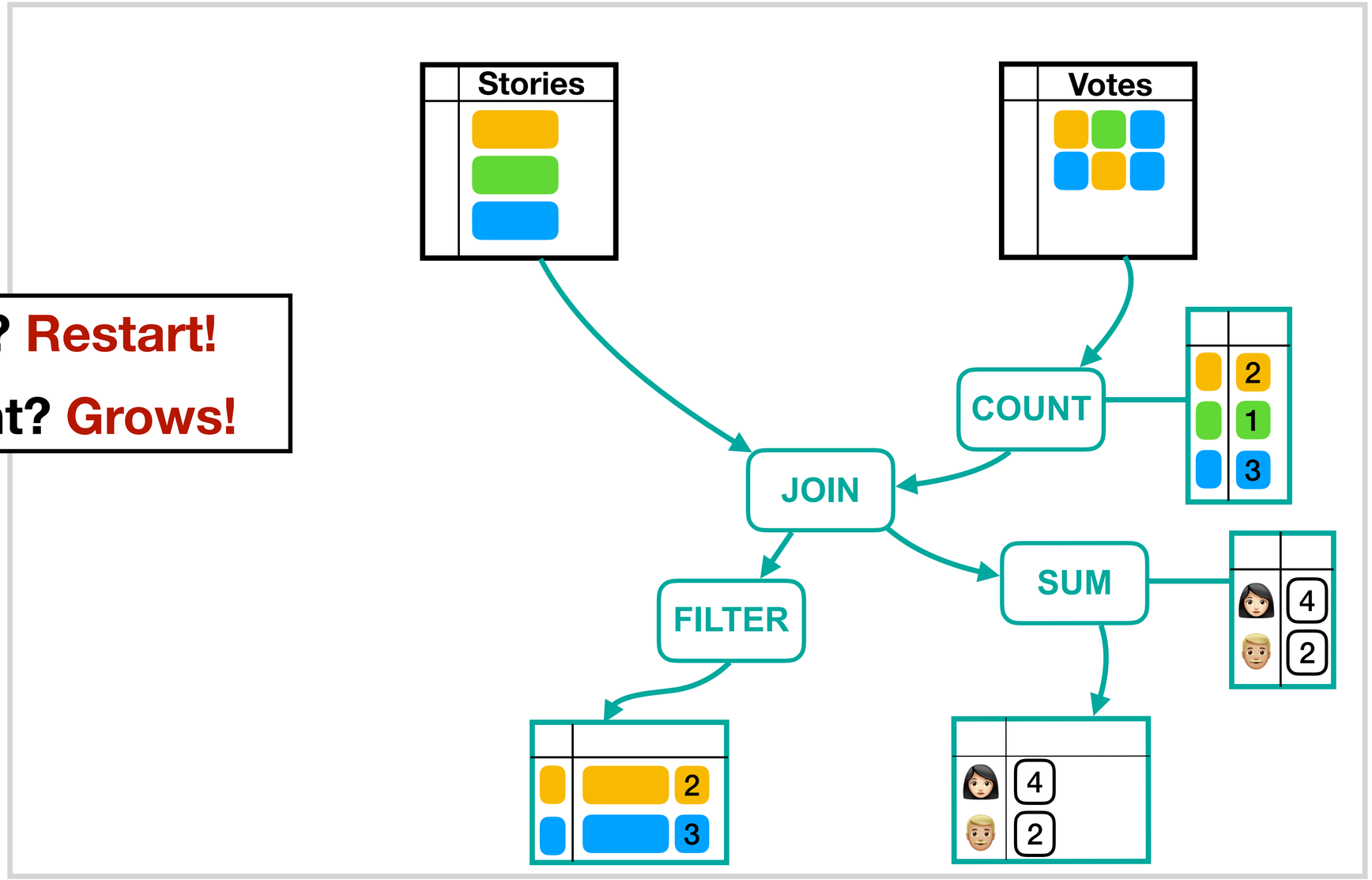


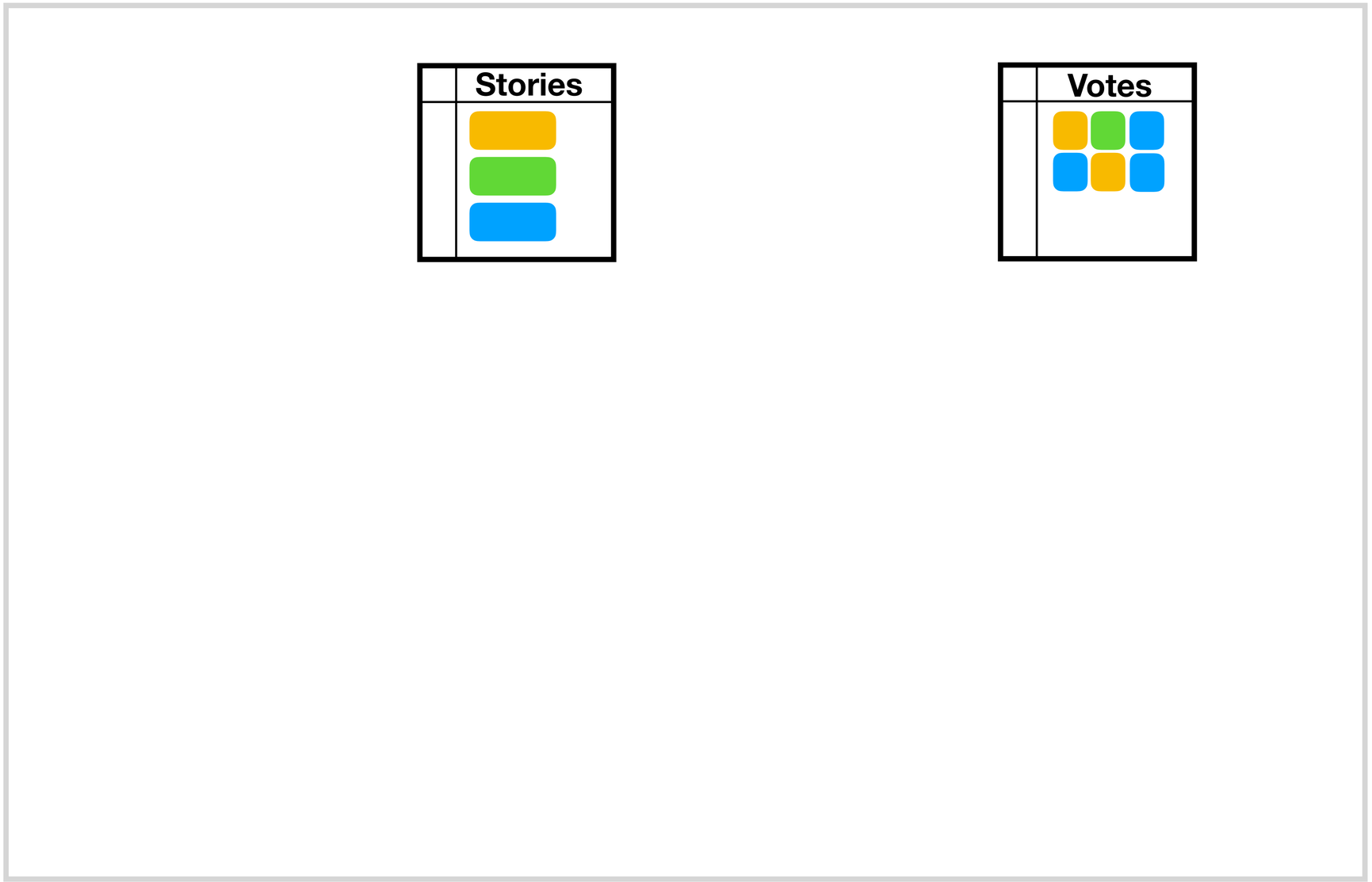
Challenges

State-of-the-art data-flow systems:

- **Change queries? Restart!**
- **Memory footprint? Grows!**



Noria



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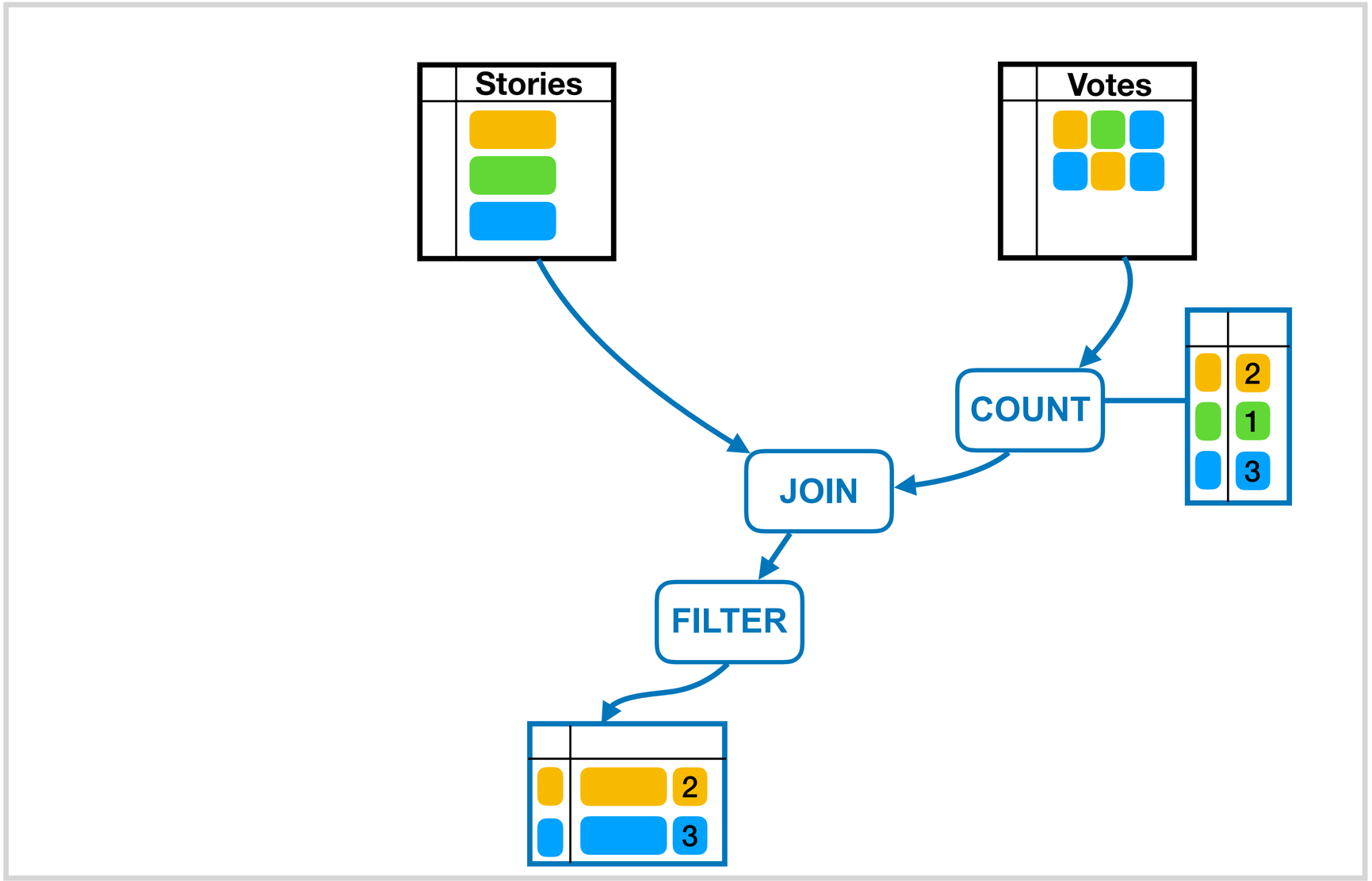
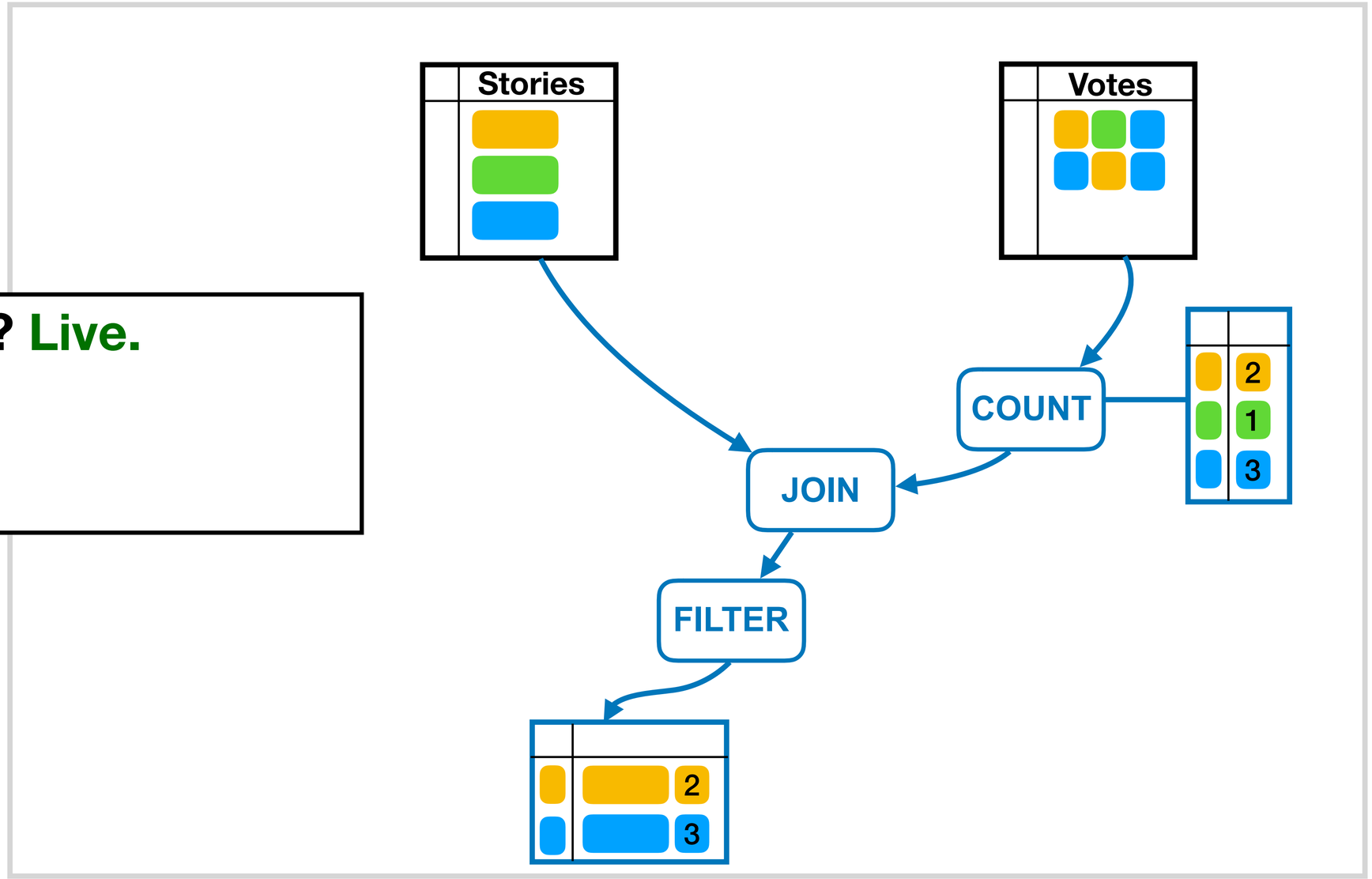


Frontend

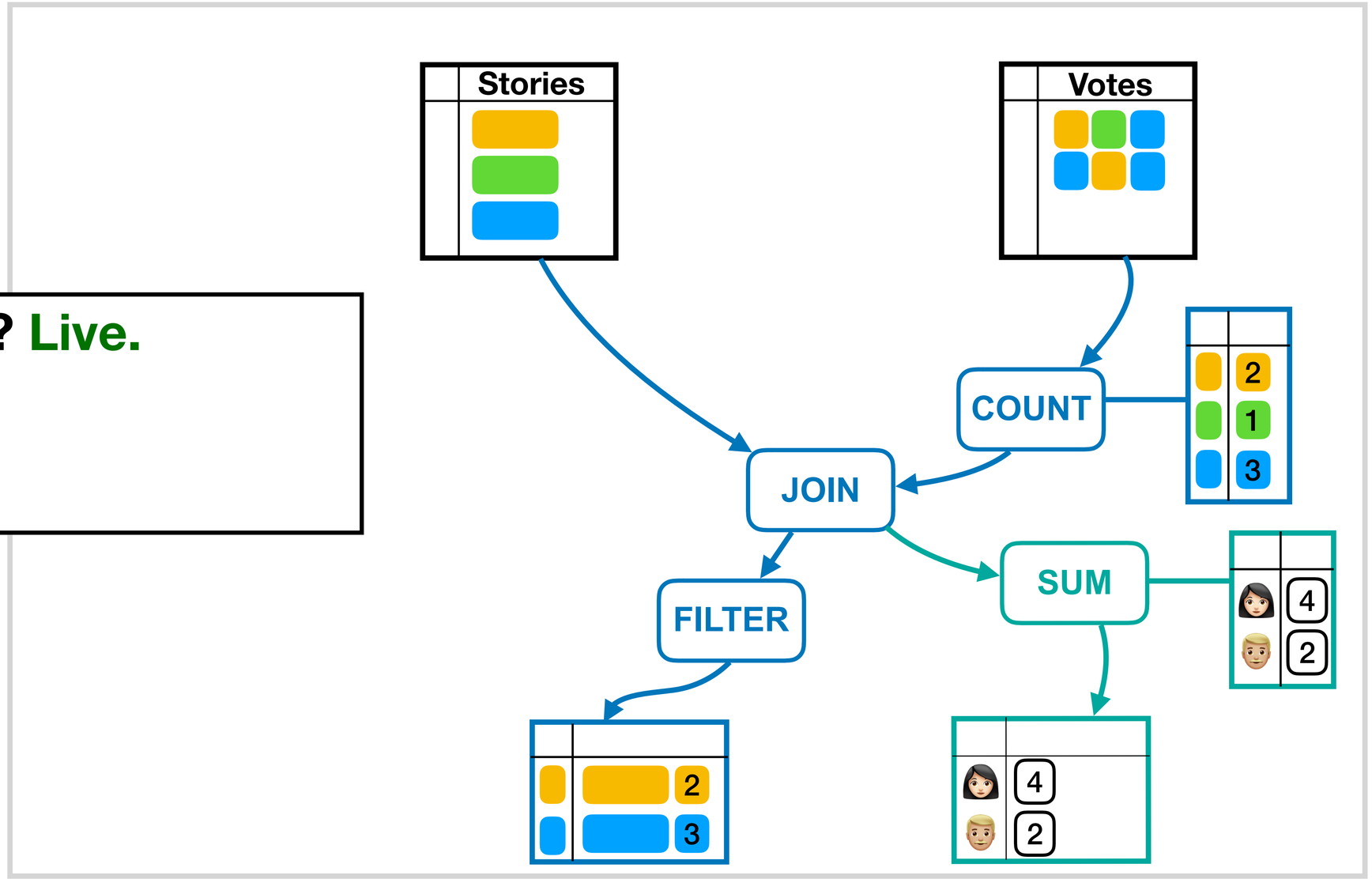
Noria

- Change queries? Live.



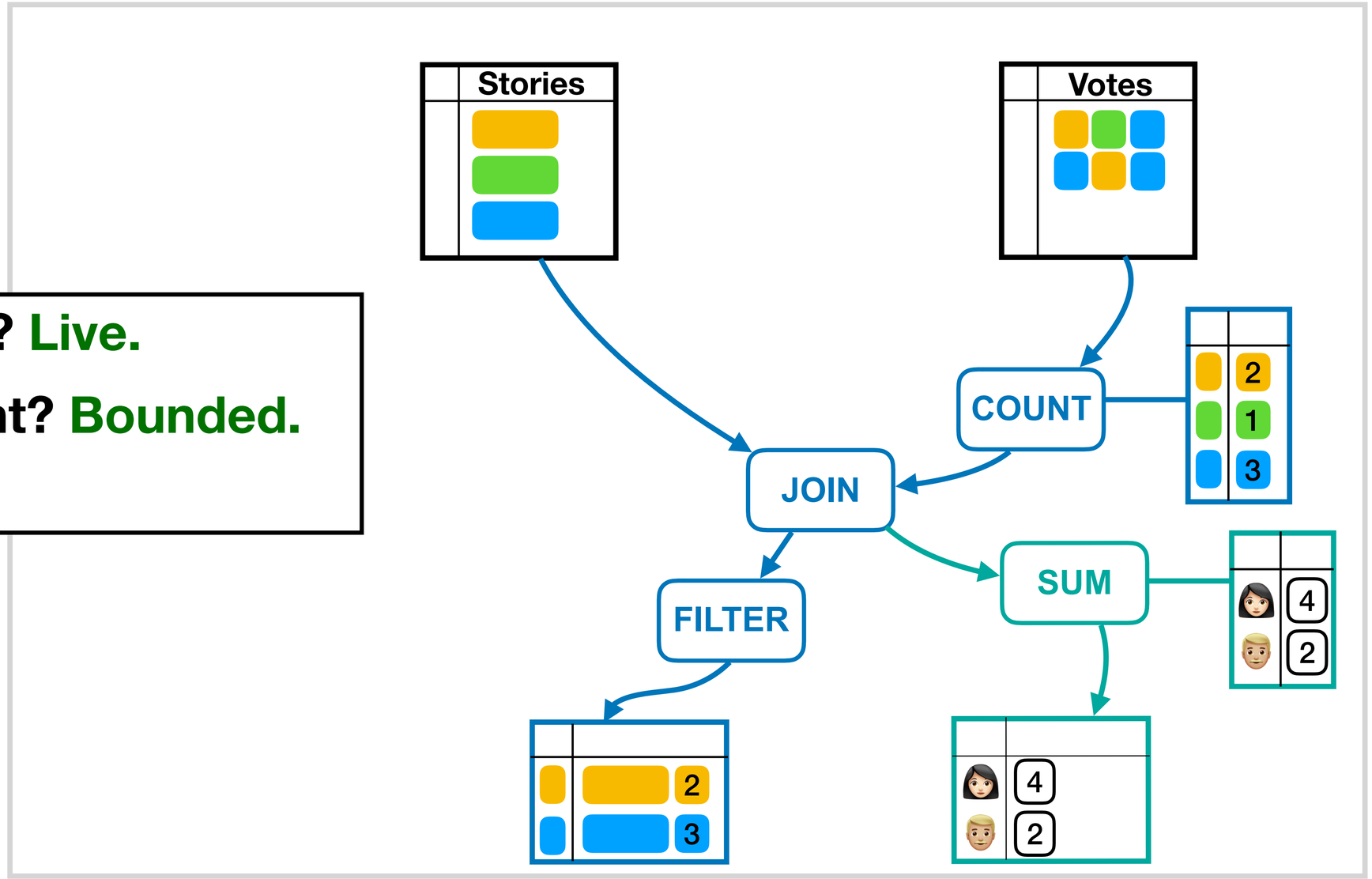
Noria

- Change queries? Live.



Noria

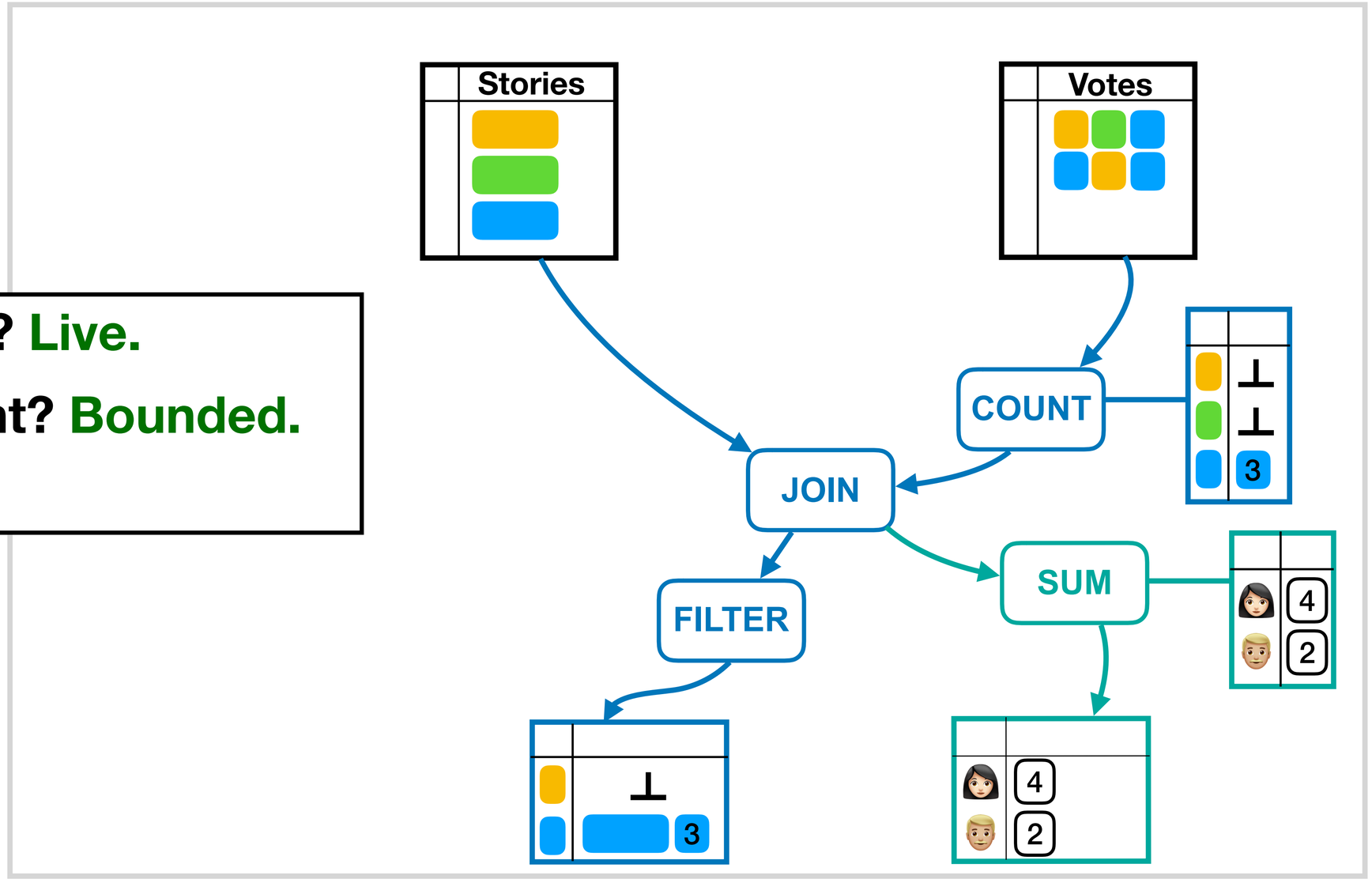
- Change queries? **Live.**
- Memory footprint? **Bounded.**



Frontend

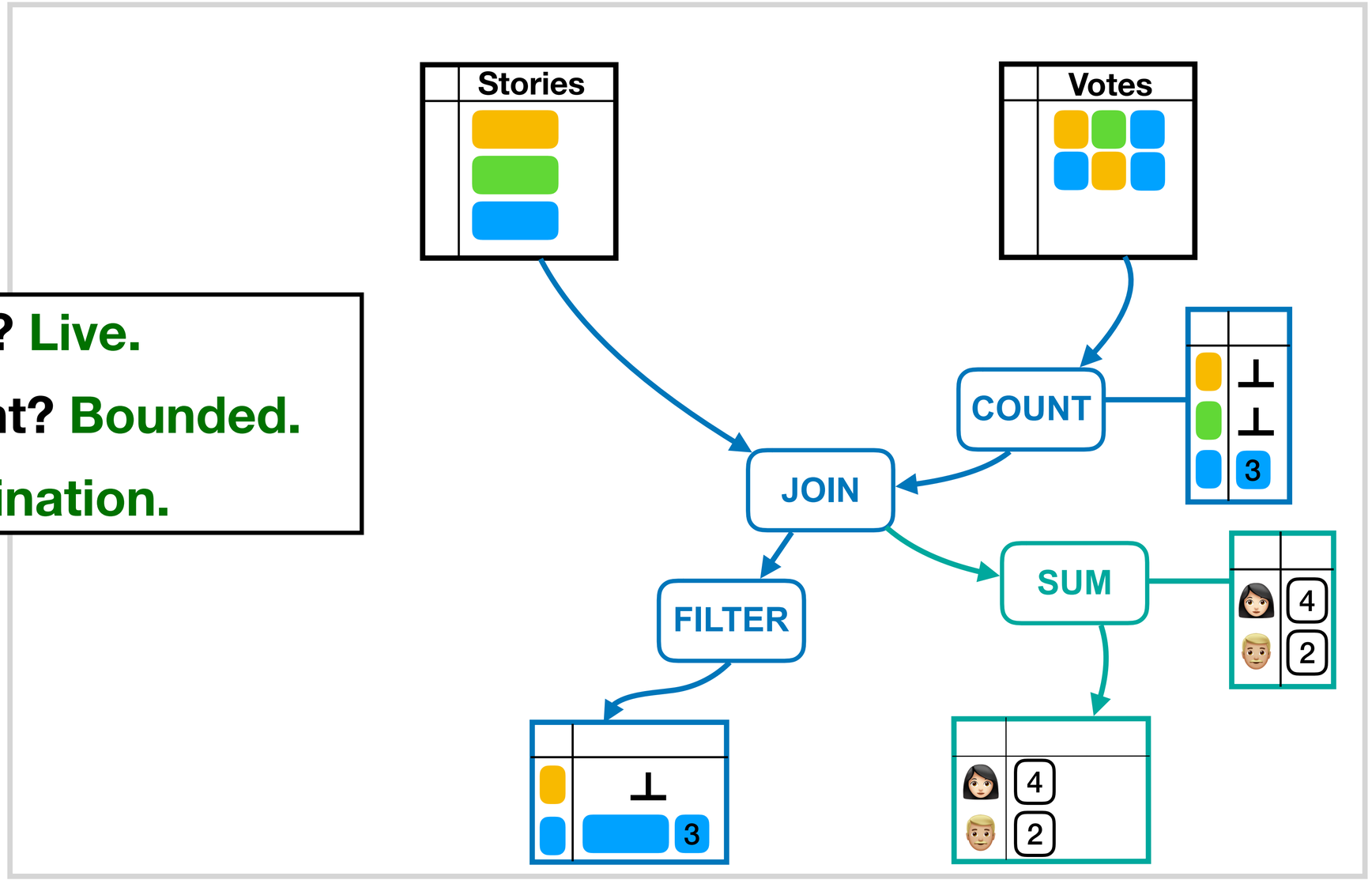
Noria

- Change queries? **Live.**
- Memory footprint? **Bounded.**



Noria

- Change queries? Live.
- Memory footprint? Bounded.
- No global coordination.

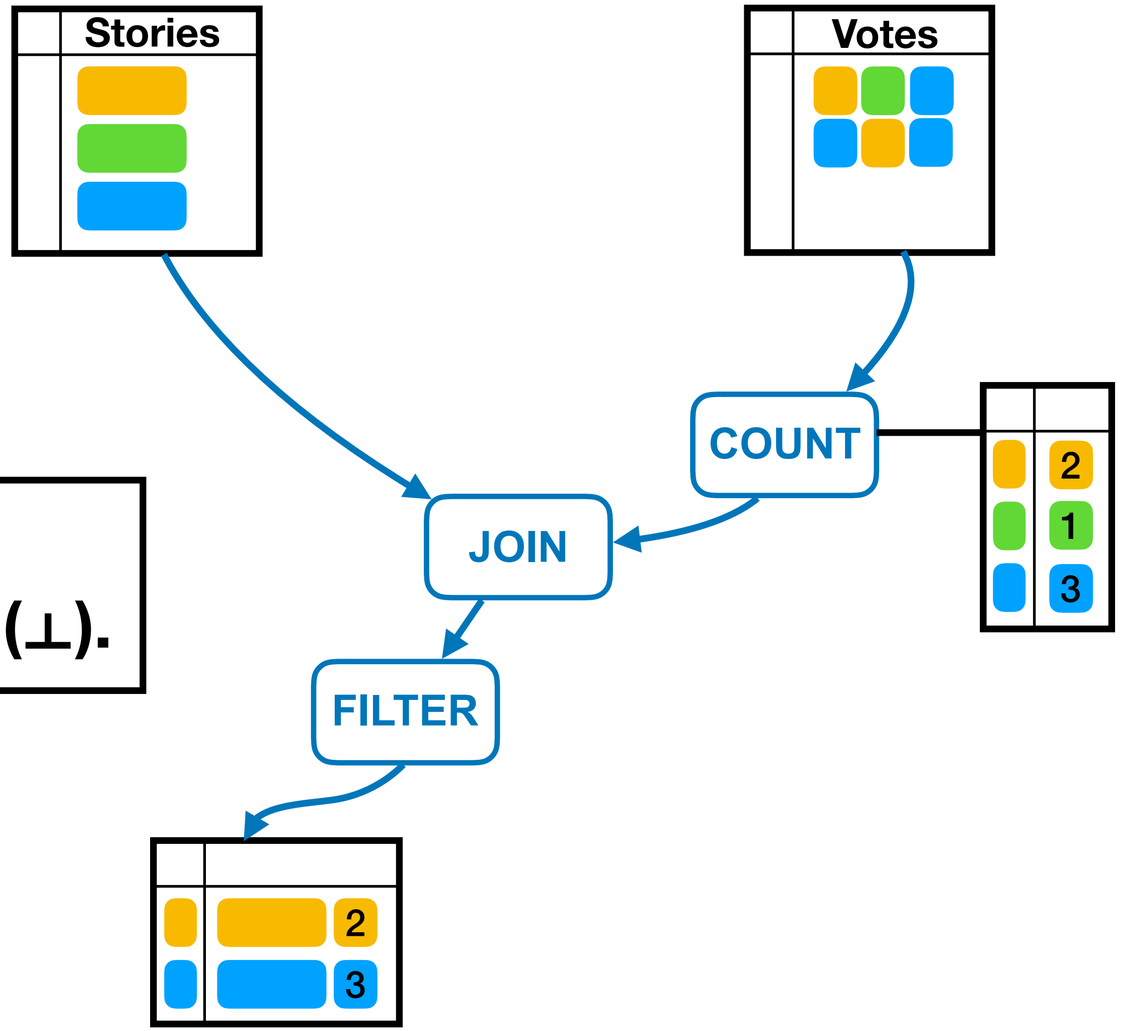


New model:

Partially-stateful data-flow

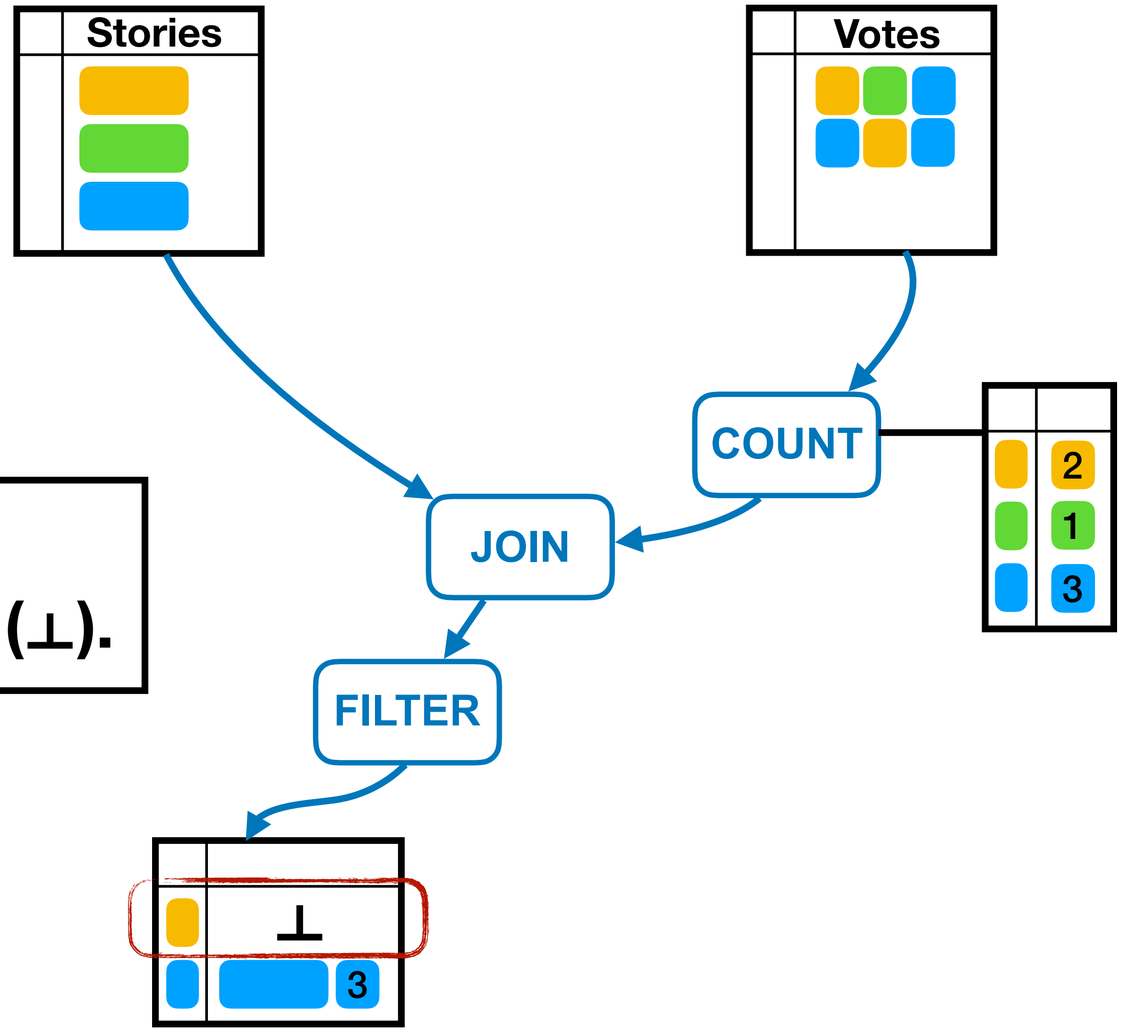
Partially-stateful data-flow

Data-flow state is *partial*:
entries for some keys are absent (⊥).



Frontend

Partially-stateful data-flow



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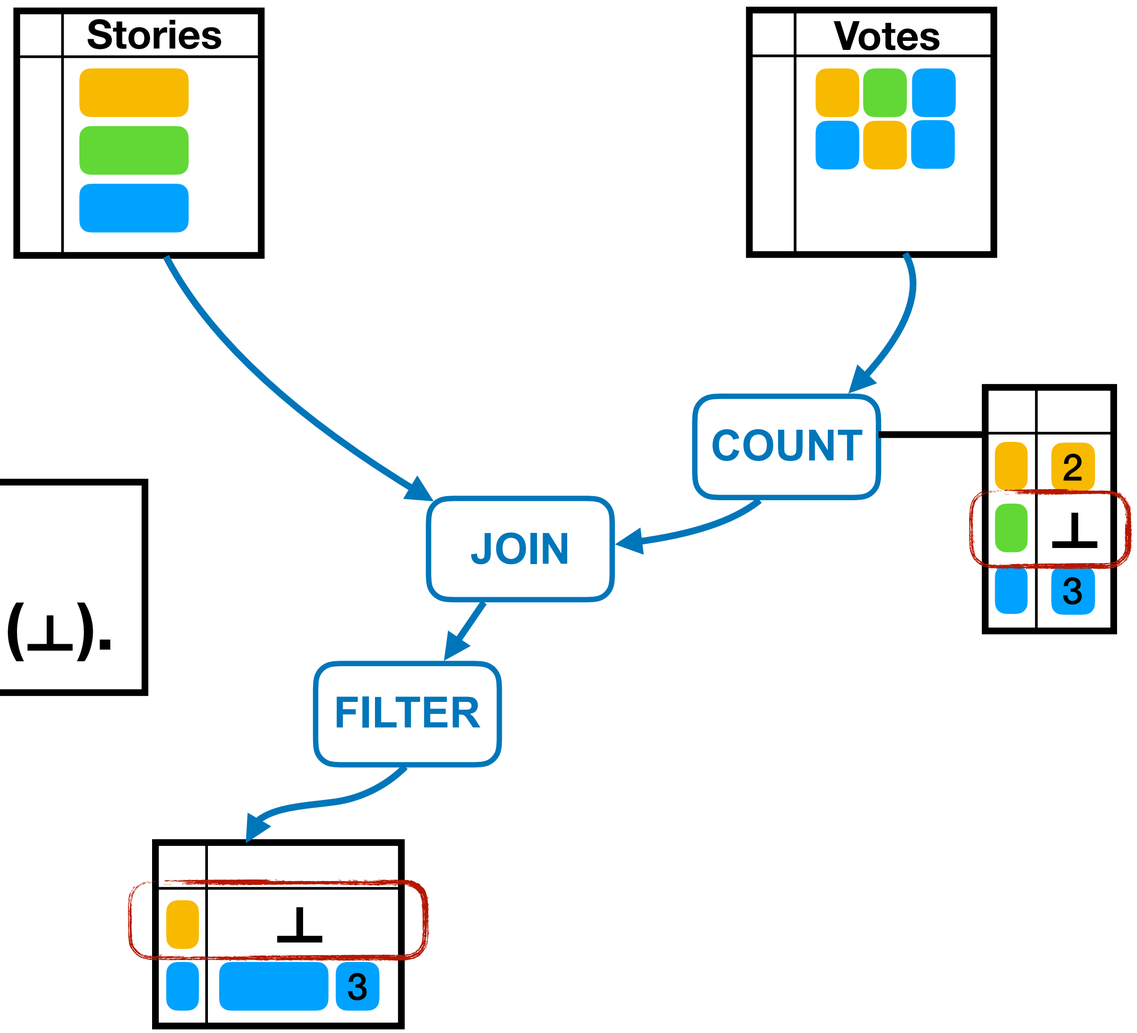

Frontend

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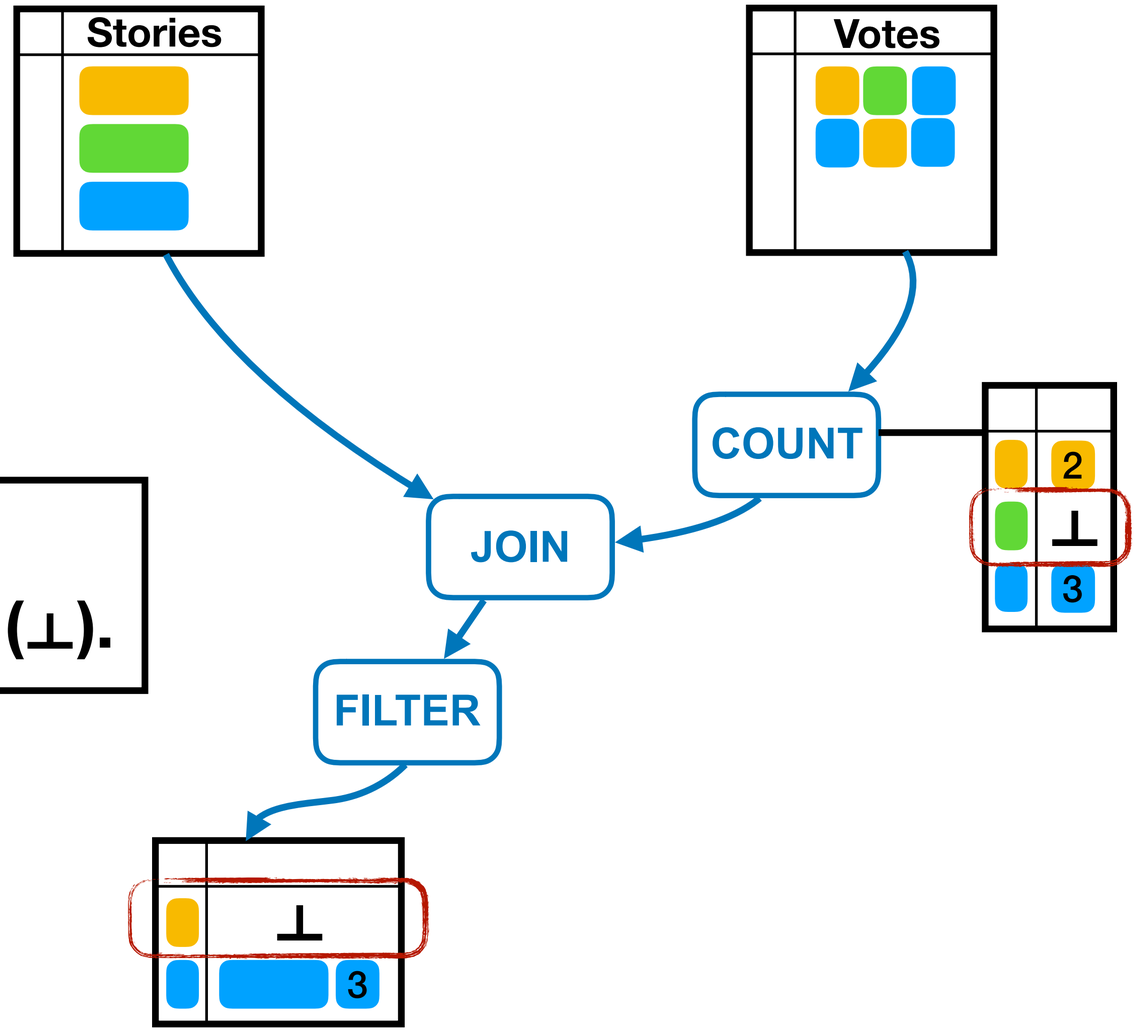

Data-flow state is *partial*:
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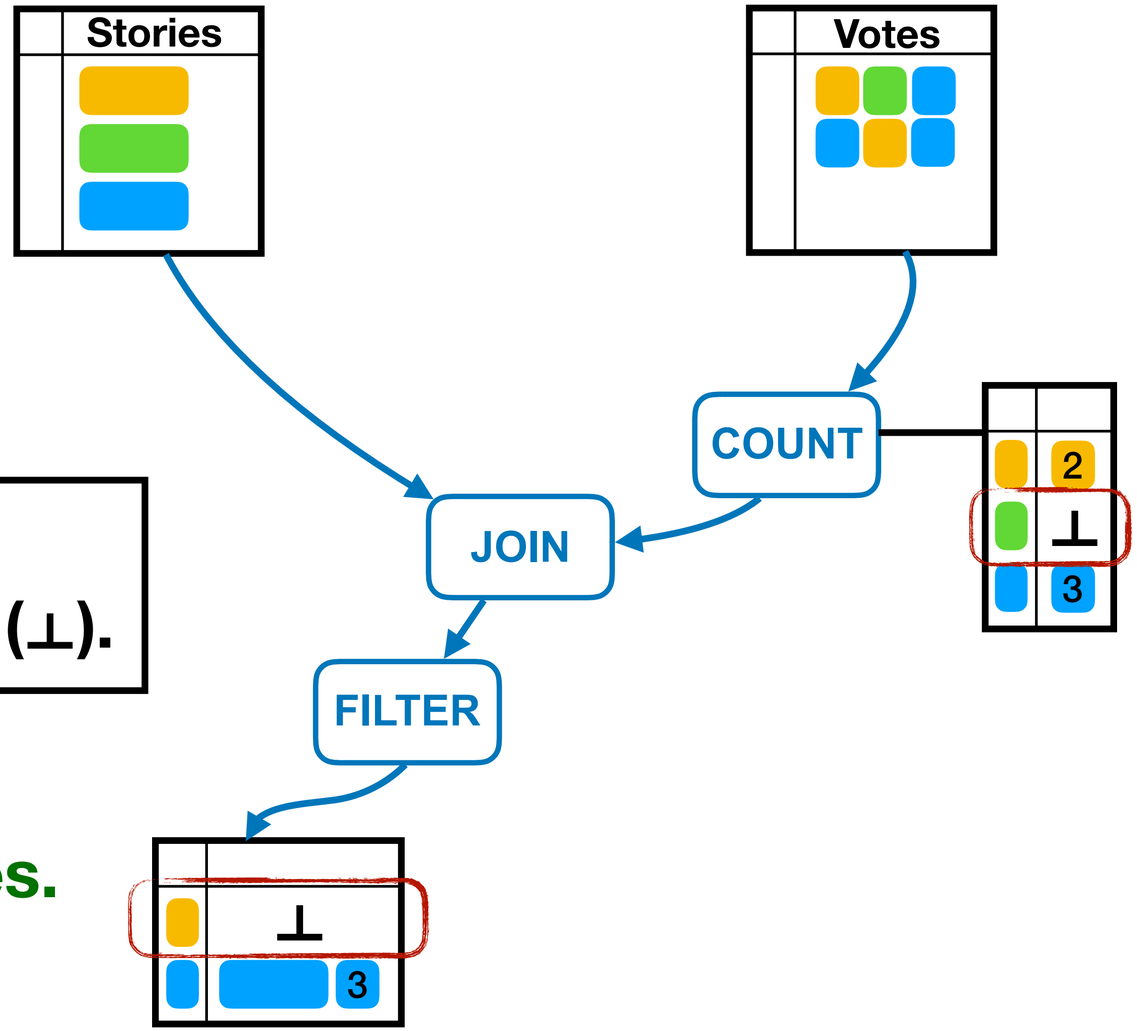
Lower memory footprint.

Frontend

Partially-stateful data-flow



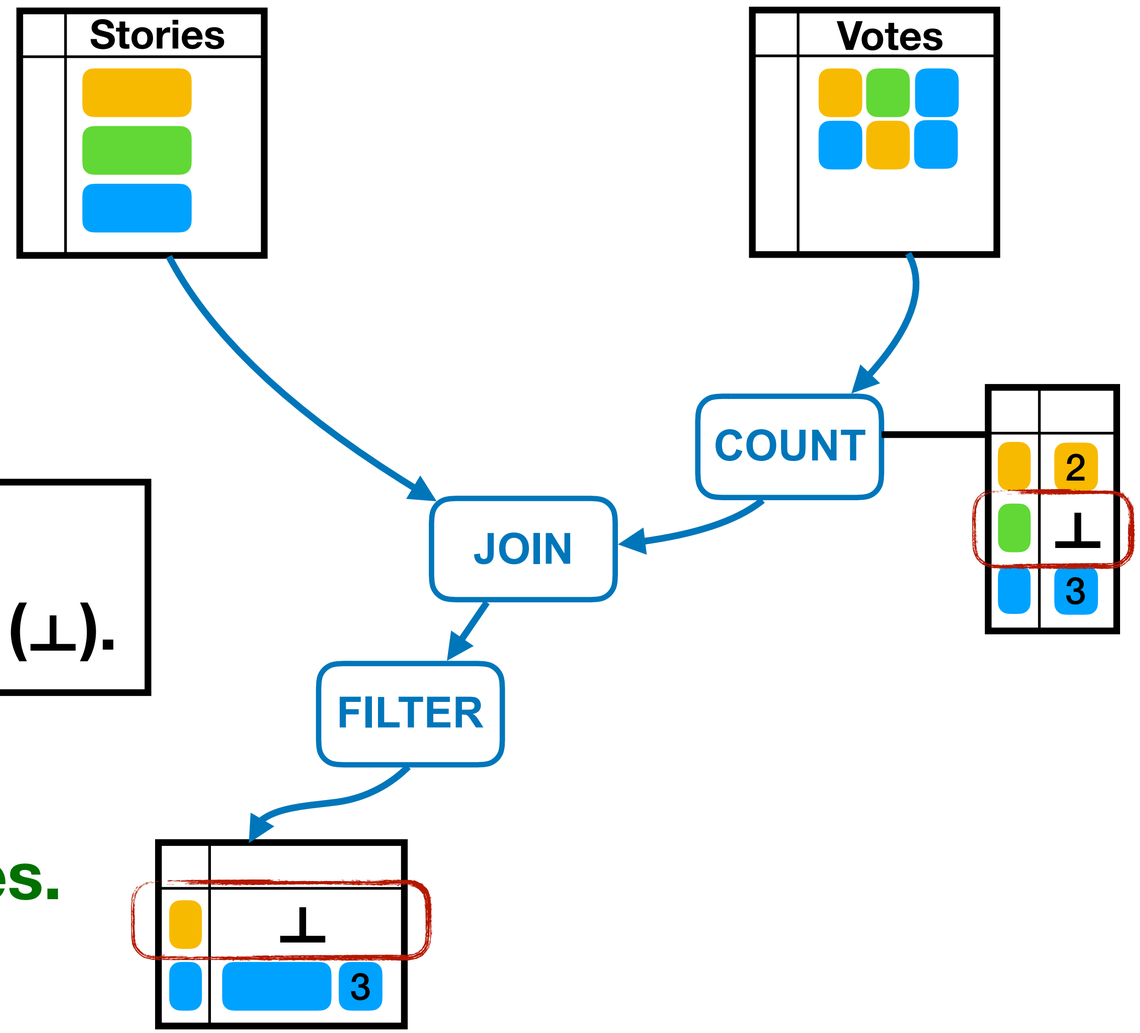
Data-flow state is *partial*:
entries for some keys are absent (⊥).

Lower memory footprint.

No need to update absent entries.

Frontend

Partially-stateful data-flow



Data-flow state is *partial*:
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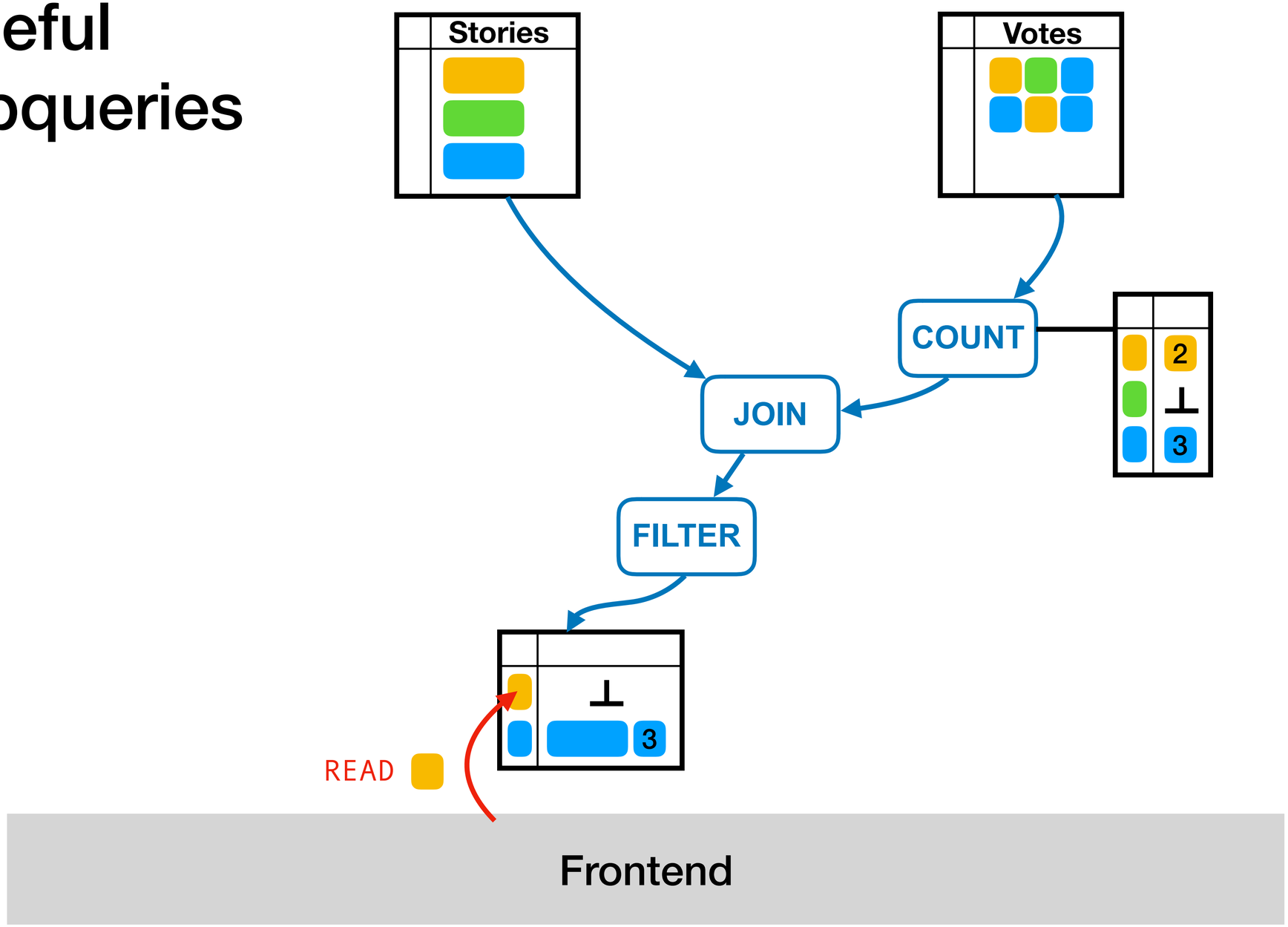
Lower memory footprint.

No need to update absent entries.

Enables live data-flow changes.

Frontend

Partially-stateful data-flow: upqueries



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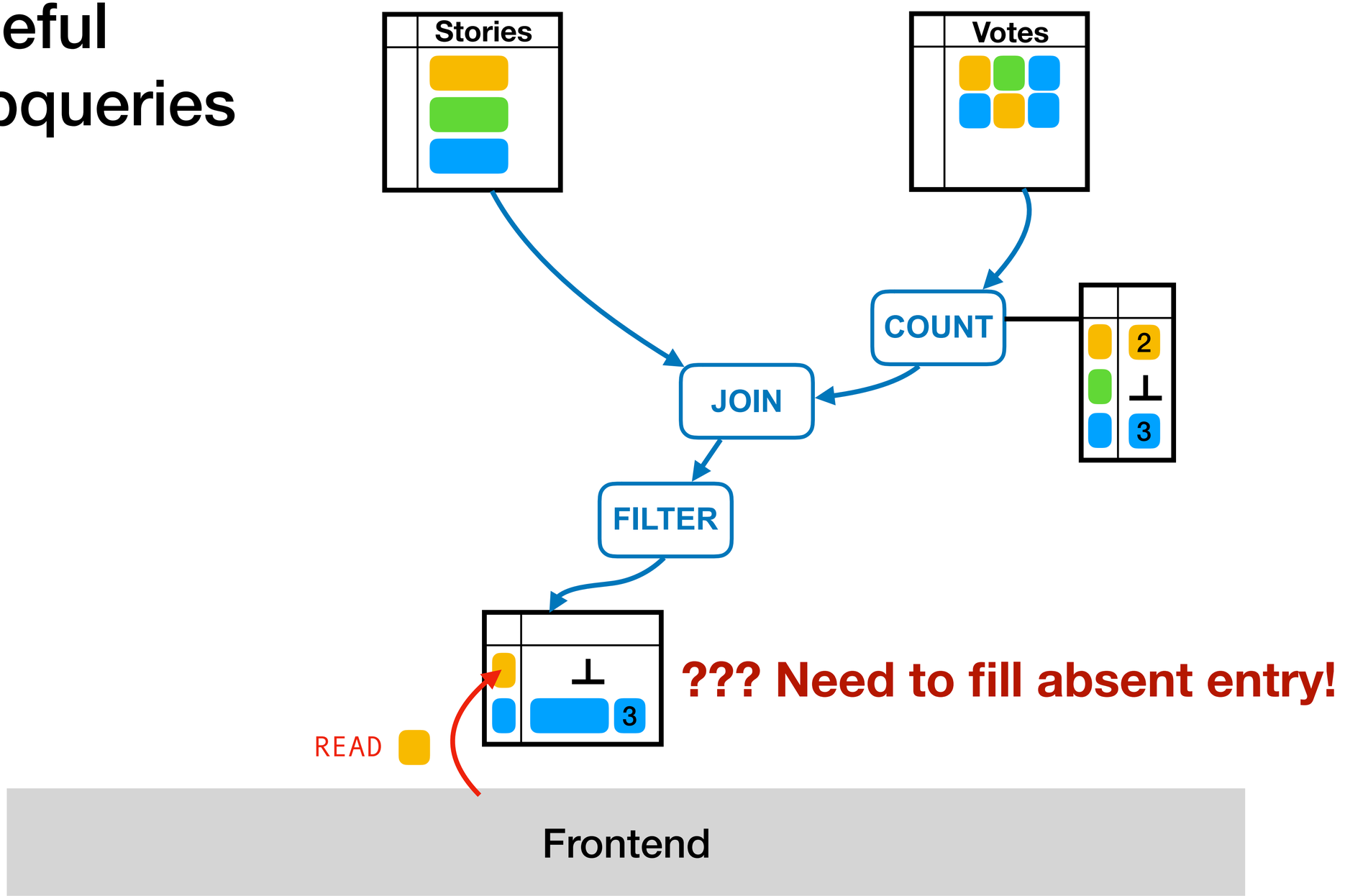

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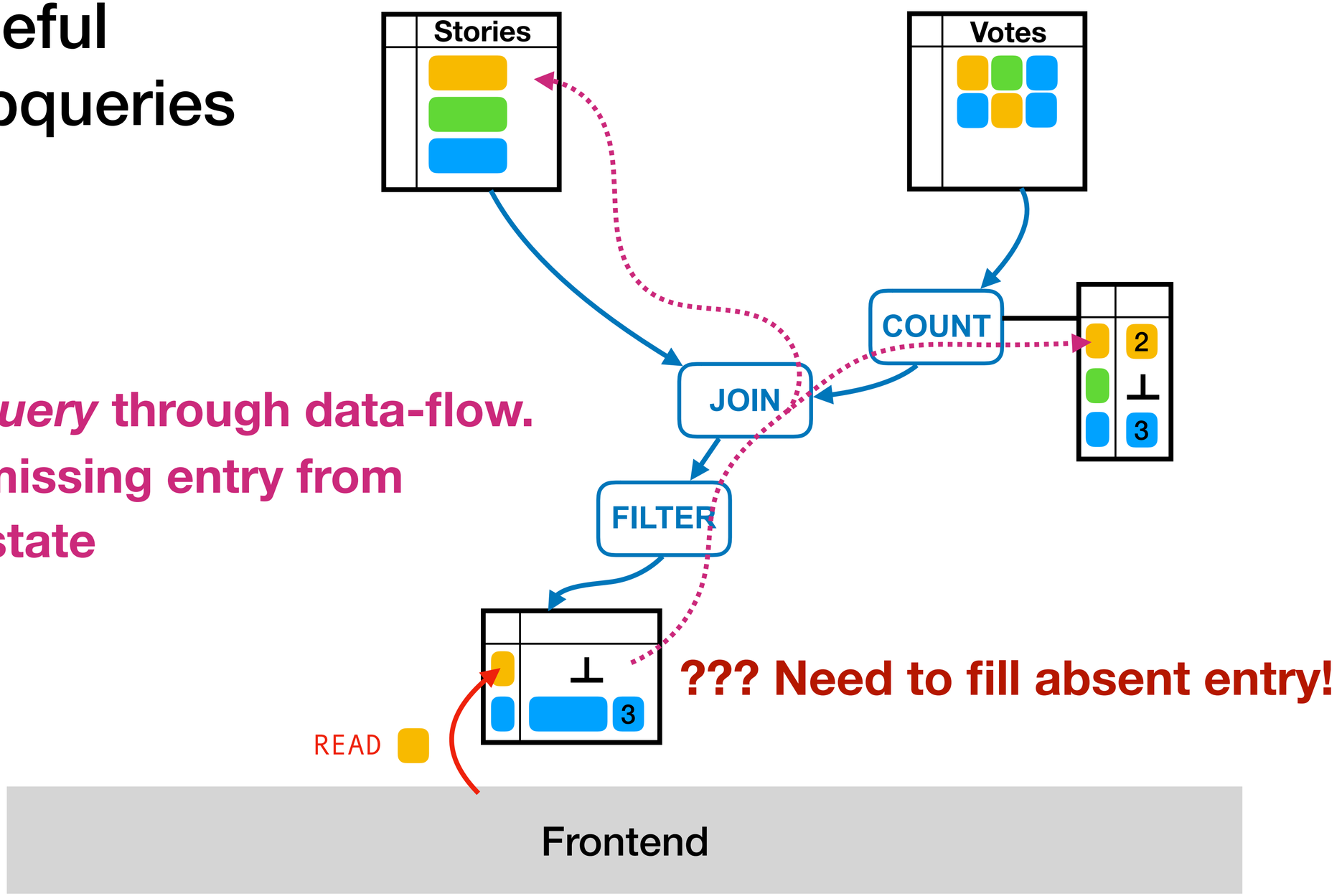
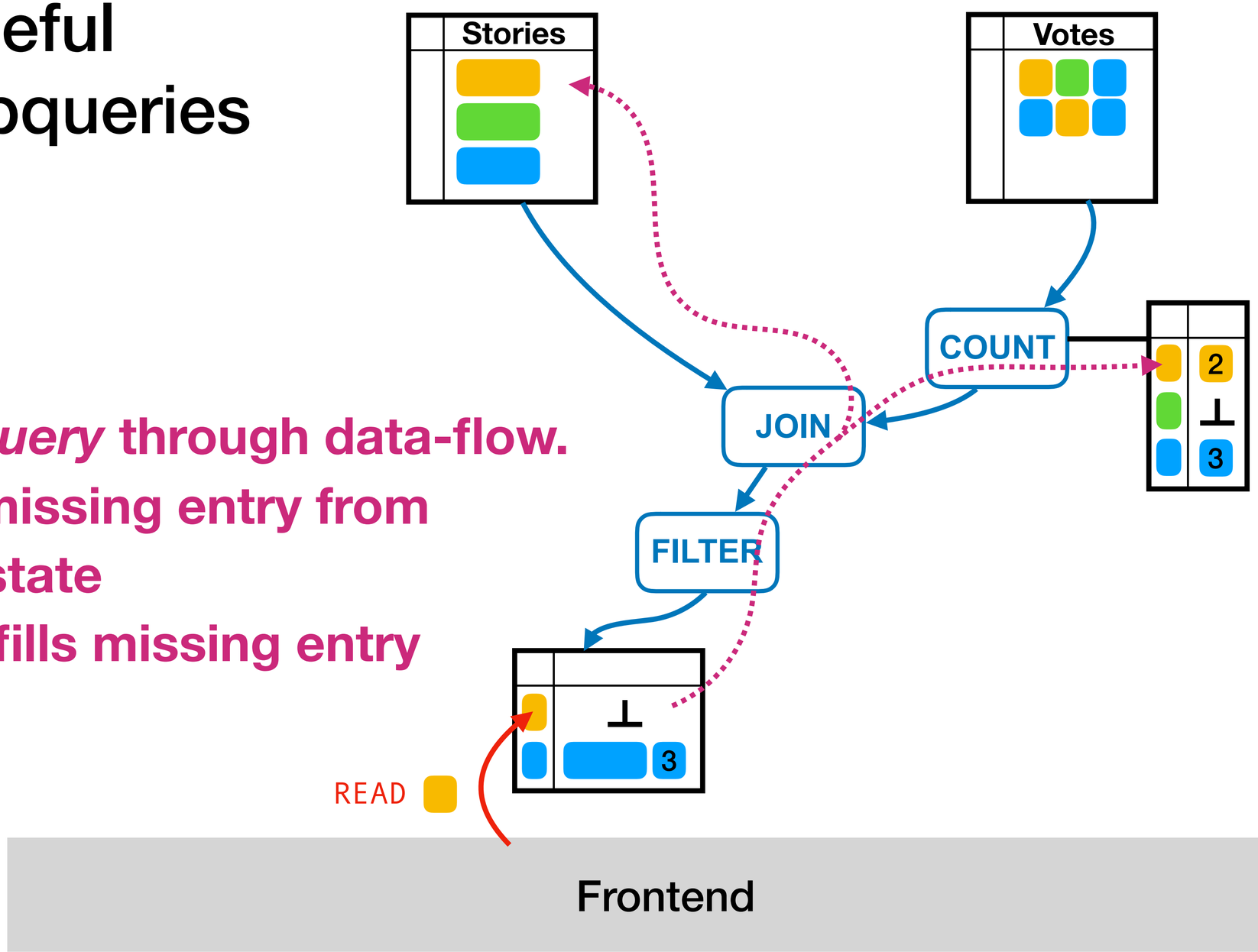


Solution: *upquery* through data-flow.

- **Compute missing entry from upstream state**

??? Need to fill absent entry!

Partially-stateful data-flow: upqueries



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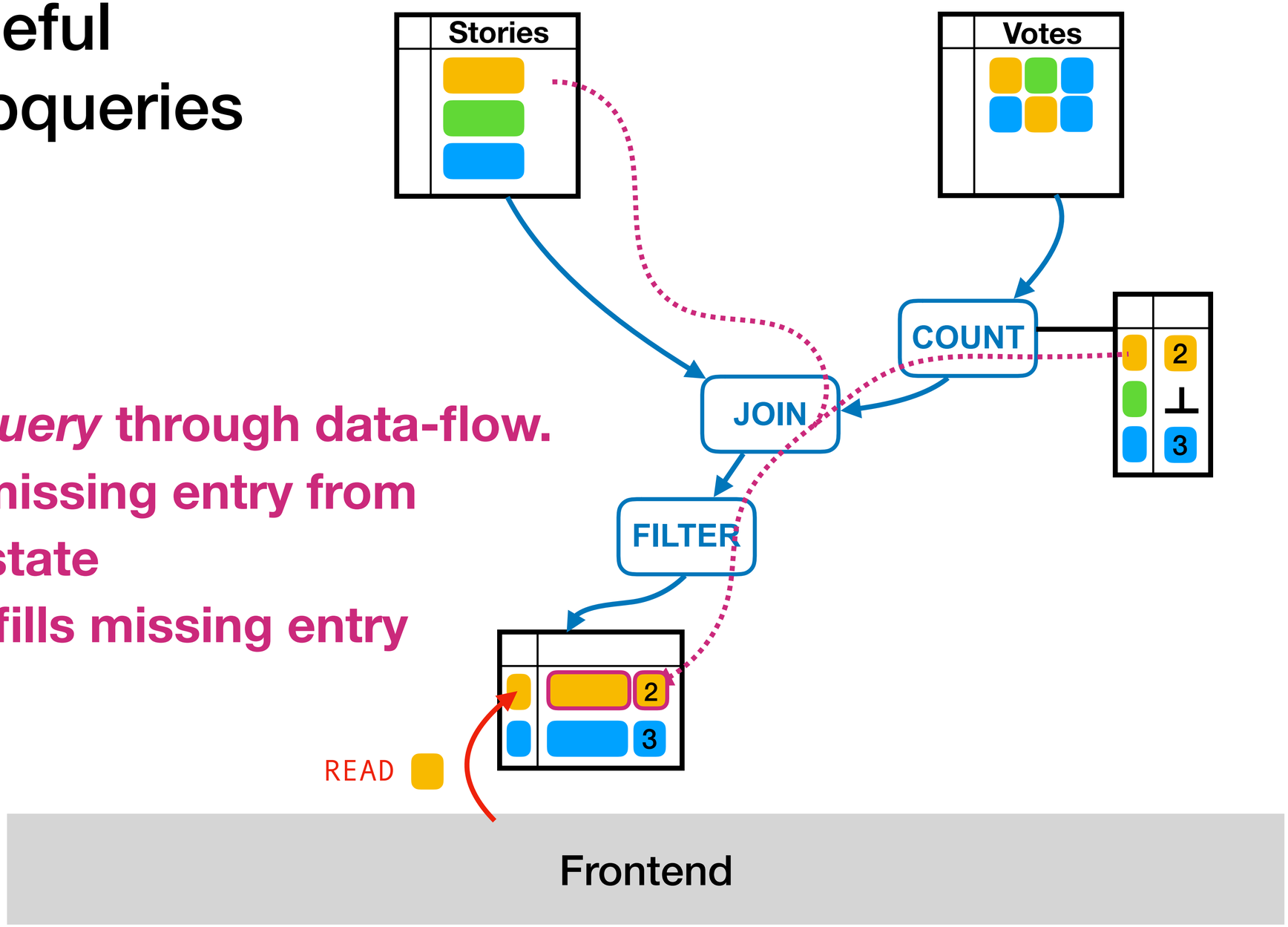

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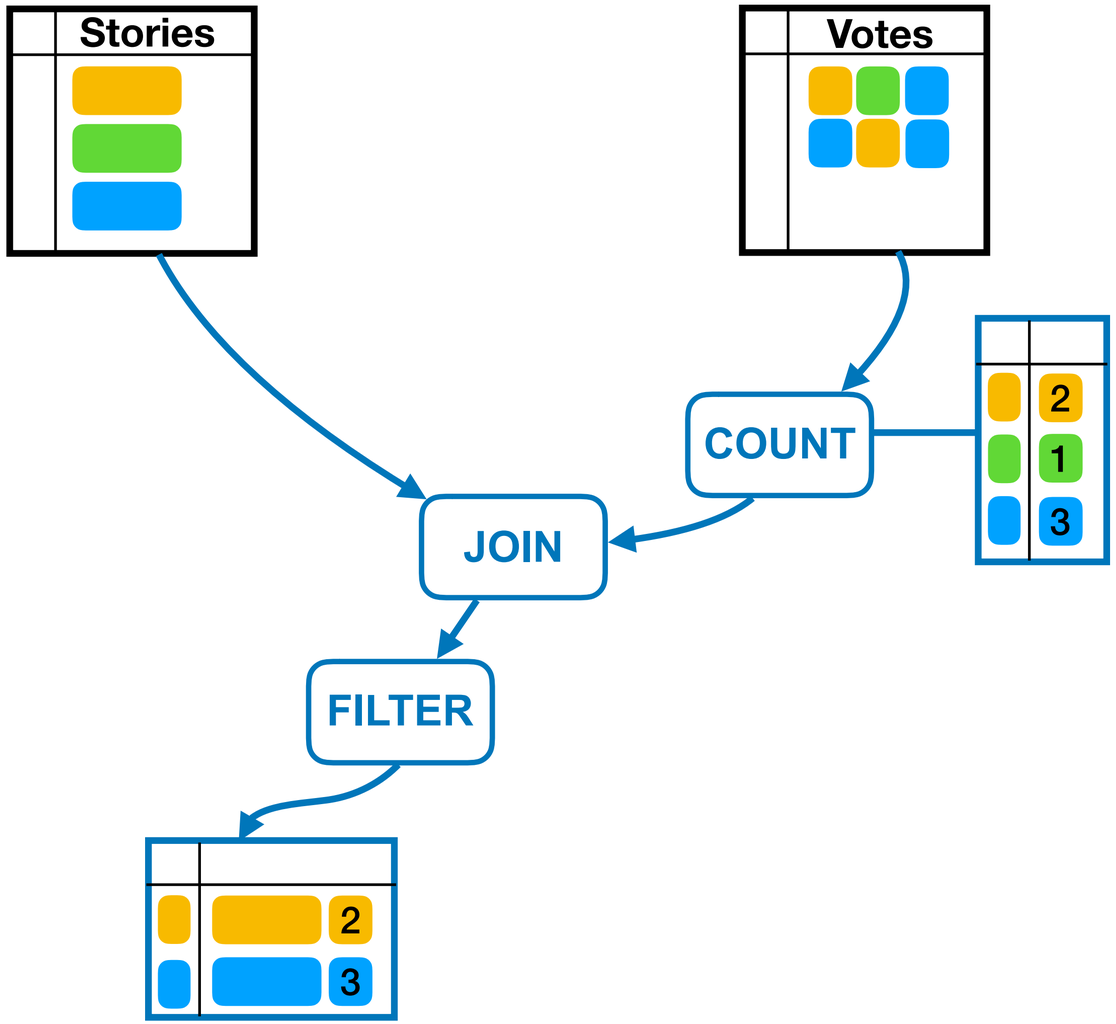
Partial state enables live data-flow changes

Start new views and operator state **empty**, **fill via upqueries**.



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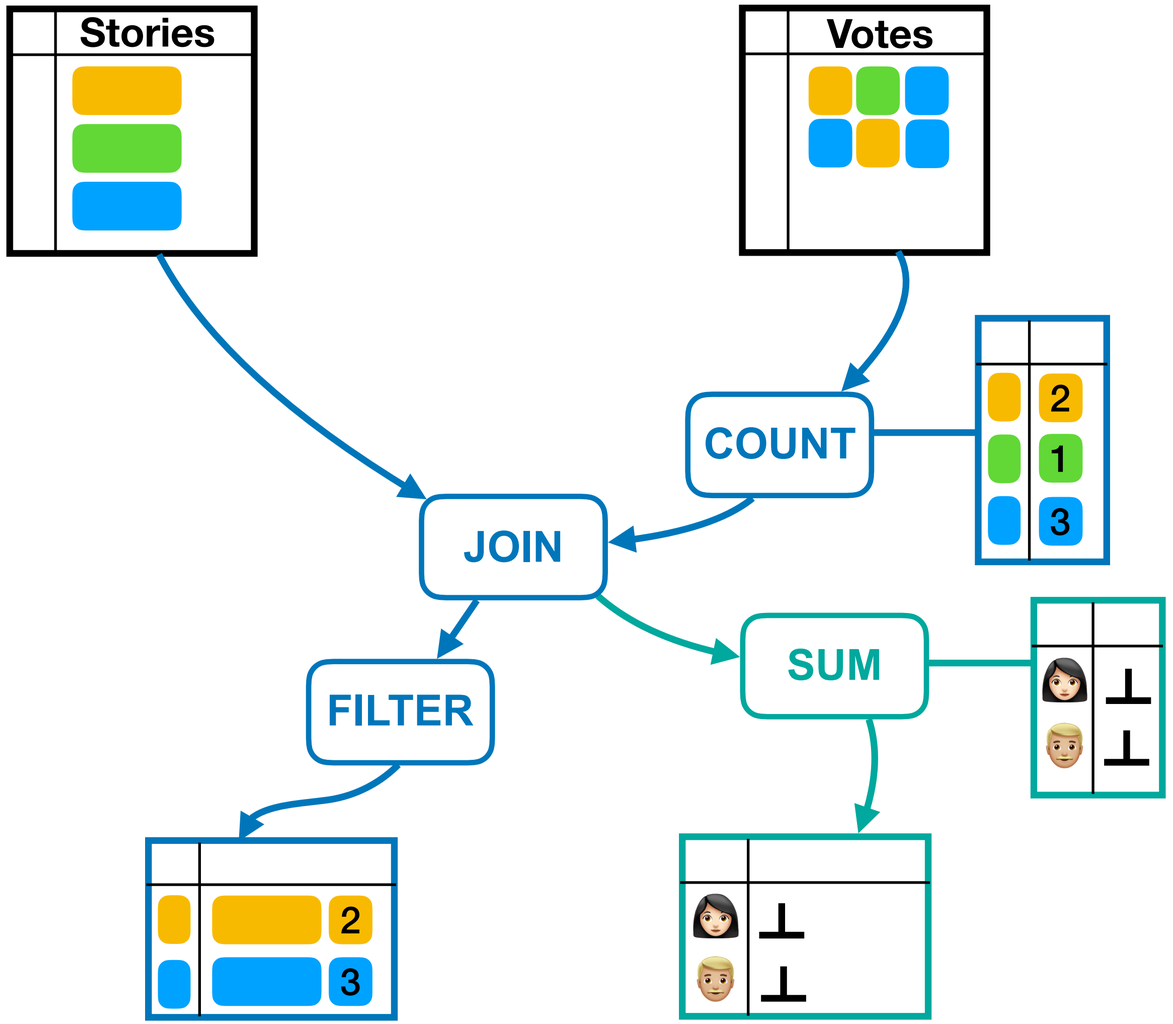

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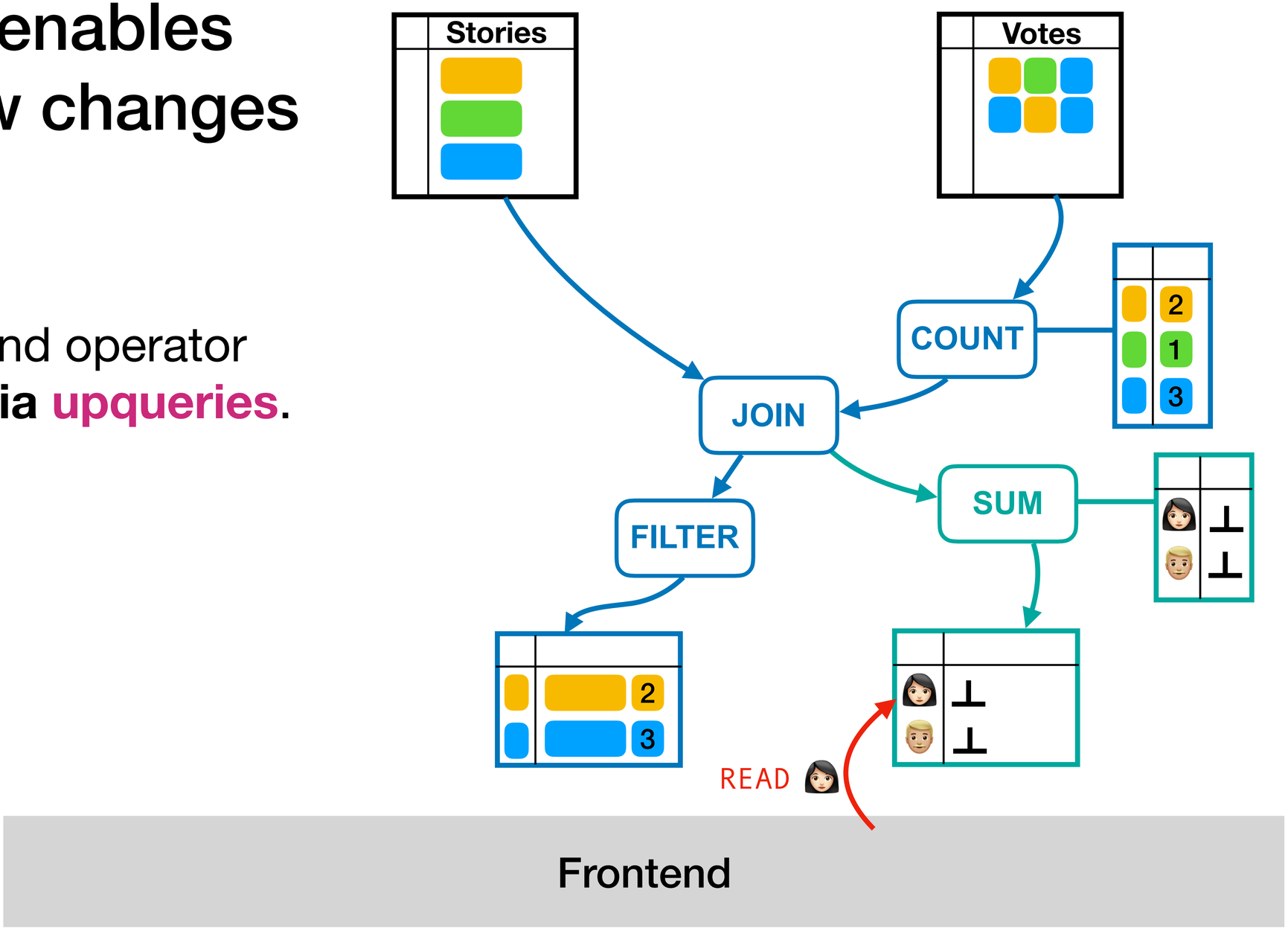

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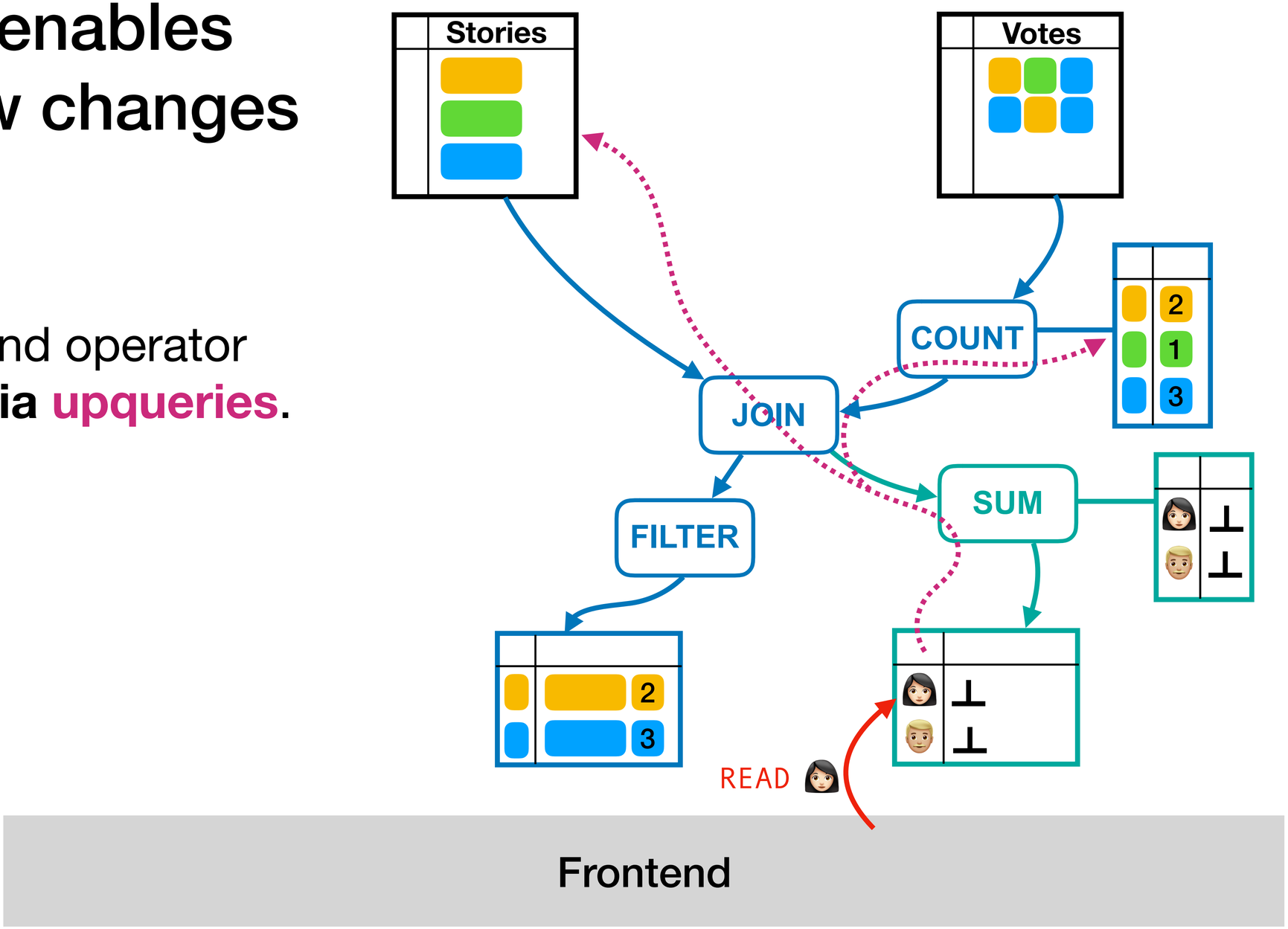

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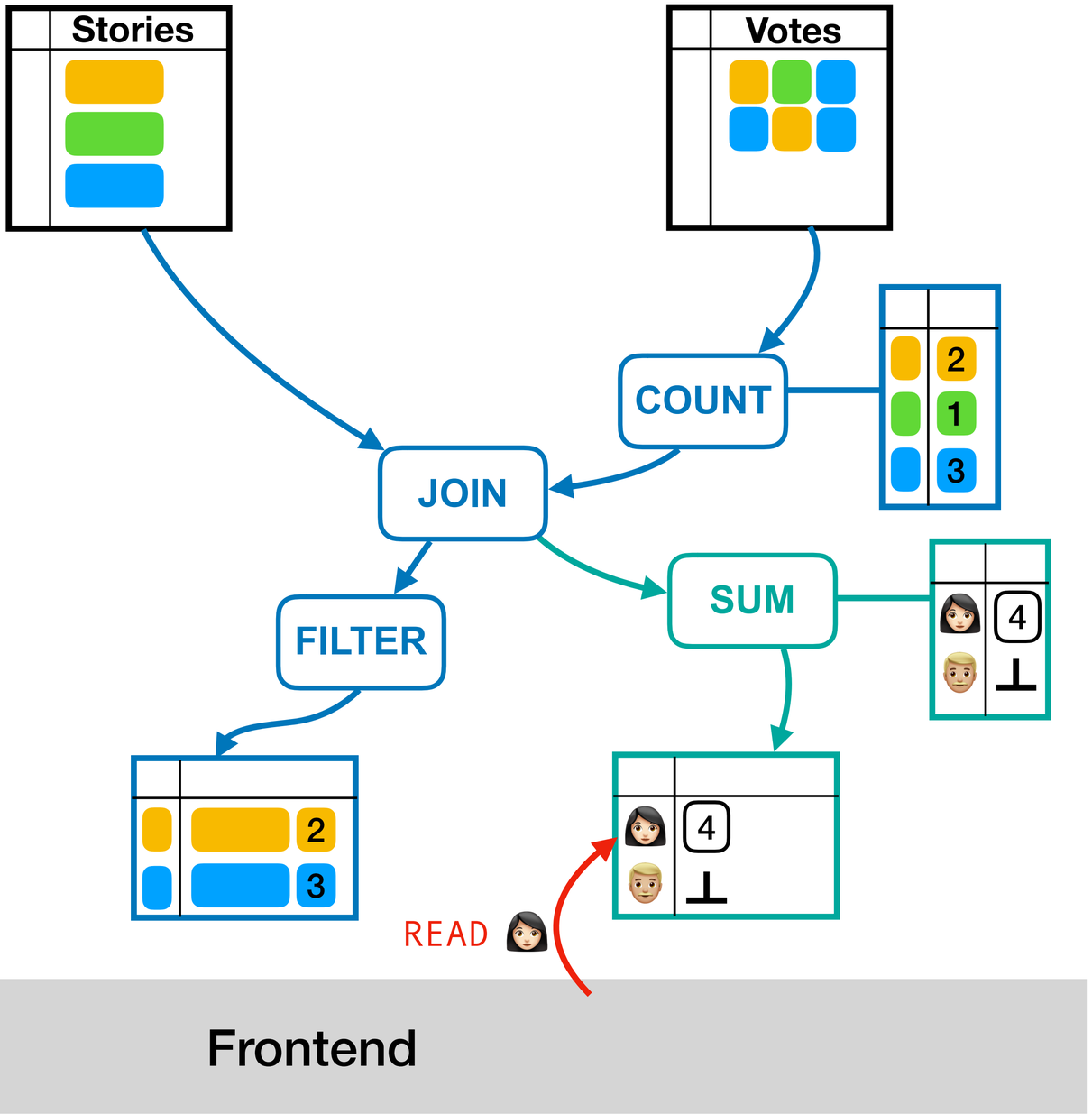
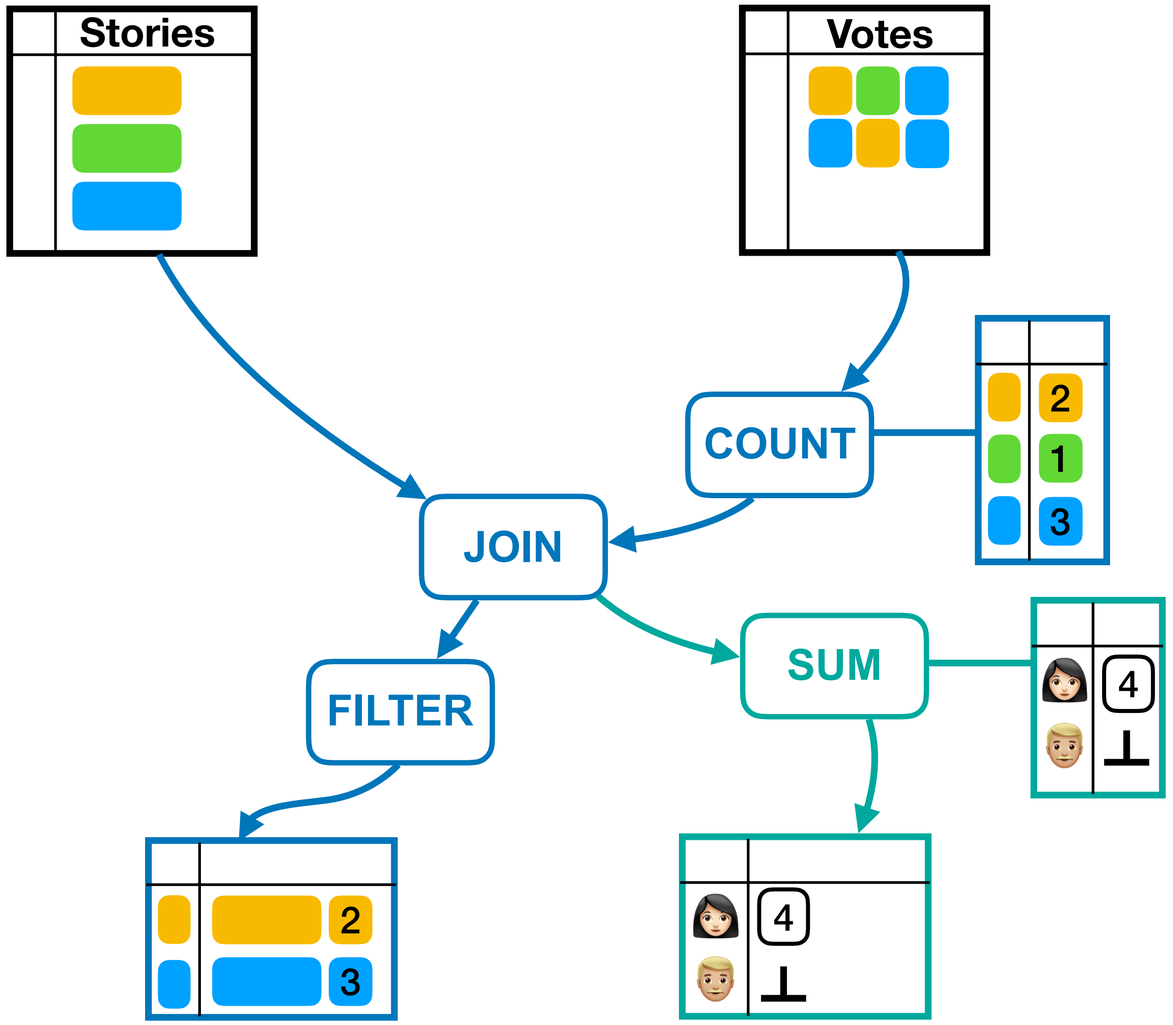


High performance requires concurrency

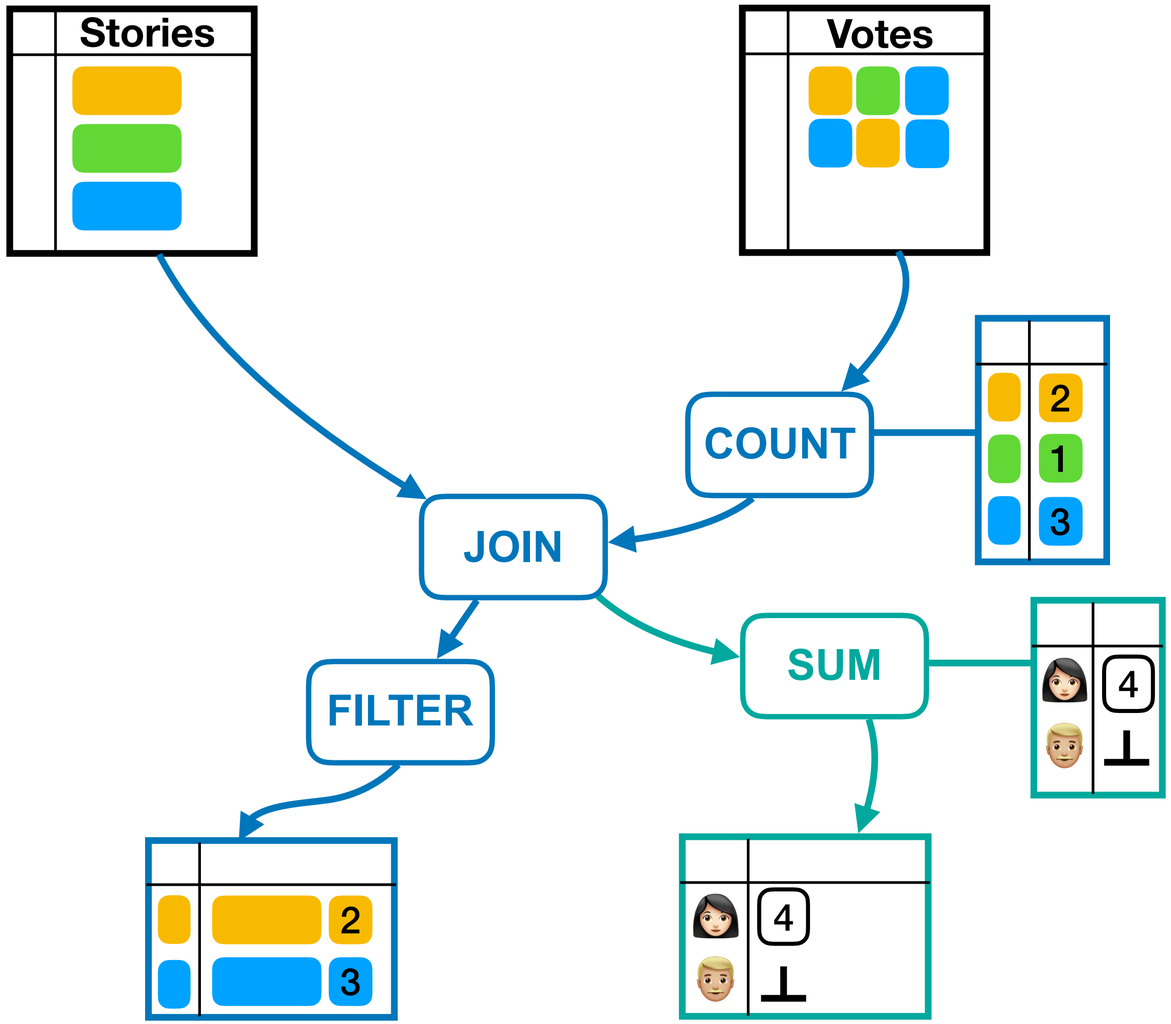


Frontend

High performance requires concurrency

Process operators concurrently.
Read from views concurrently.
Process shards concurrently.

Without global coordination!



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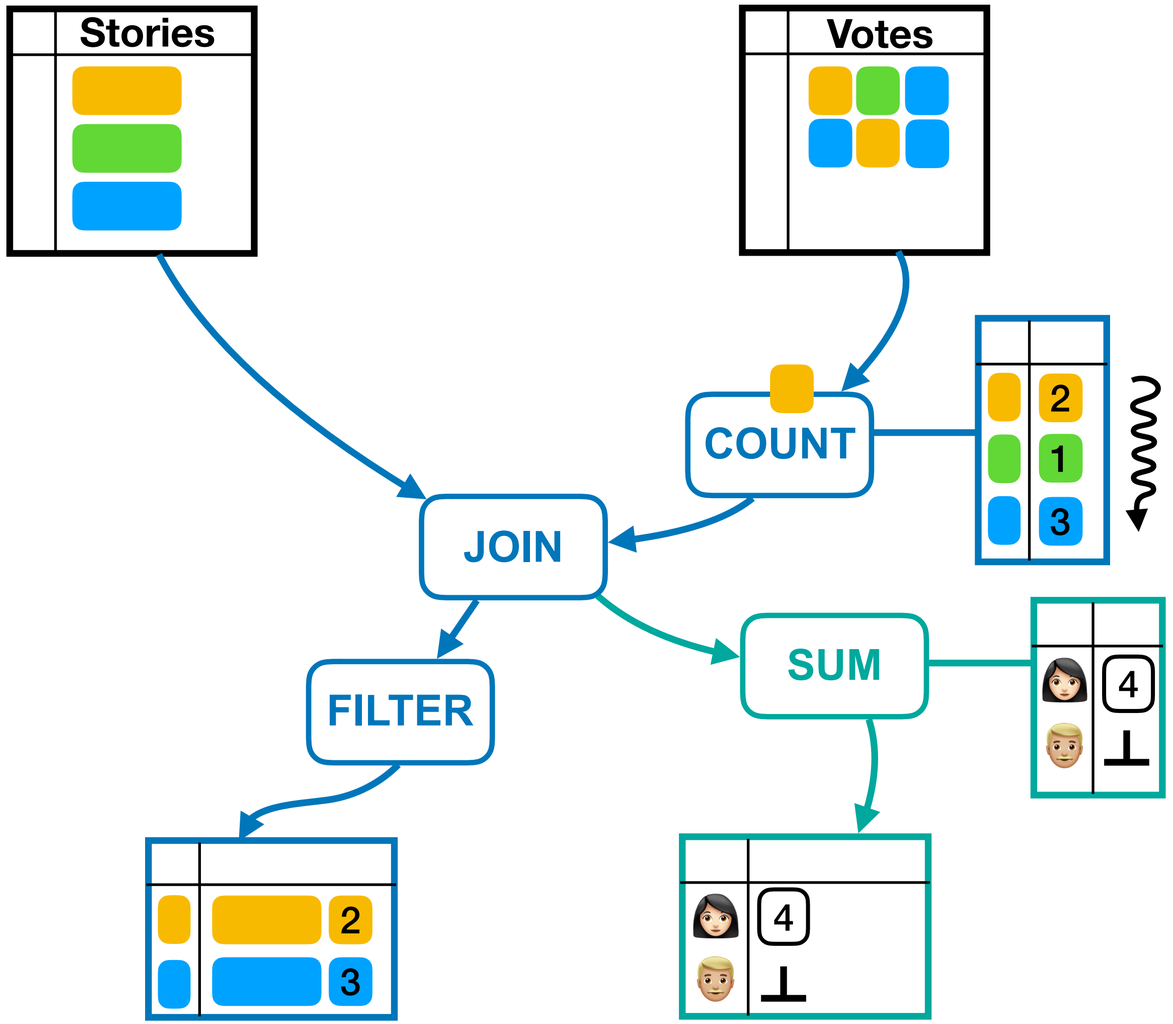

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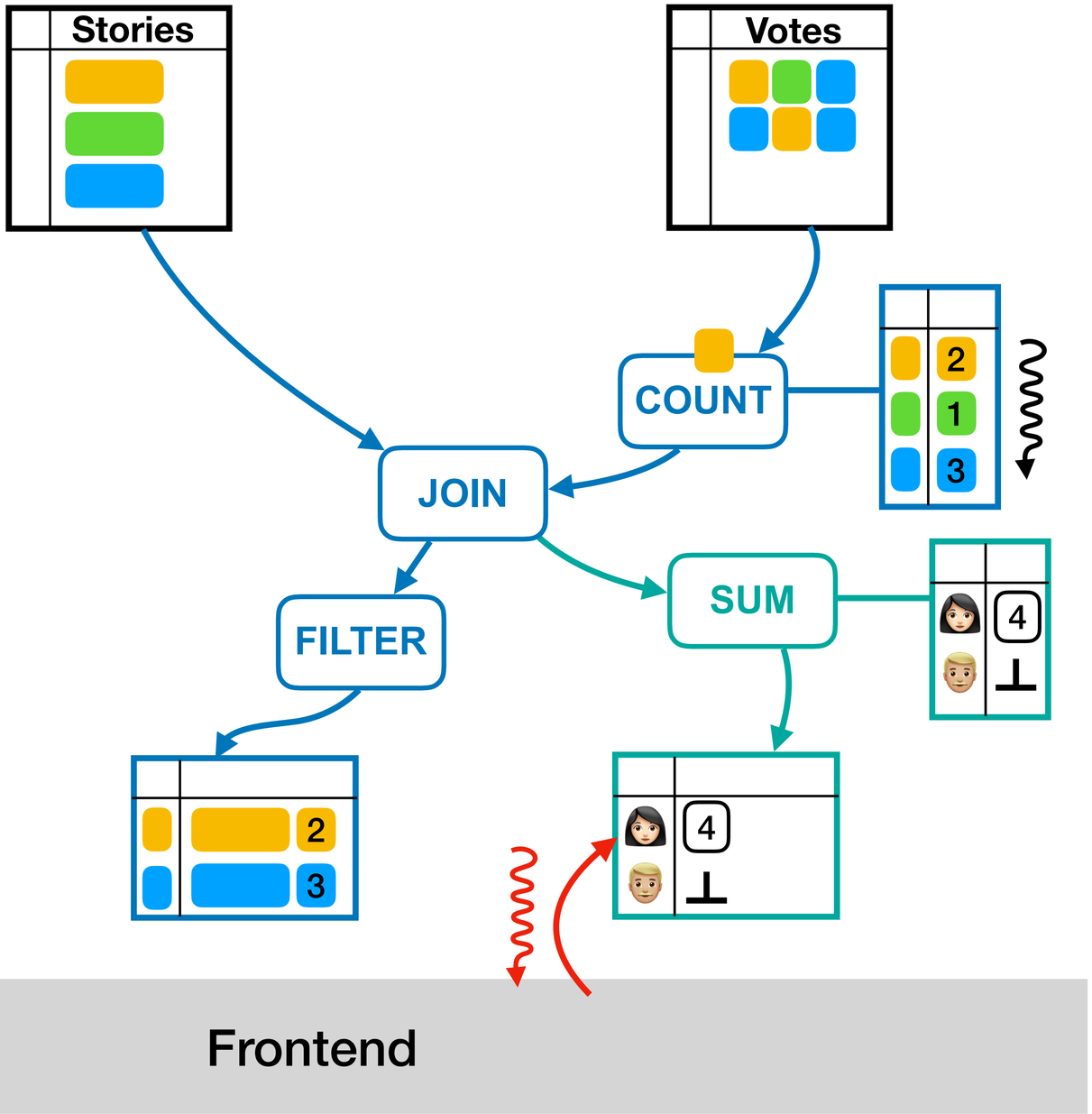

Challenges implementing partially-stateful data-flow

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1. Concurrent upqueries and forward processing — races!

Must maintain **correctness** under concurrency!

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Goal: upquery restores state as if present all along.

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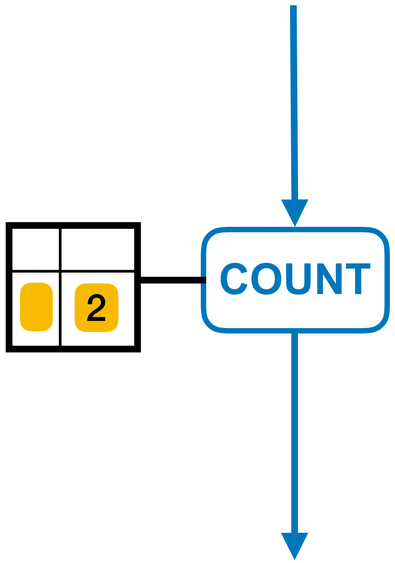

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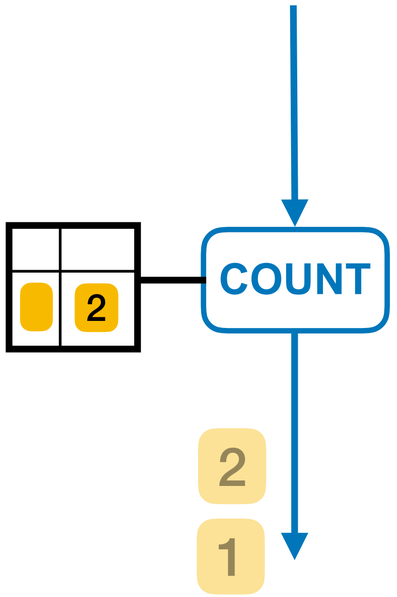
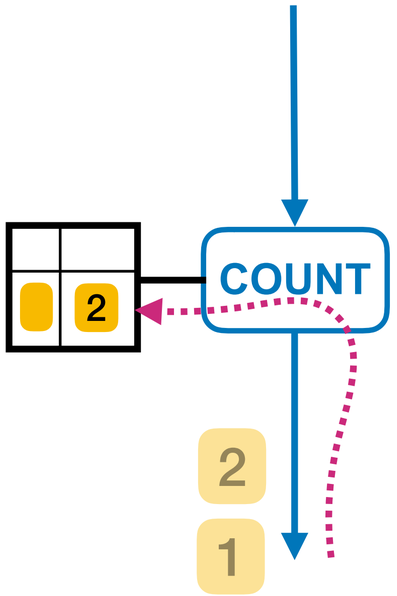


Correctness under concurrency

Goal: upquery restores state as if present all along.



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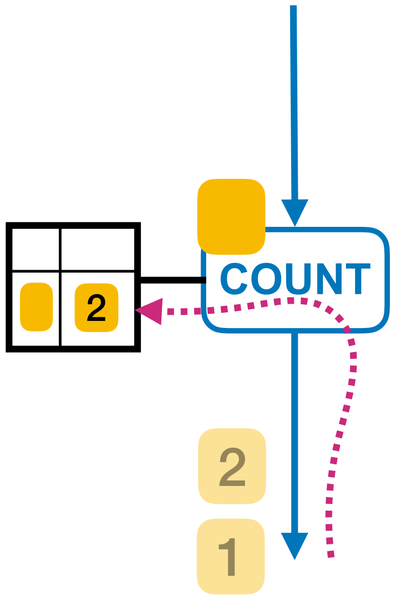

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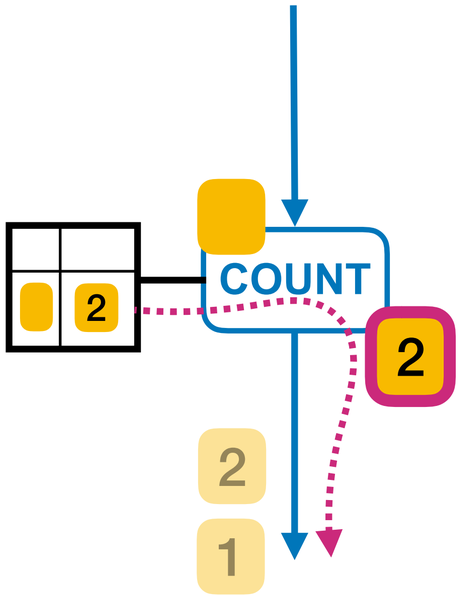

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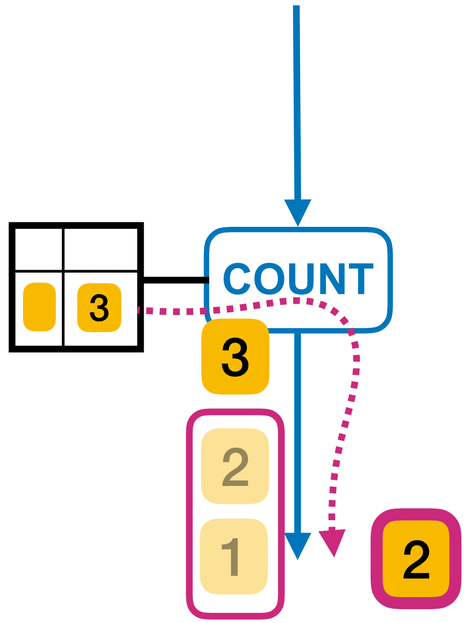

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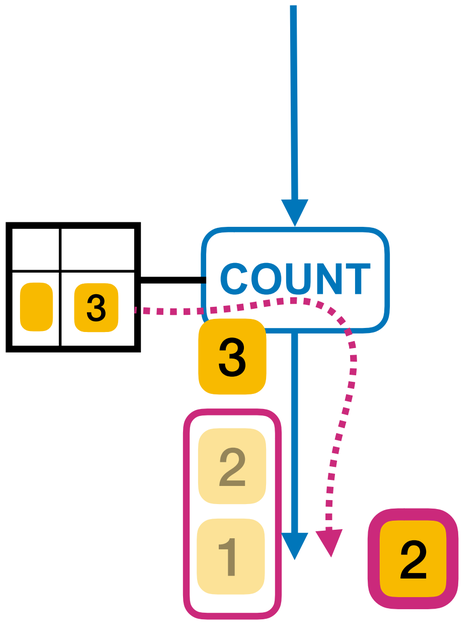

Upquery response is a **snapshot** of state

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Solution: Maintain **order** of upquery response and surrounding updates, despite lack of global coordination.

Upquery responses in total order with updates

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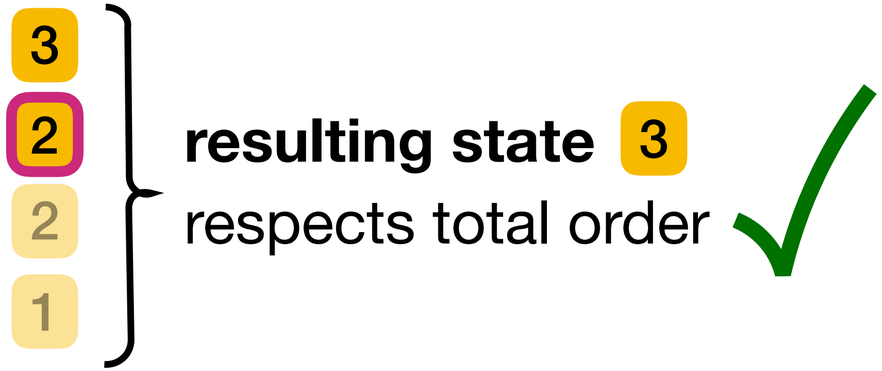

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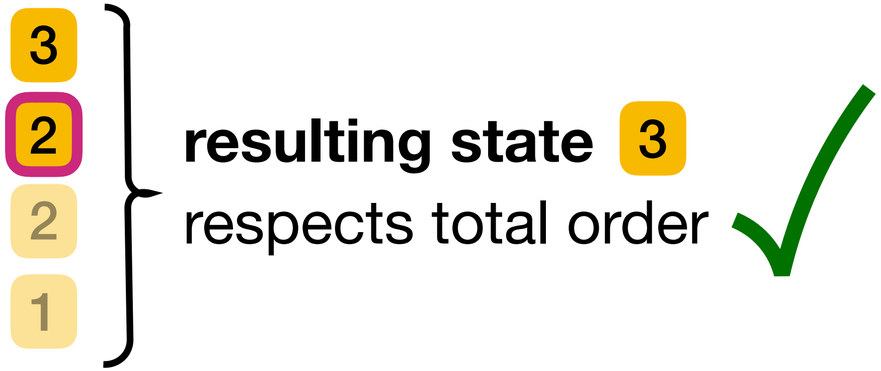

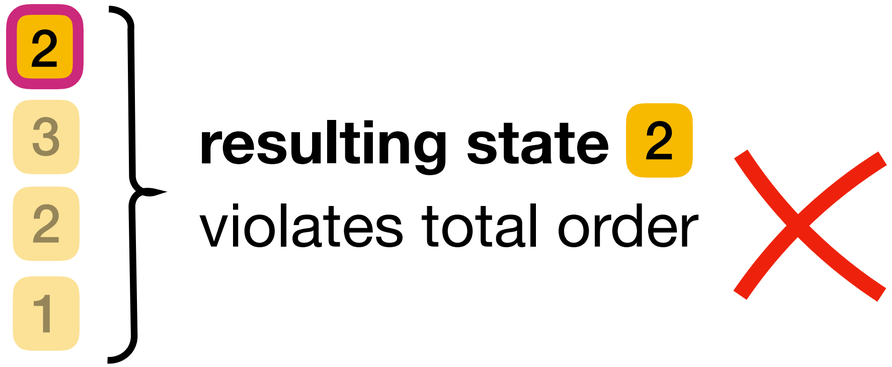

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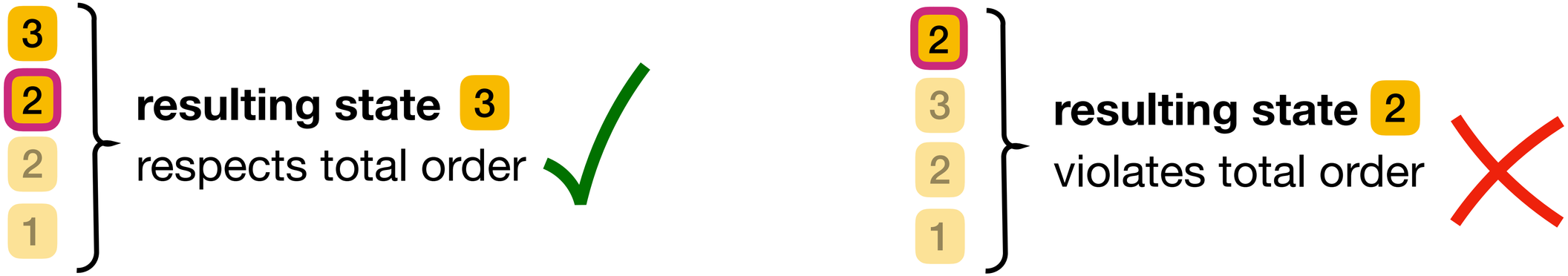

More complex cases: merged upquery responses, evictions (**Paper**).

Challenges implementing partially-stateful data-flow

1. Concurrent upqueries and forward processing — races!

Must maintain **correctness** under concurrency!

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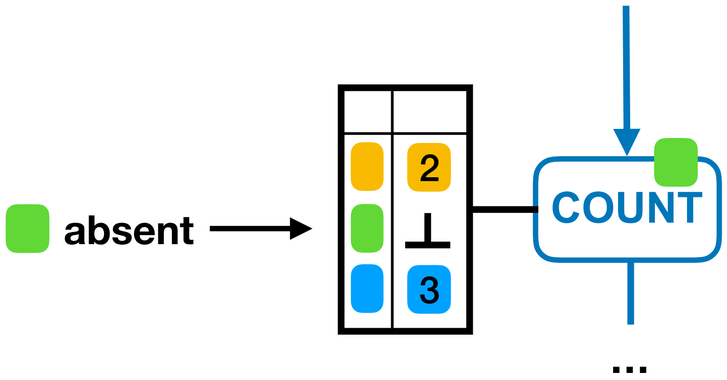

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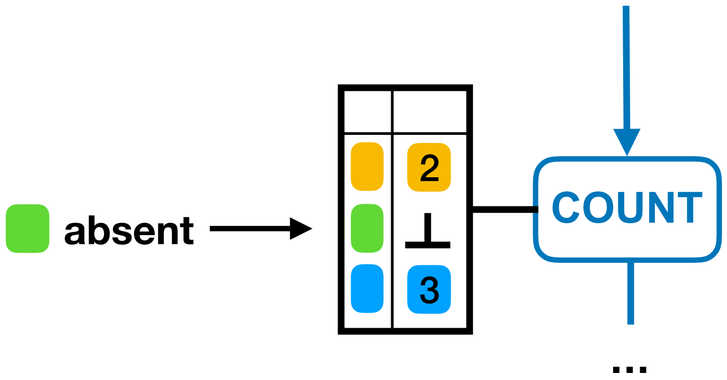

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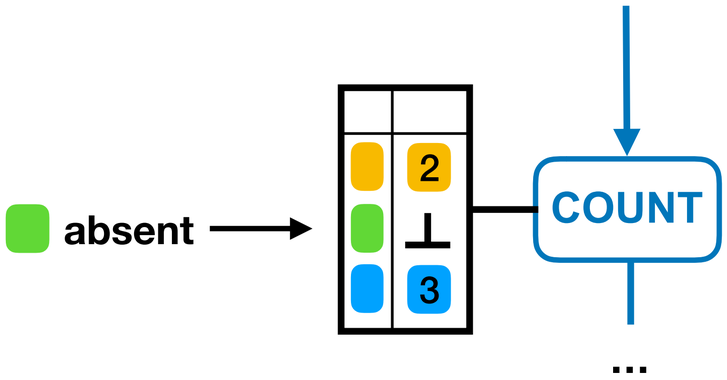

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Noria implementation

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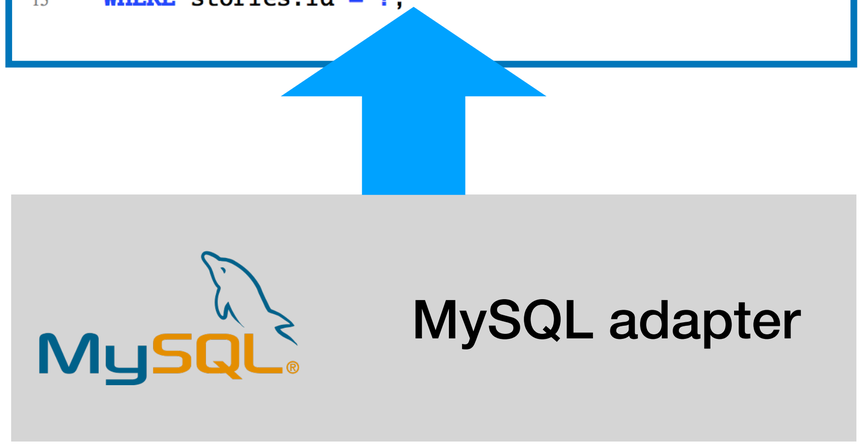

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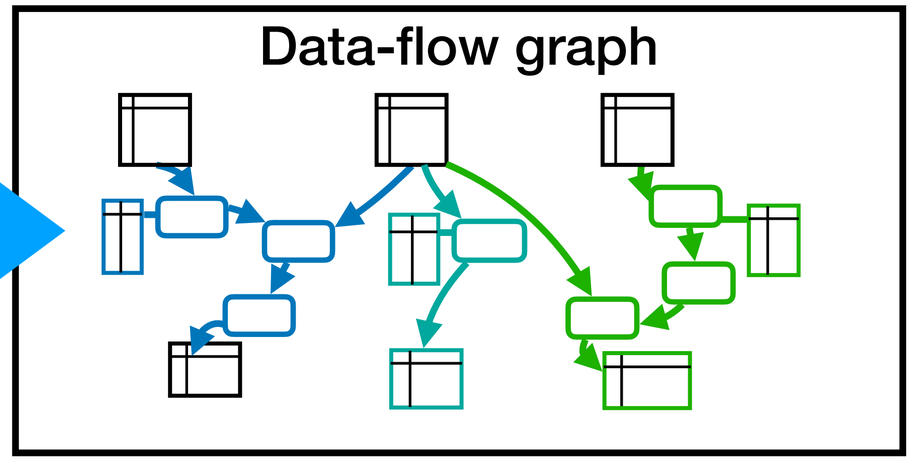

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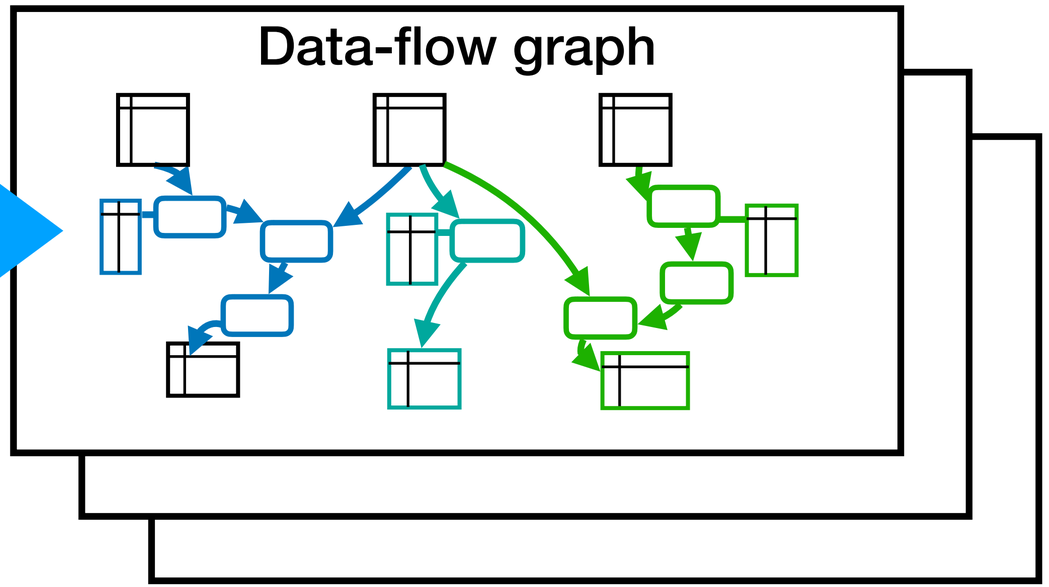

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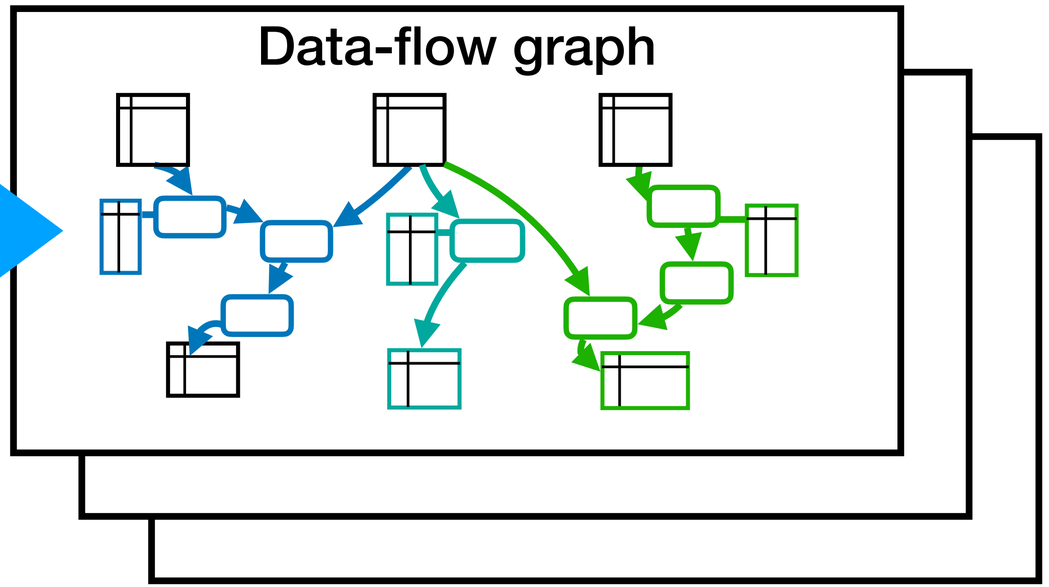

MySQL adapter

- 45k lines of Rust + 15k libraries
- RocksDB for base table storage
- ZooKeeper for leader election

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multi-machine experiments
comparison with differential dataflow } **see Paper**

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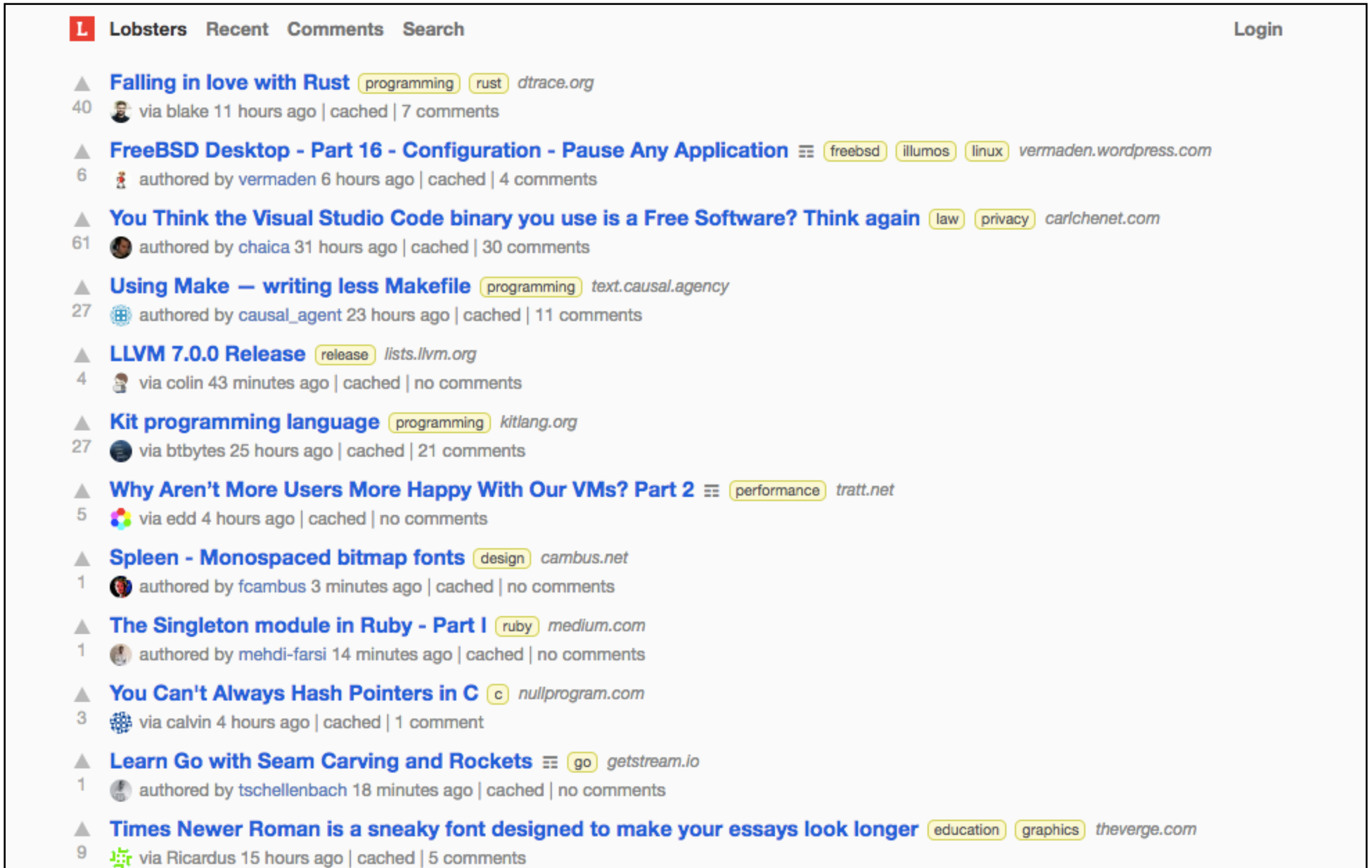

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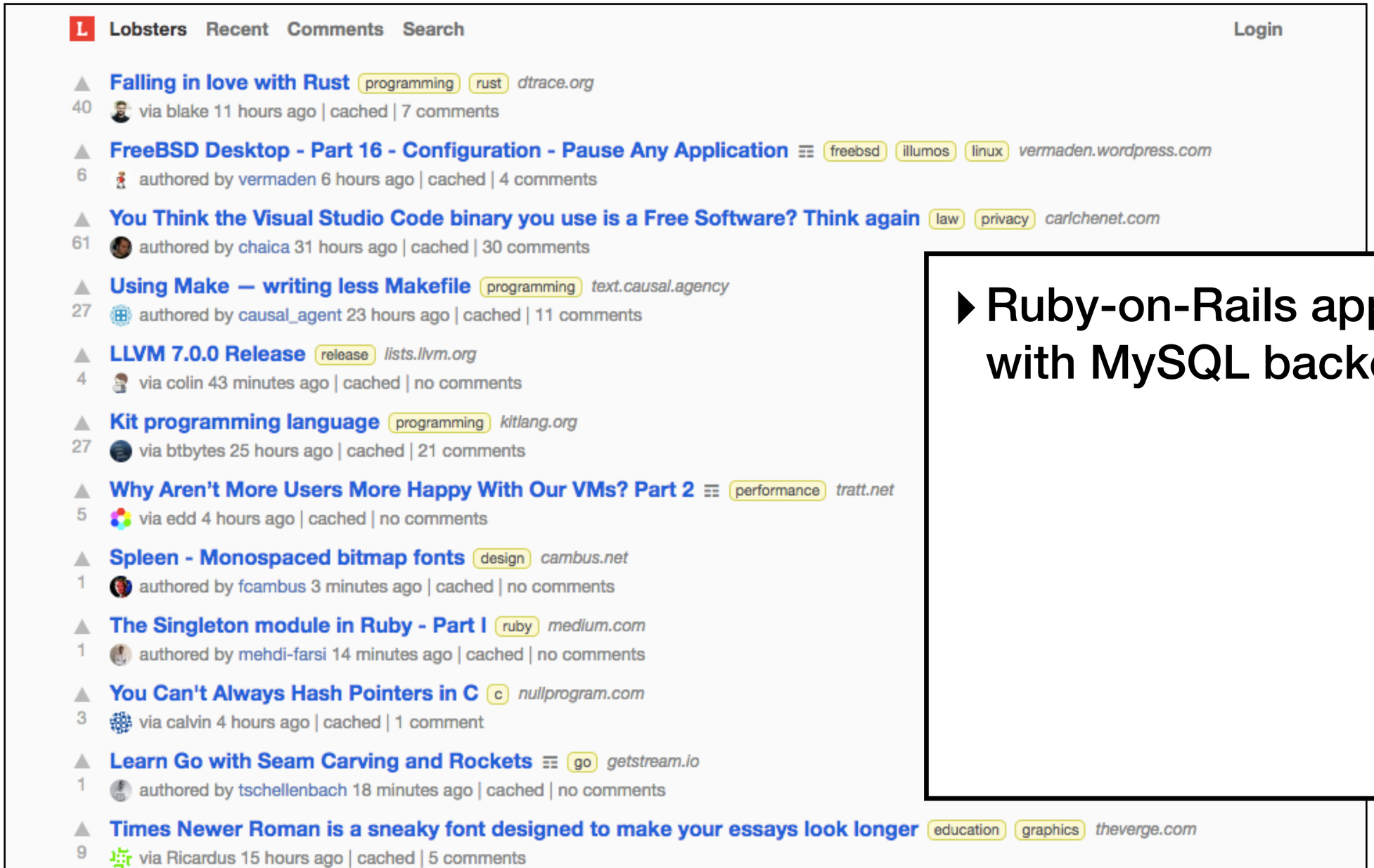

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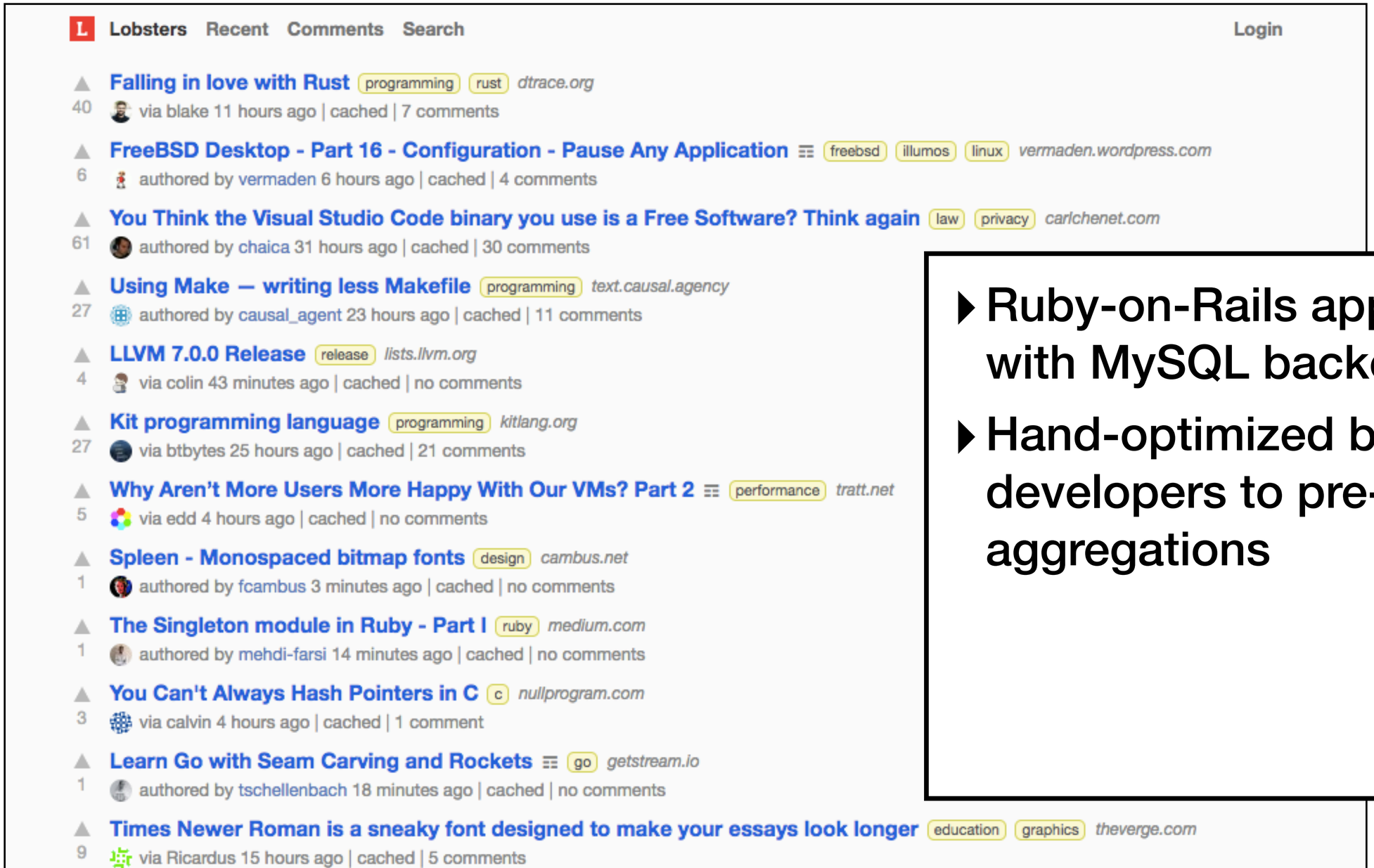

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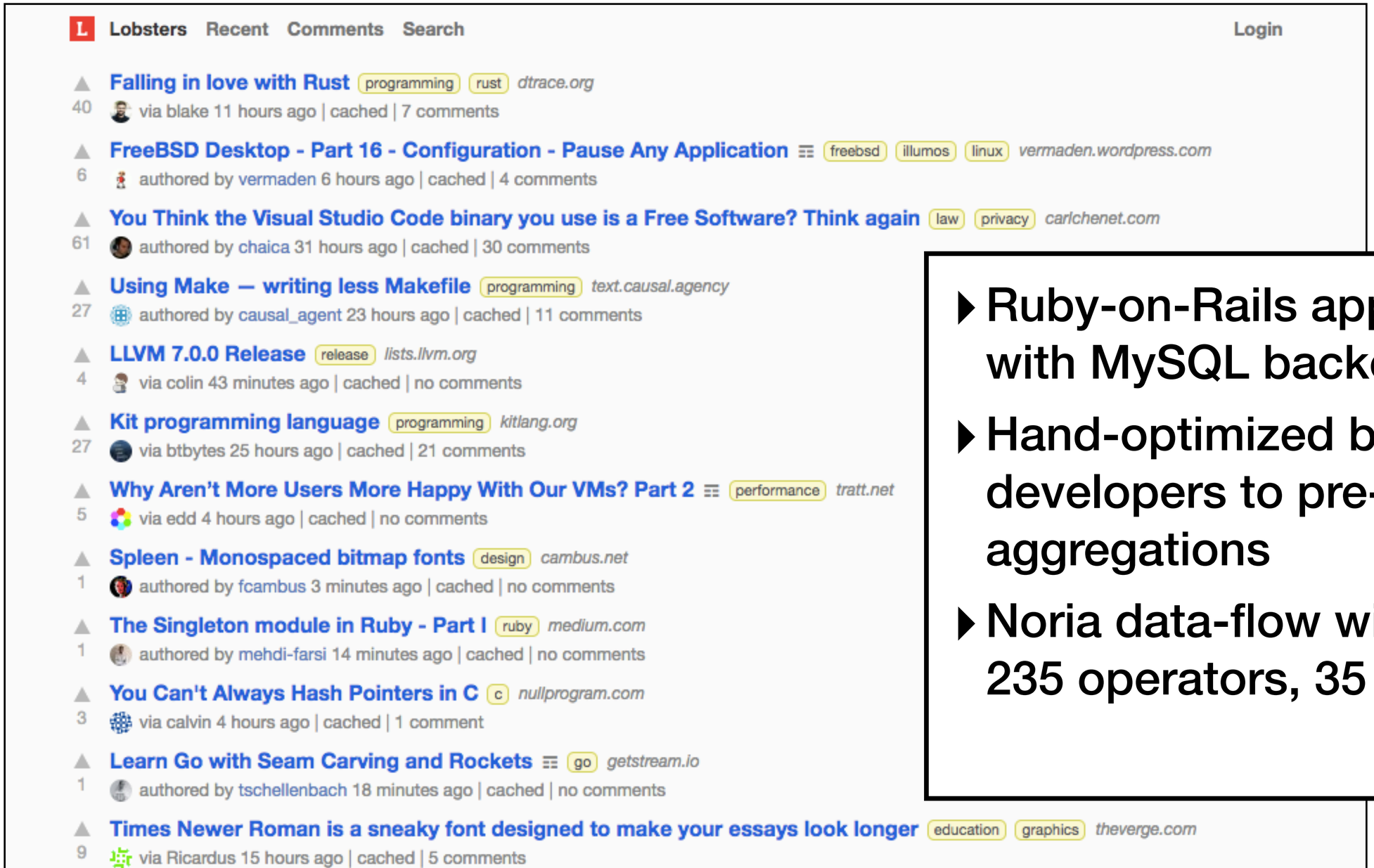

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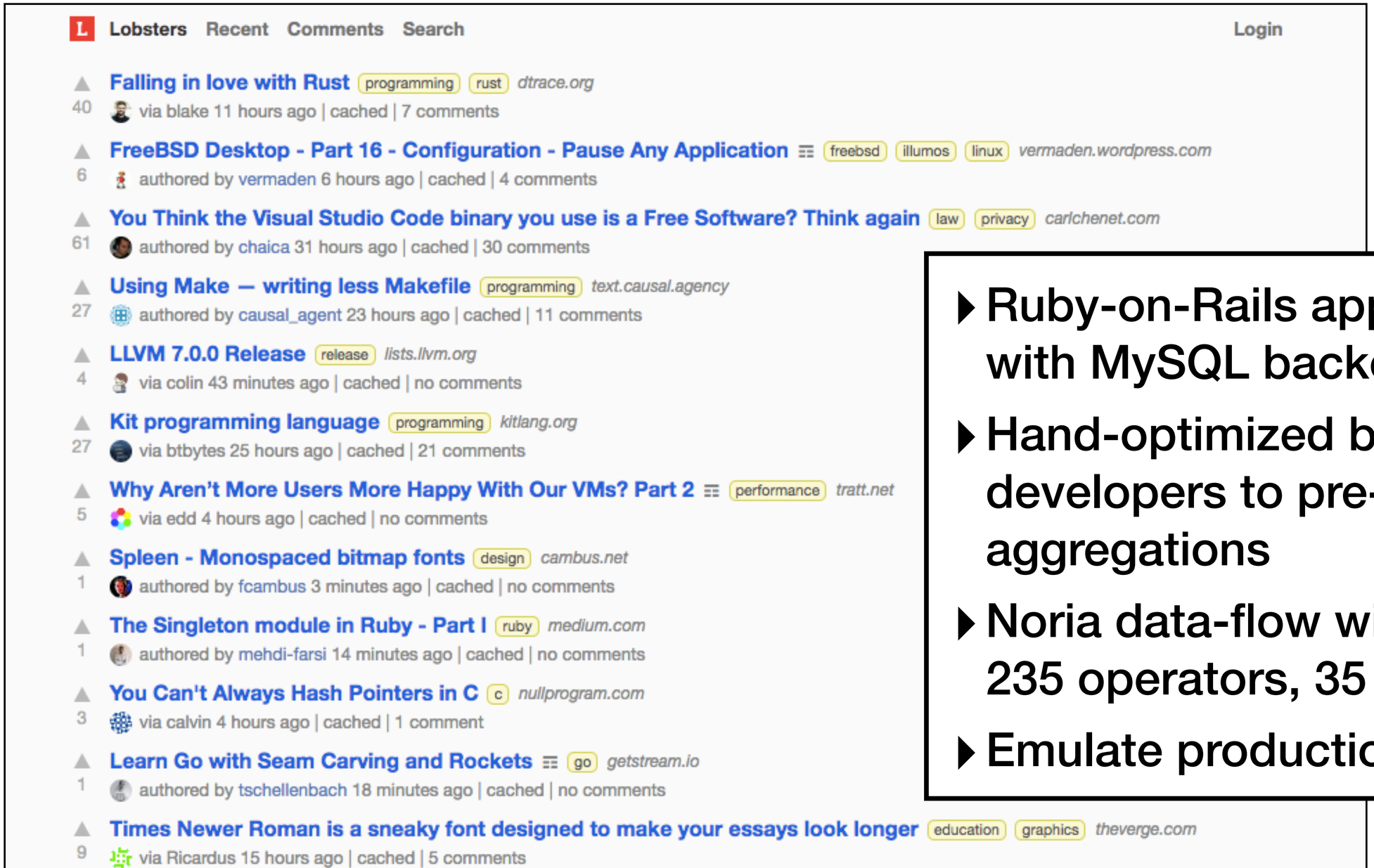

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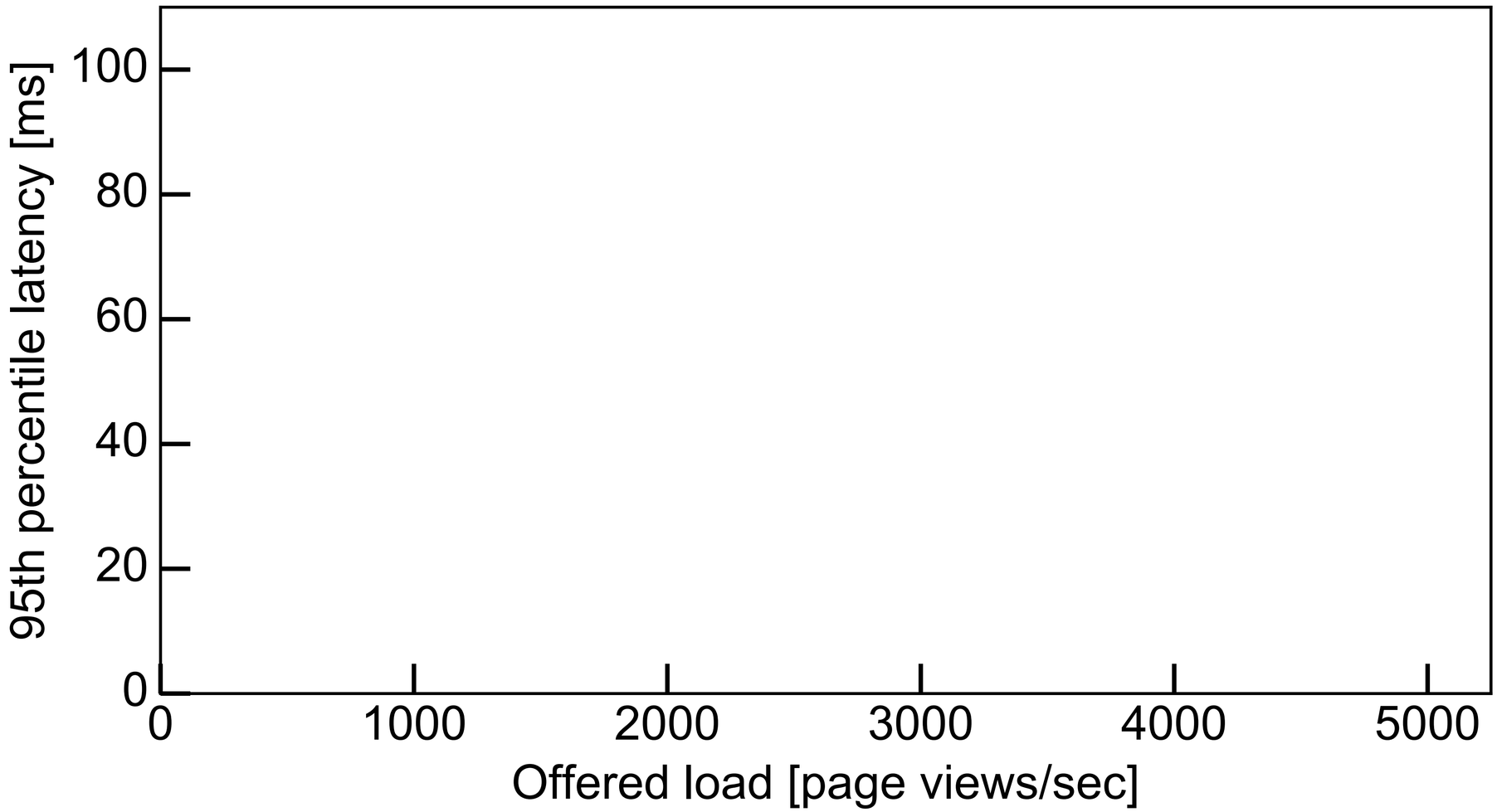
Can Noria improve Lobsters' performance?



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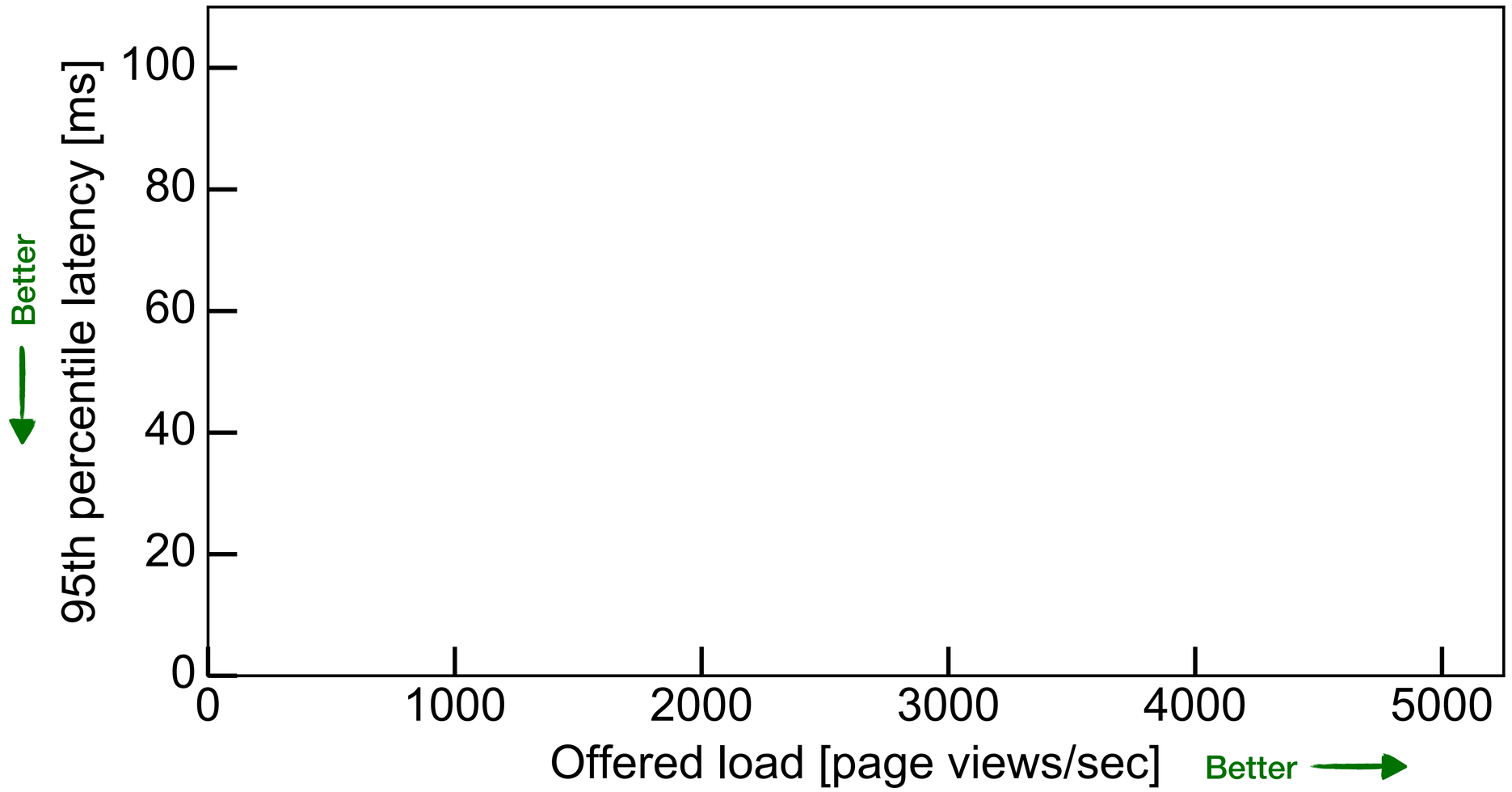

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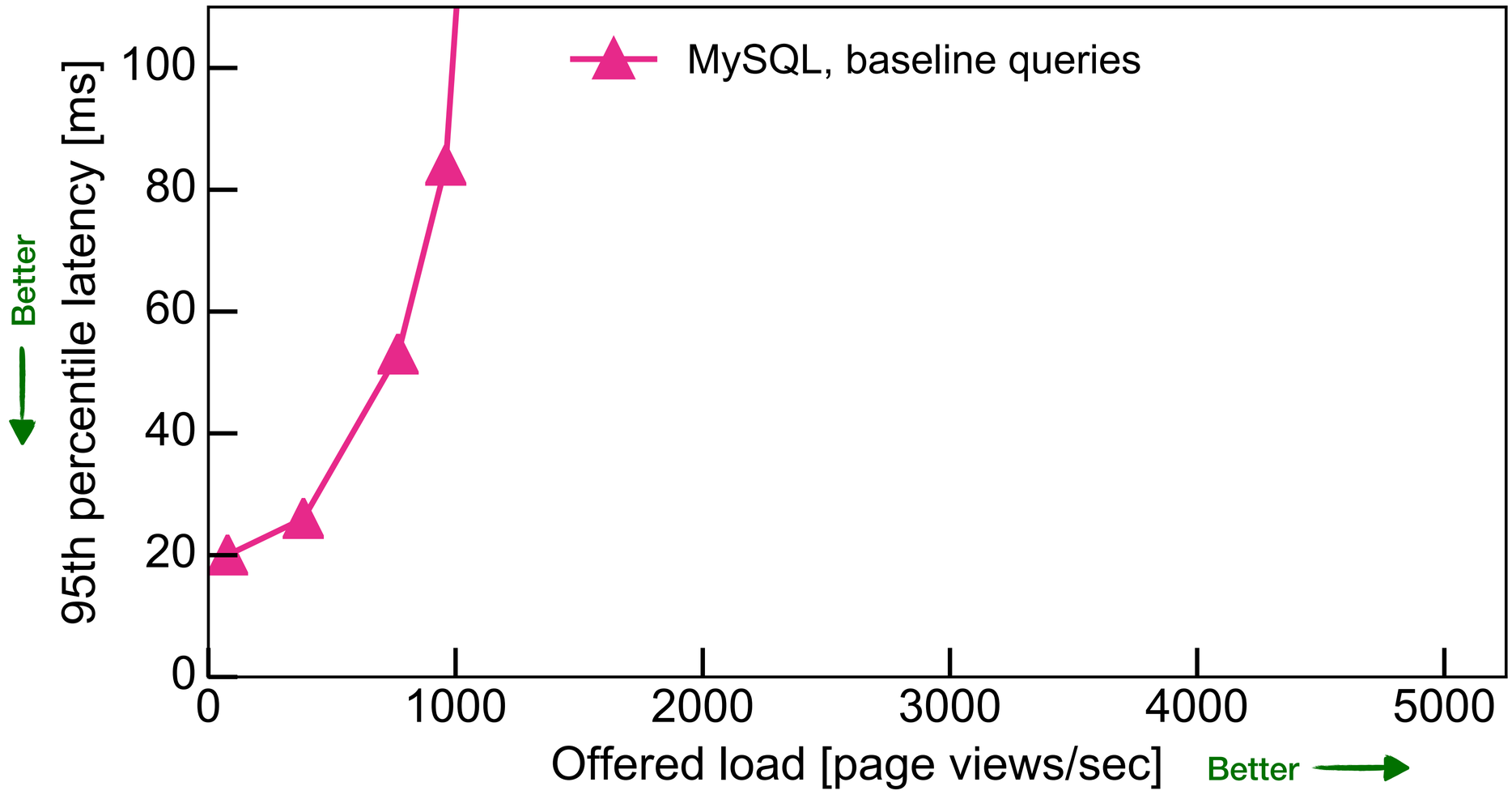

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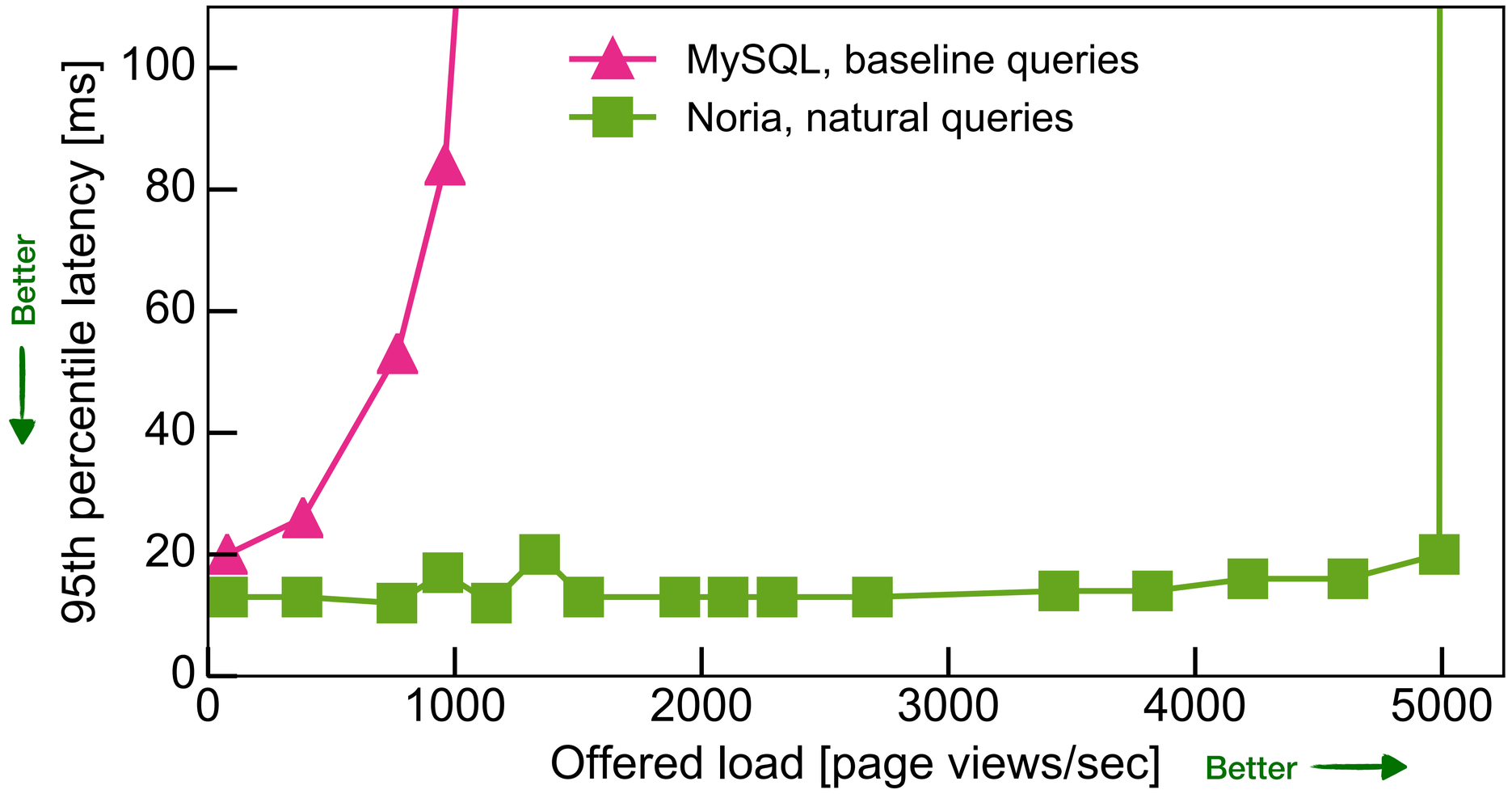

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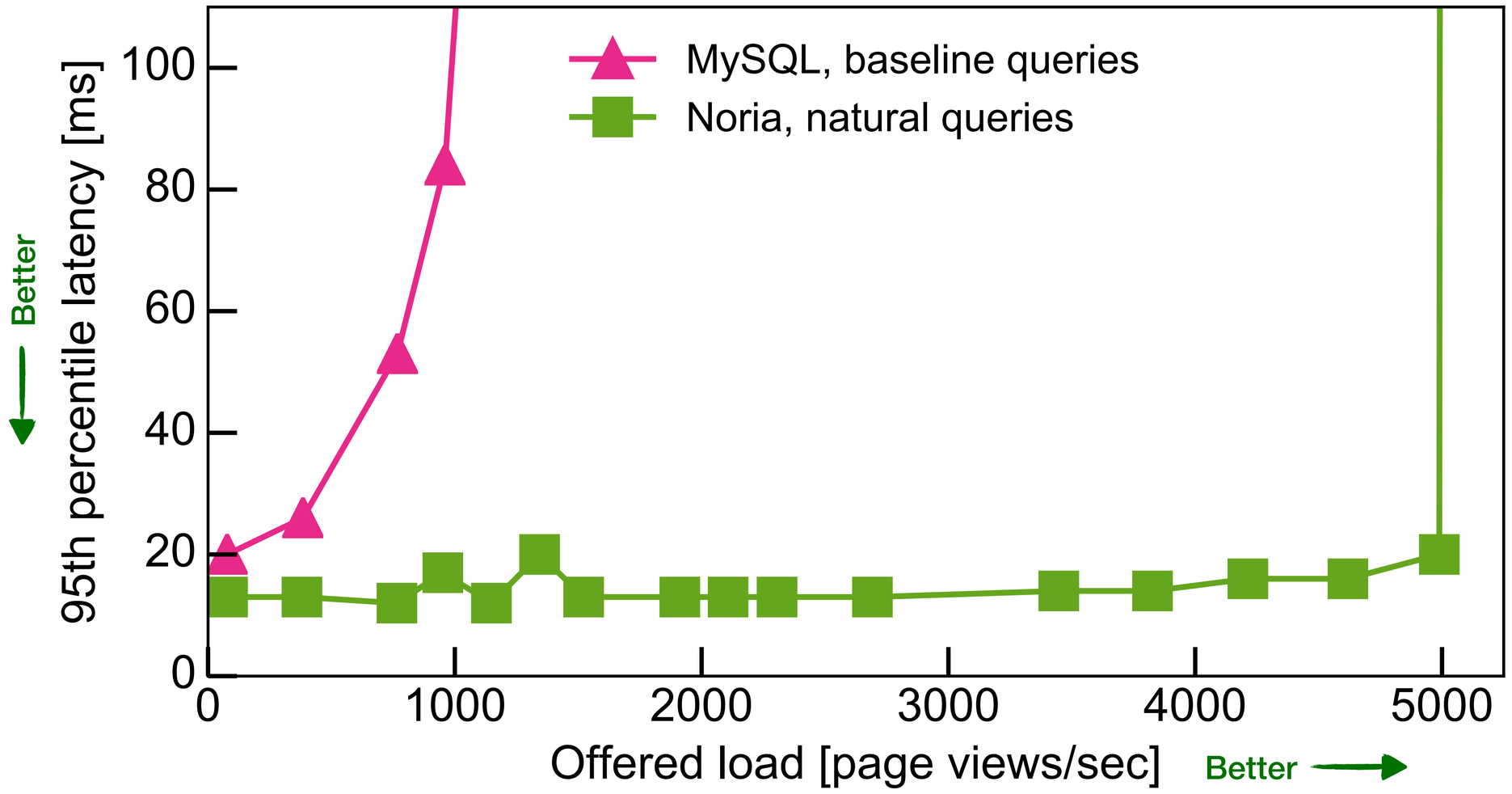
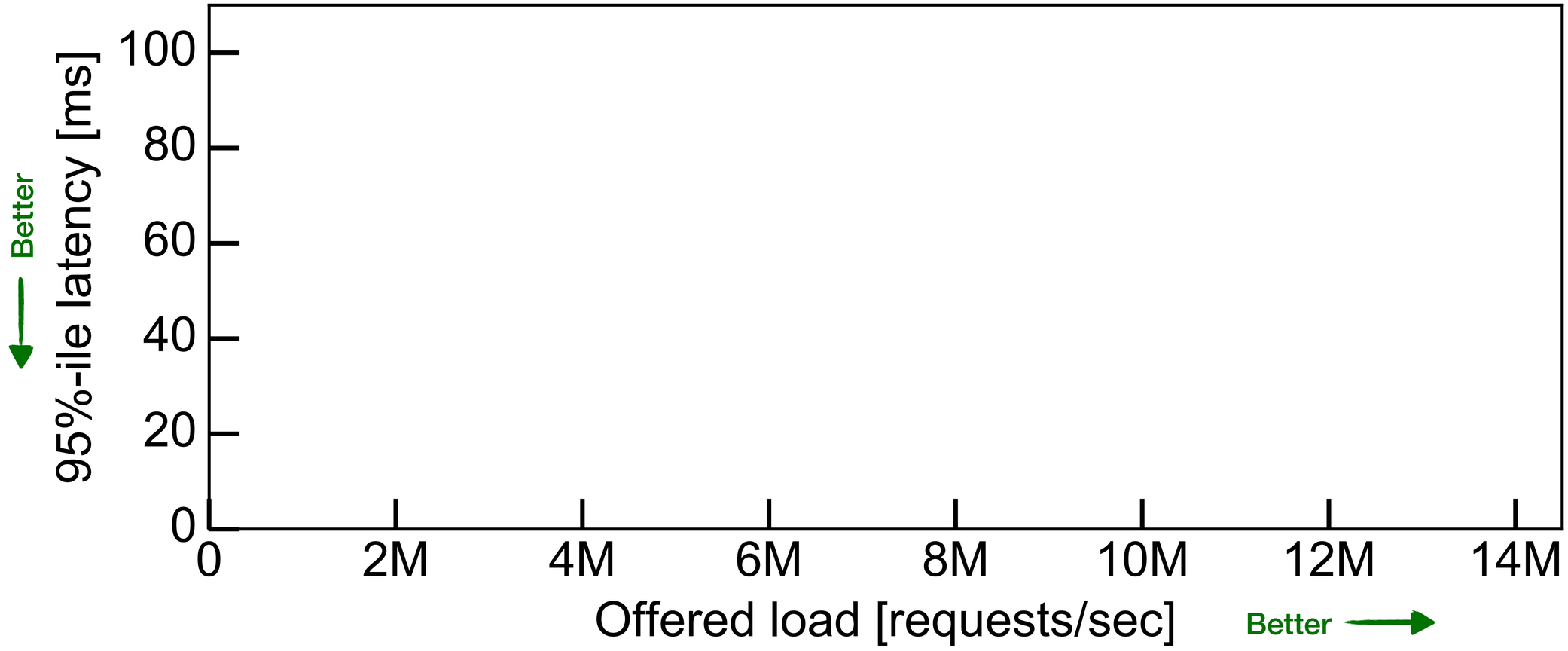


Noria with **natural queries** supports **5x** MySQL's throughput.

How does Noria compare to alternatives?



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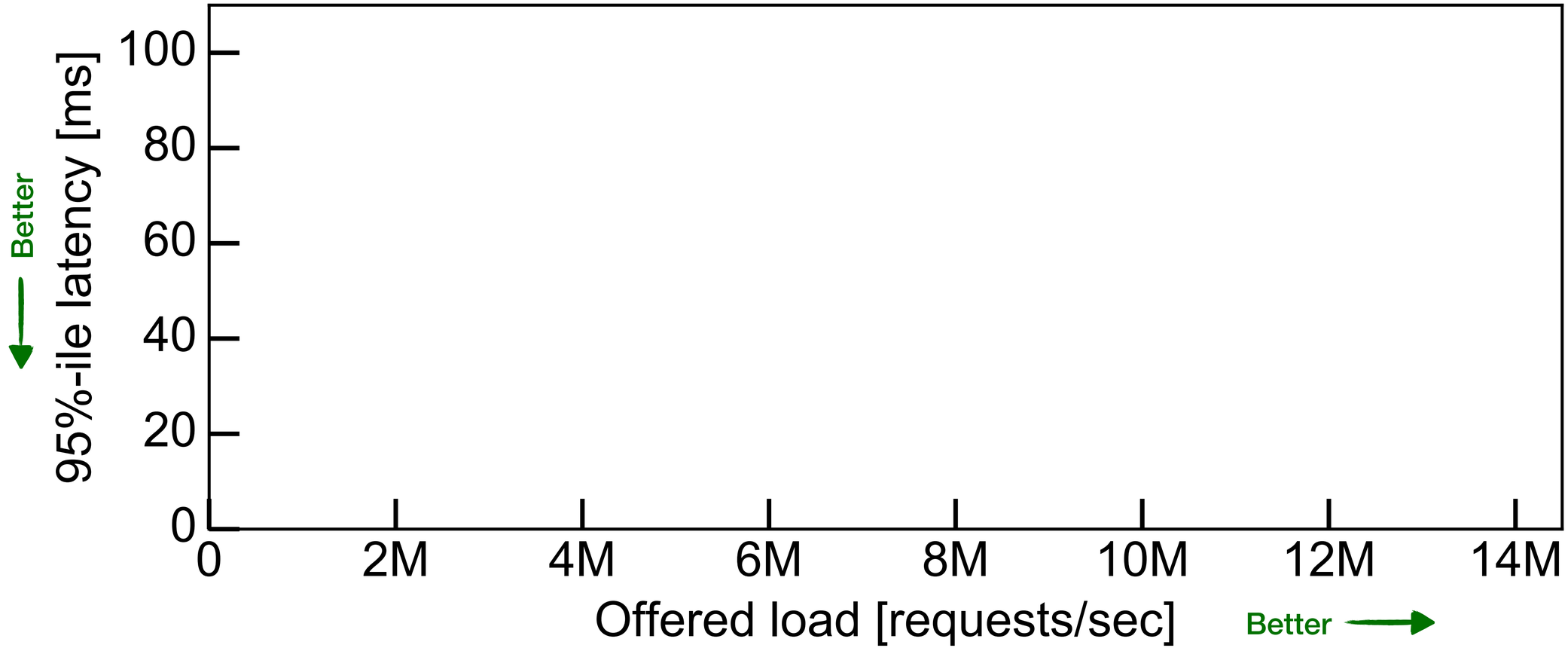

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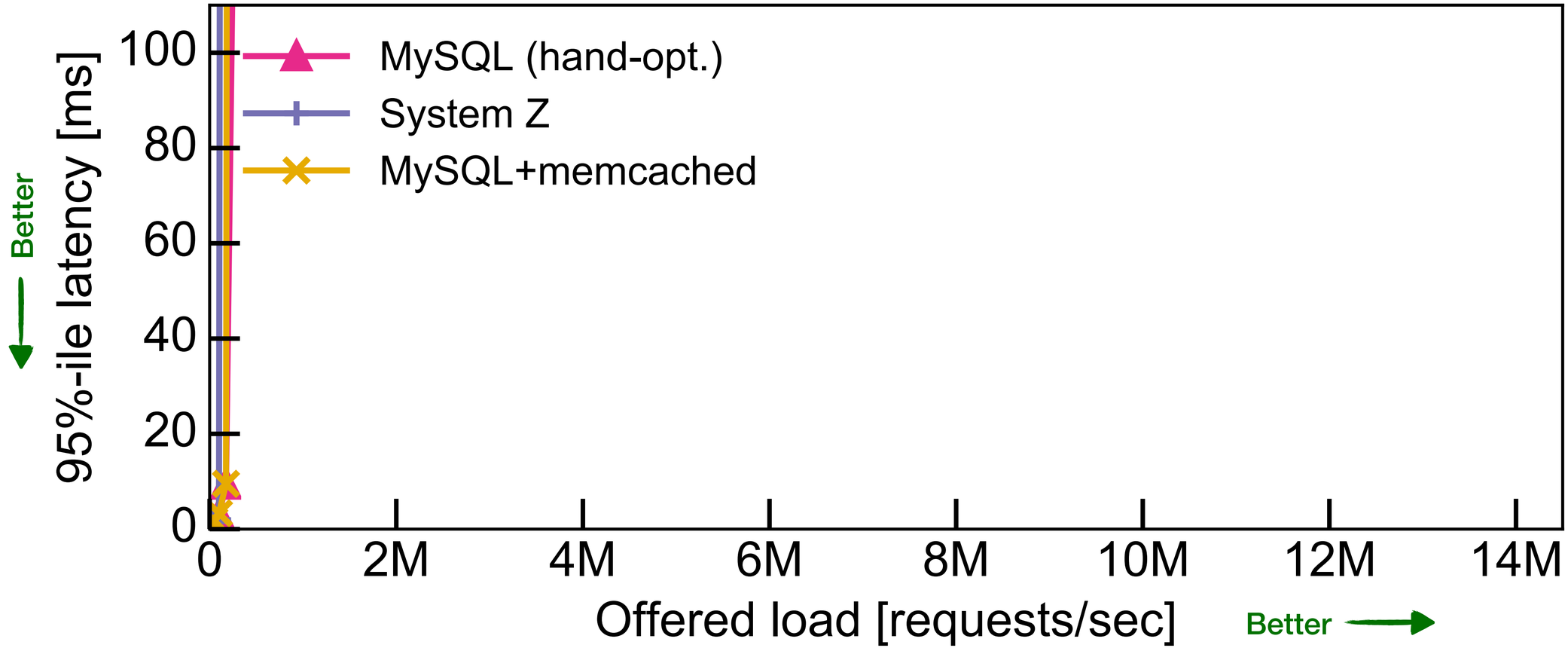

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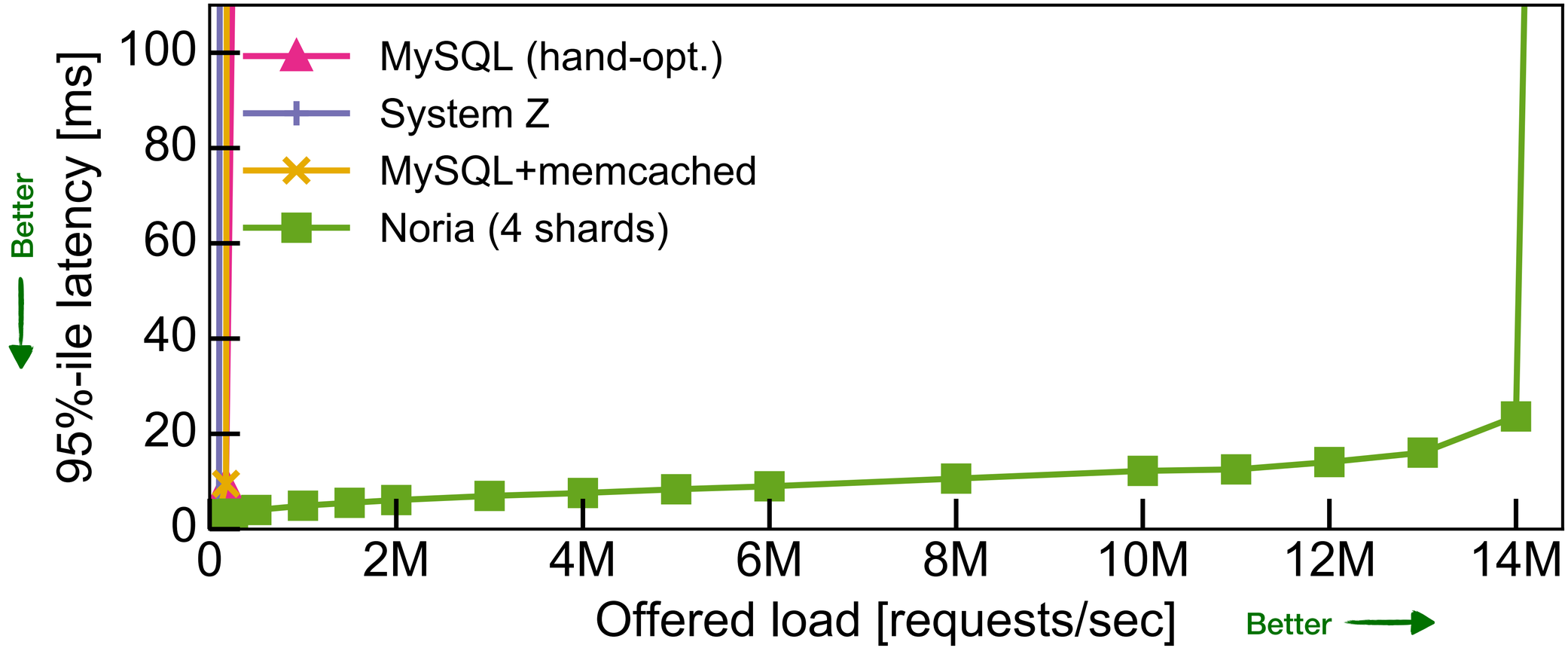

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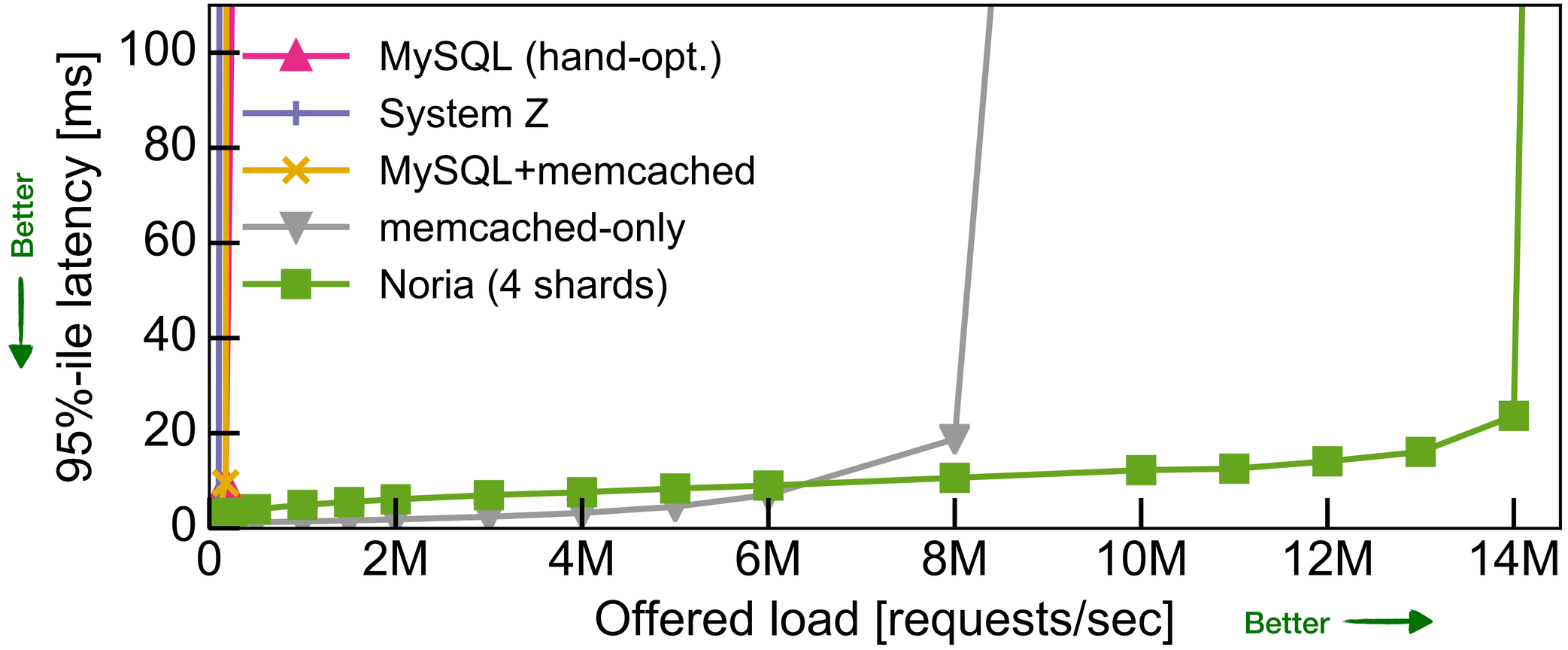

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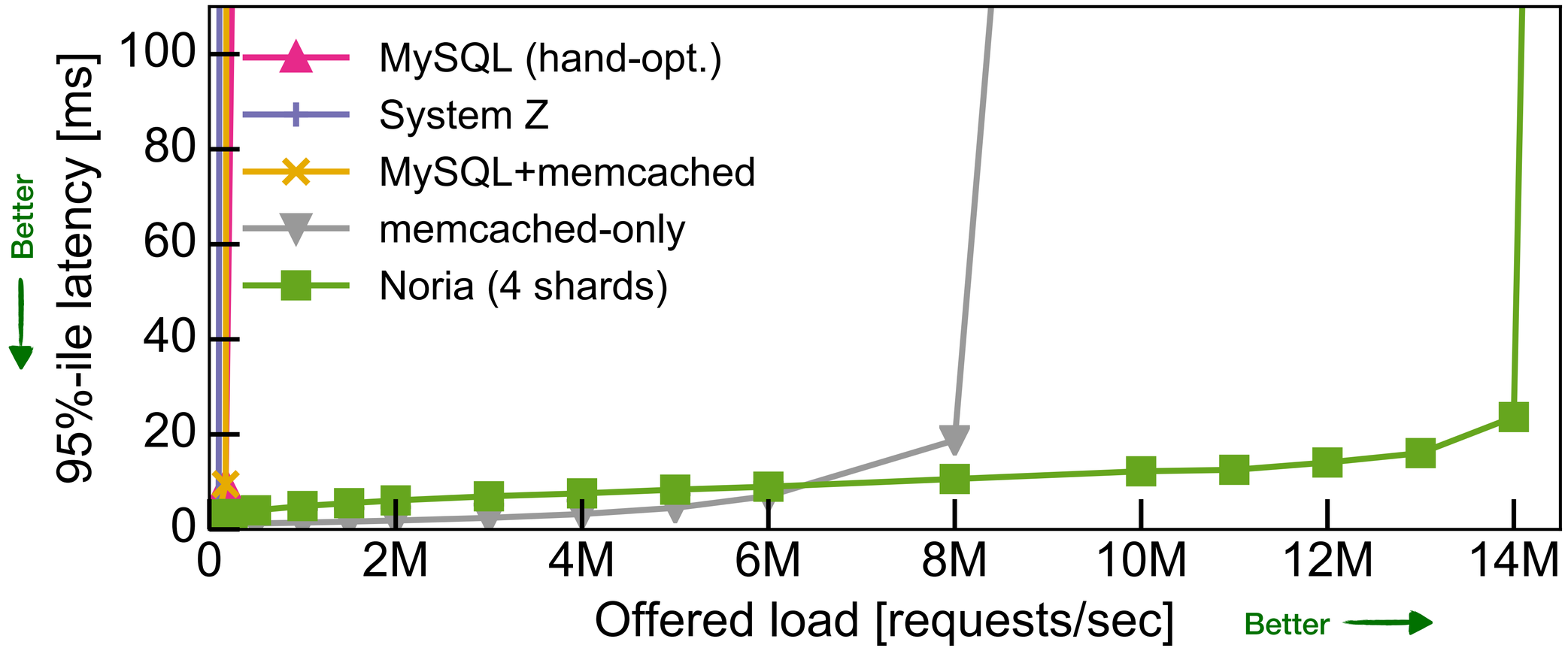

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Noria outperforms an **in-memory key-value store** and simplifies its interface.

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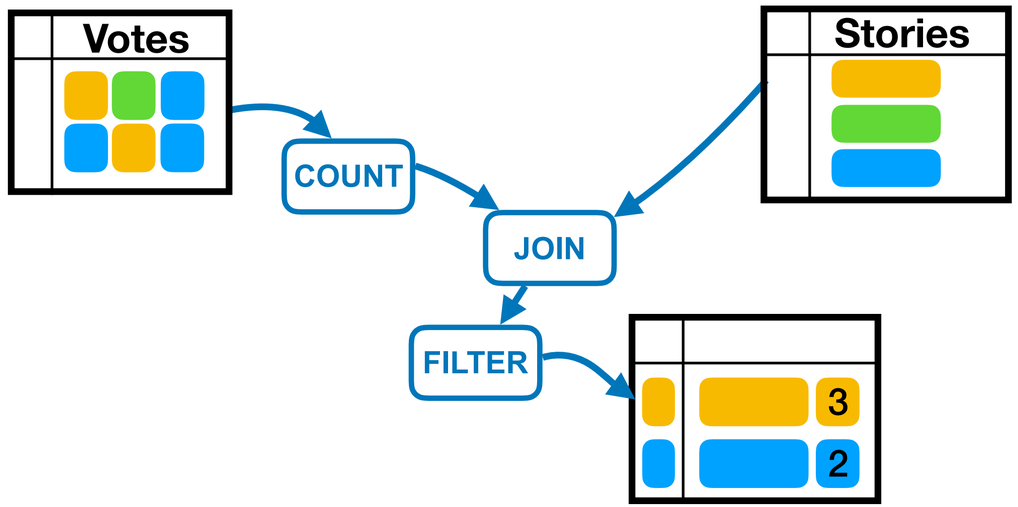

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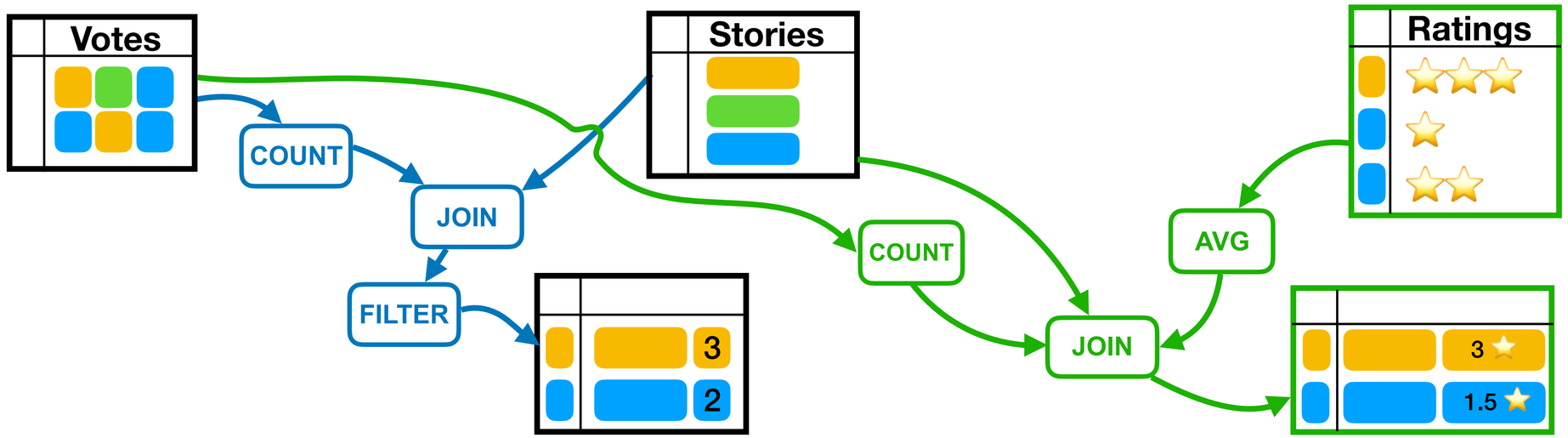

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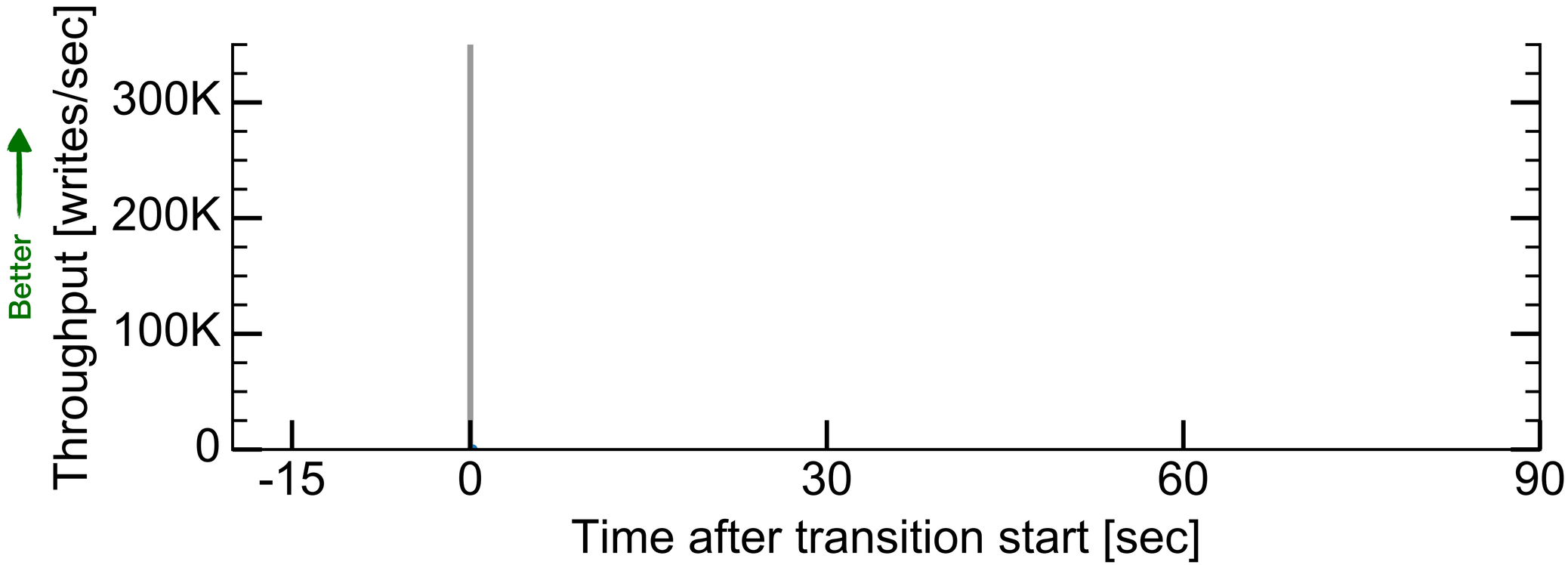

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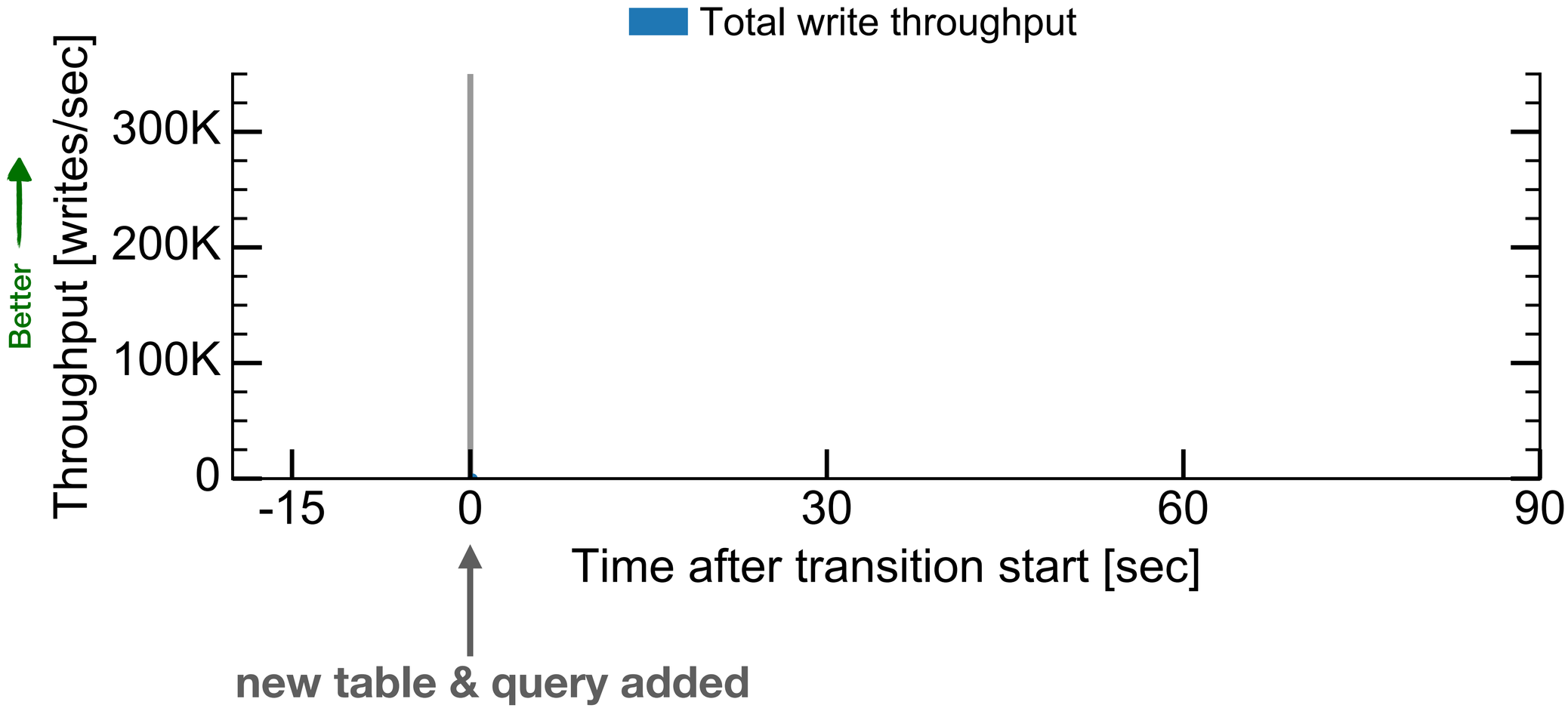

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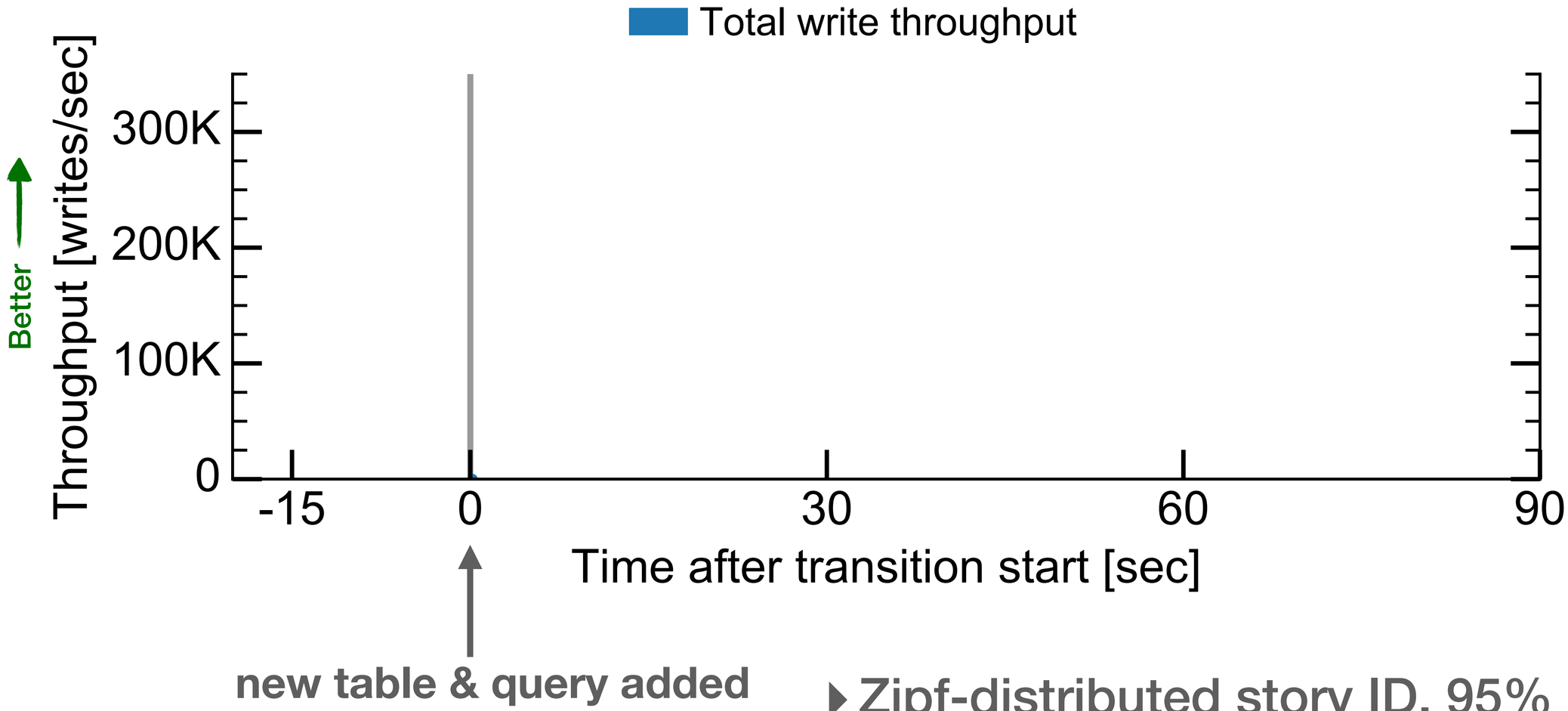

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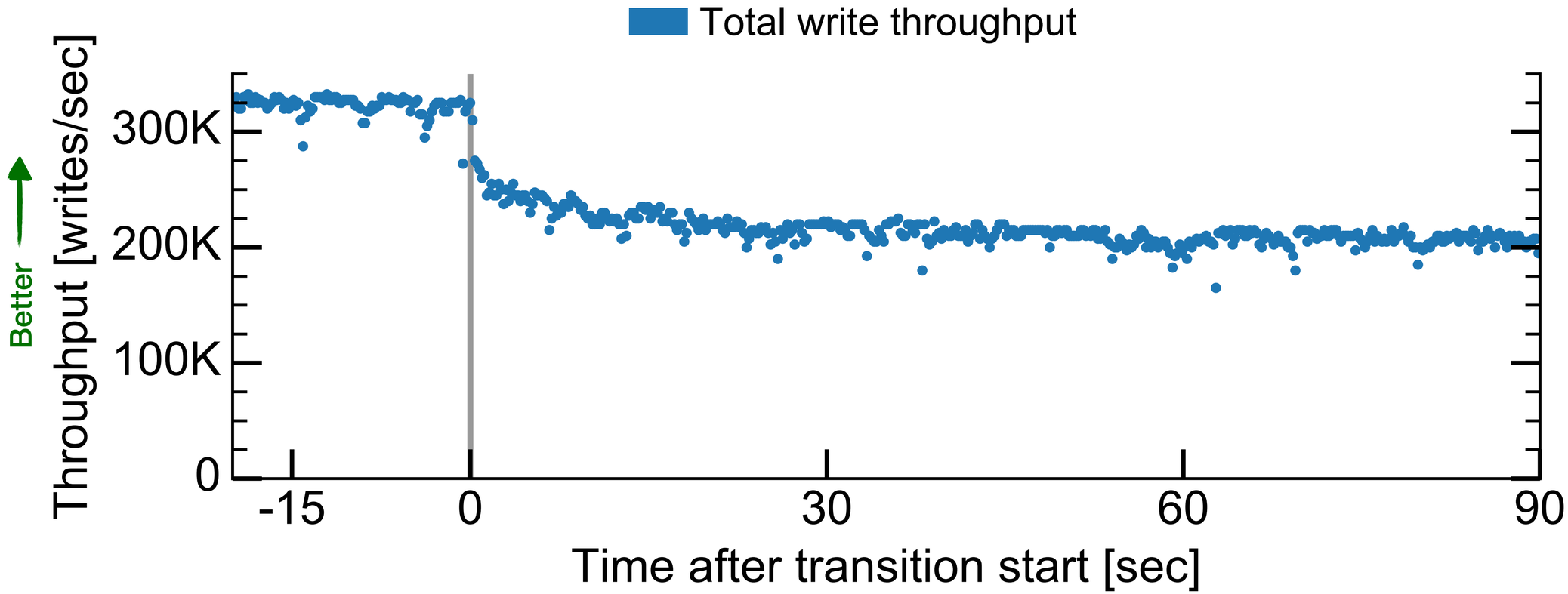

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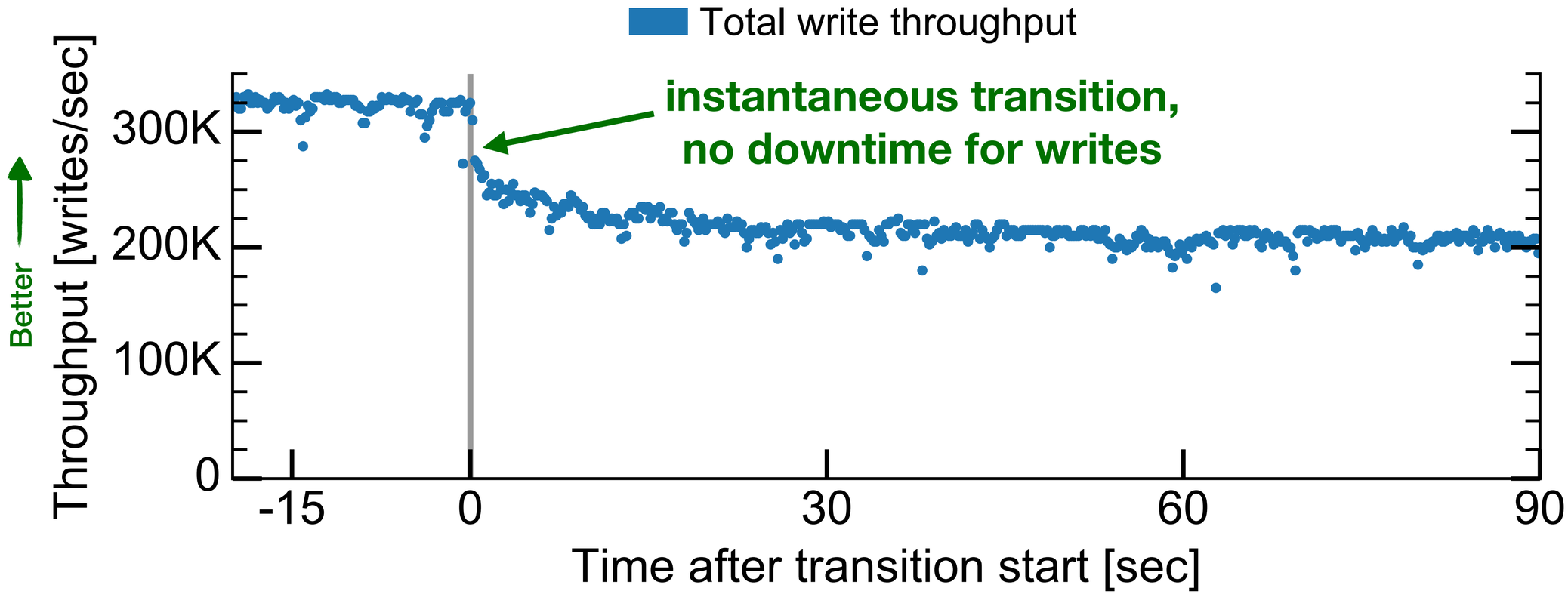

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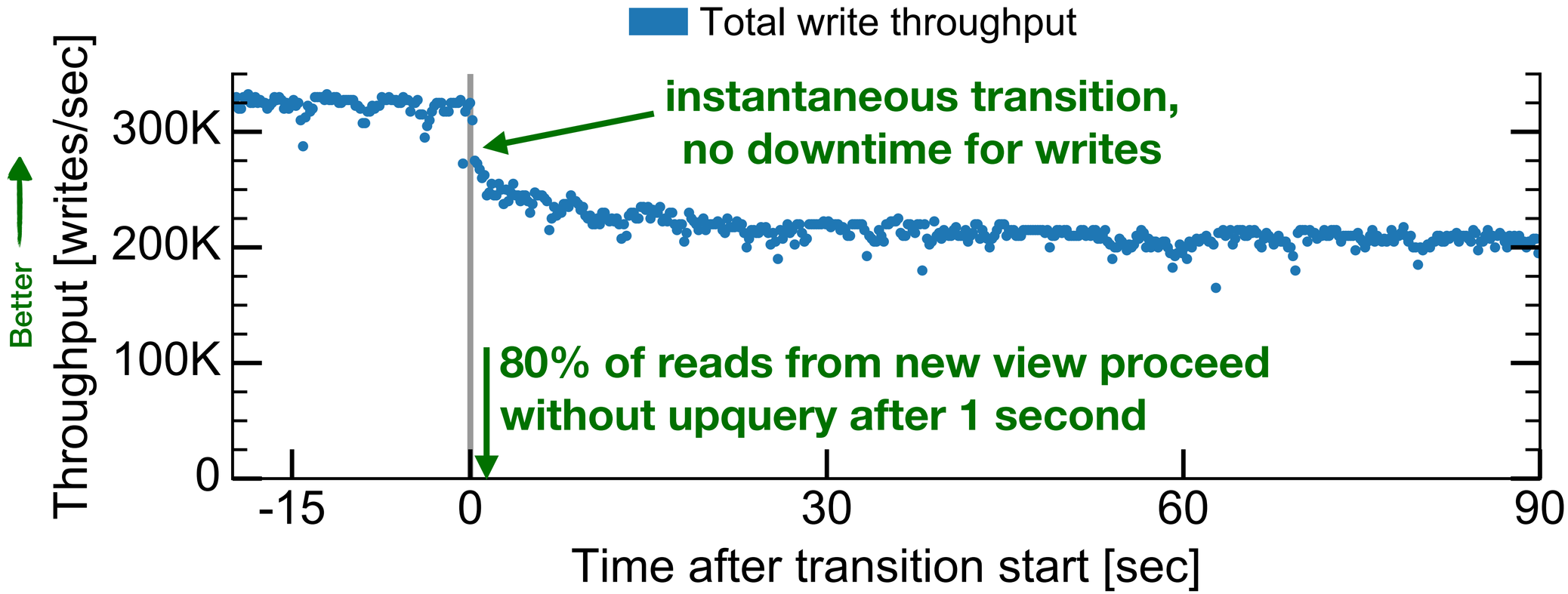

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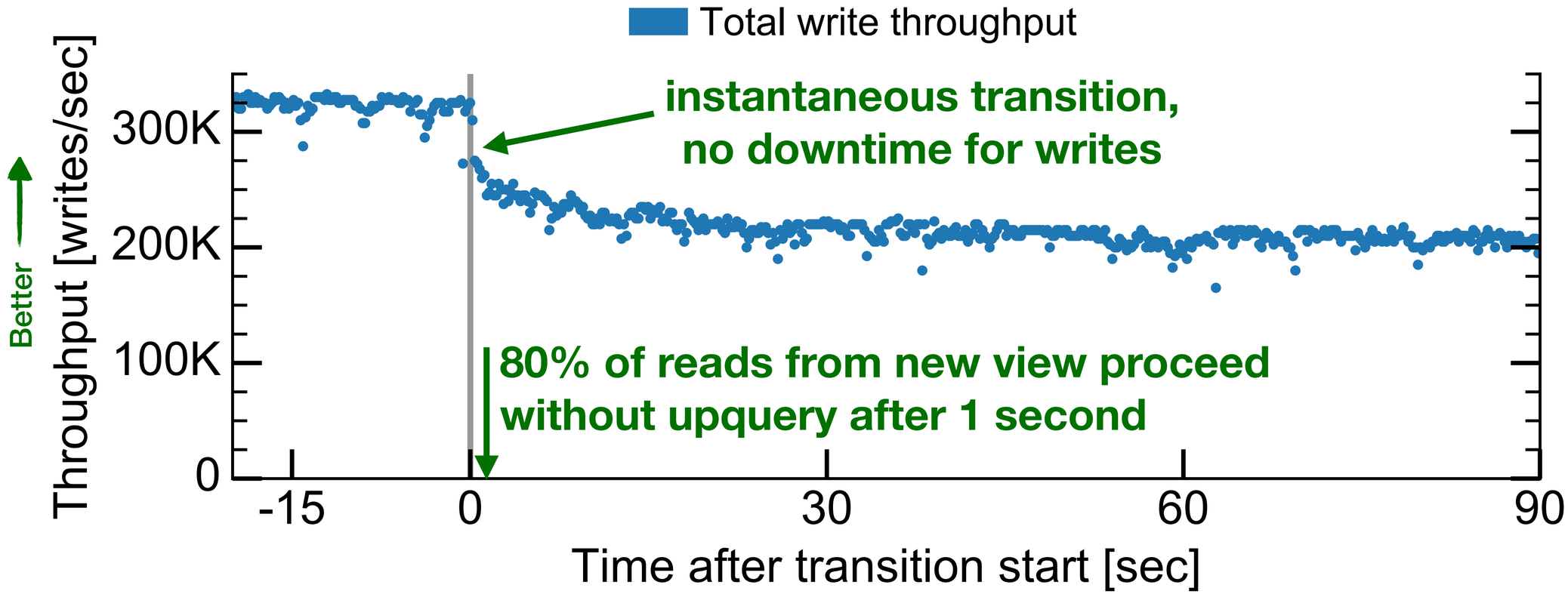

Noria **achieves downtime-free query change** with partial state.

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Noria — Summary

- New partially-stateful data-flow model.
- Noria: new web application backend based on data-flow.
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https://pdos.csail.mit.edu/noria

(see our demo at poster #37 today!)